

OSQP.jl

A Julia wrapper for the Operator Splitting QP solver

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joint work with Goran Banjac,

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JuMP Developers Meetup, 13 Jun 2017

Why quadratic programming?

AN ALGORITHM FOR QUADRATIC PROGRAMMING

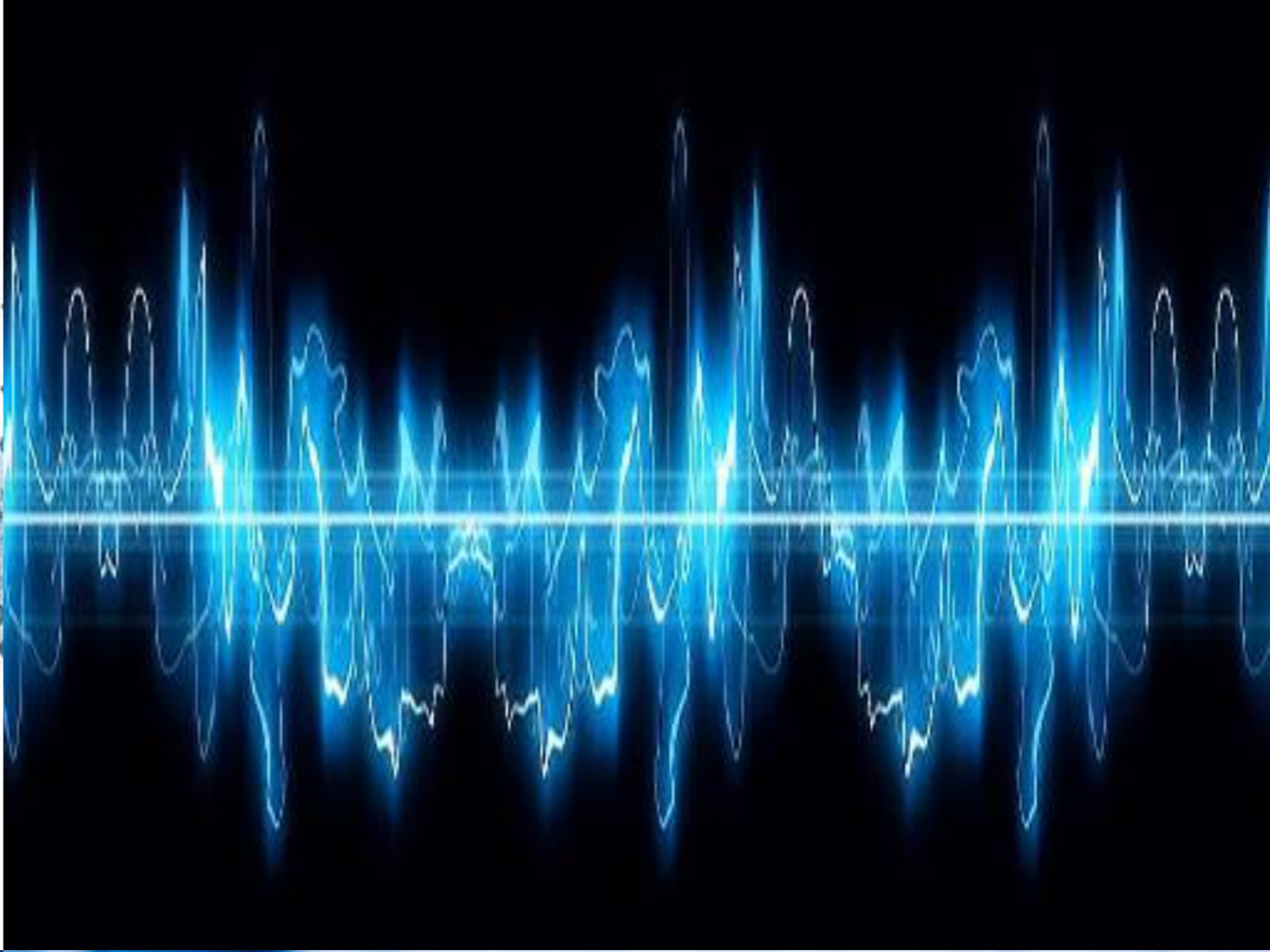
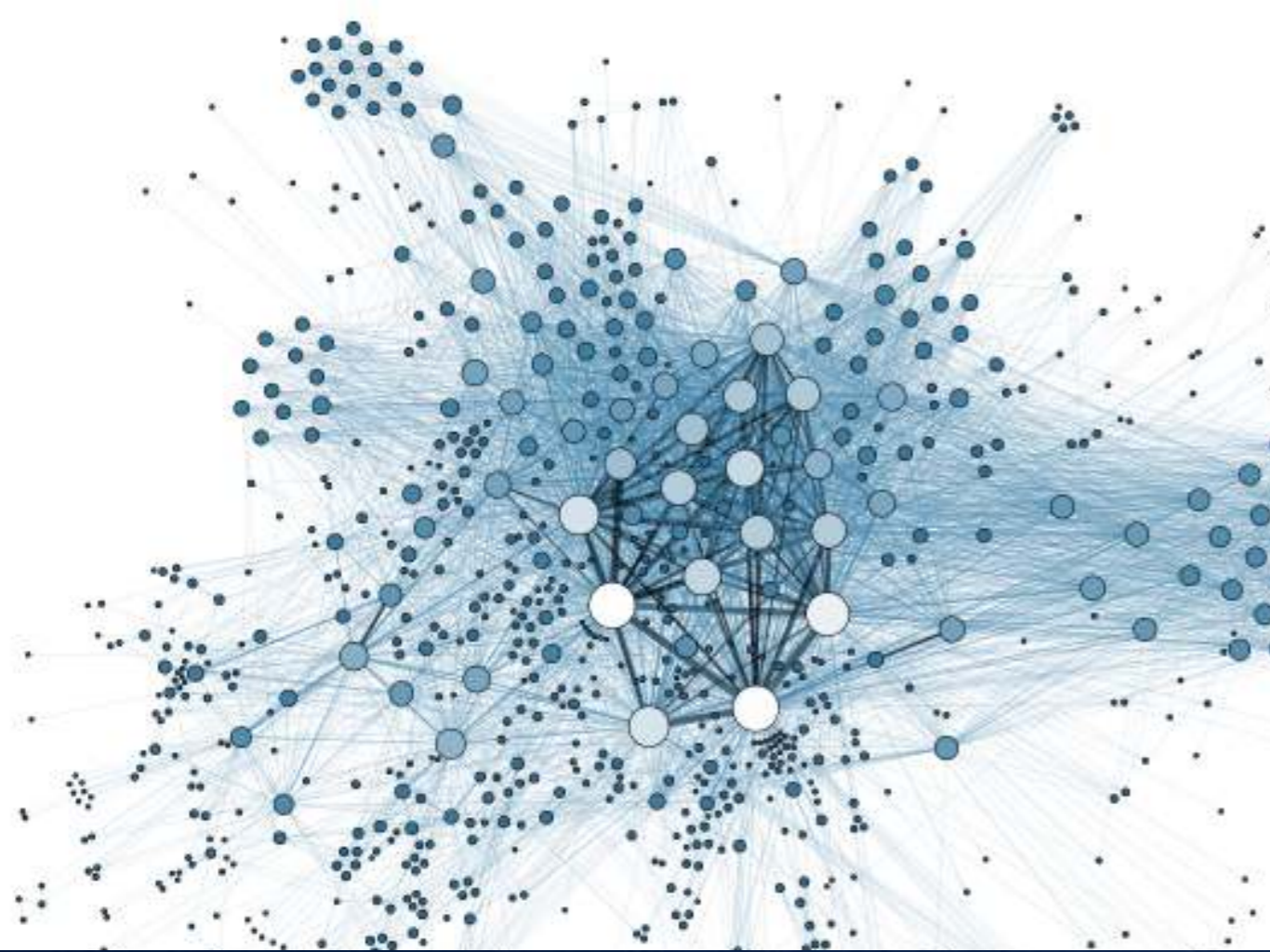
Marguerite Frank and Philip Wolfe¹
Princeton University

A finite iteration method for calculating the solution of quadratic programming problems is described. Extensions to more general non-linear problems are suggested.

1. INTRODUCTION

The problem of maximizing a concave quadratic function whose variables are subject to linear inequality constraints has been the subject of several recent studies, from both the computational side and the theoretical (see Bibliography). Our aim here has been to develop a method for solving this non-linear programming problem which should be particularly well adapted to high-speed machine computation.

March 1956!



First-order methods

Pros

Warm starting

Handle large-scale problems

Embeddable

Cons

Low accuracy solutions

Don't detect infeasibility

Problem data dependent

General Purpose QP Solver

Based on first-order
methods

Robust

Accurate

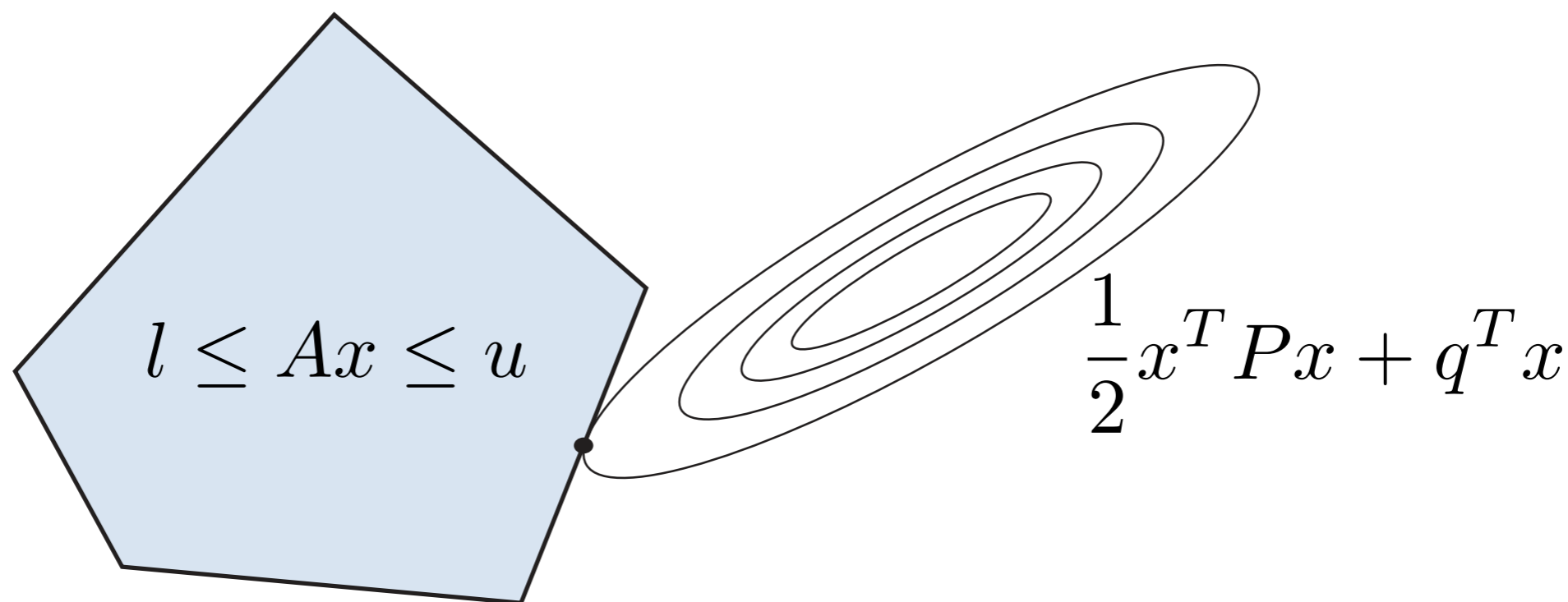
Detects
Infeasibility

The OSQP Solver

The problem

Quadratic Program

$$\begin{array}{ll} \text{minimize} & \frac{1}{2}x^T P x + q^T x \\ \text{subject to} & \underline{l} \leq A x \leq \underline{u} \end{array}$$



ADMM

$$\begin{array}{l} \text{minimize} \\ \text{subject to} \end{array} f(x) + g(x) \quad \longrightarrow \quad \begin{array}{l} \text{minimize} \\ \text{subject to} \end{array} f(\tilde{x}) + g(x) \\ \tilde{x} = x$$

ADMM

$$\text{minimize } f(x) + g(x) \longrightarrow \begin{array}{l} \text{minimize } f(\tilde{x}) + g(x) \\ \text{subject to } \tilde{x} = x \end{array}$$

$$\mathbf{1} \quad \tilde{x}^{k+1} \leftarrow \underset{\tilde{x}}{\operatorname{argmin}} \left(f(\tilde{x}) + \frac{\rho}{2} \left\| \tilde{x} - x^k + \frac{y^k}{\rho} \right\|^2 \right)$$

$$\mathbf{2} \quad x^{k+1} \leftarrow \underset{x}{\operatorname{argmin}} \left(g(x) + \frac{\rho}{2} \left\| x - \tilde{x}^{k+1} - \frac{y^k}{\rho} \right\|^2 \right)$$

$$\mathbf{3} \quad y^{k+1} \leftarrow y^k + \rho (\tilde{x}^{k+1} - x^{k+1})$$

How to split the QP?

$$\begin{aligned} &\text{minimize} && \frac{1}{2}x^T P x + q^T x \\ &\text{subject to} && Ax = z \\ &&& \underline{l} \leq z \leq \underline{u} \end{aligned}$$

$$\begin{aligned} &\text{minimize} && \frac{1}{2}\tilde{x}^T P \tilde{x} + q^T \tilde{x} + \mathcal{I}_{Ax=z}(\tilde{x}, \tilde{z}) + \mathcal{I}_{[l,u]}(z) \\ &\text{subject to} && (\tilde{x}, \tilde{z}) = (x, z) \end{aligned}$$

How to split the QP?

$$\begin{array}{ll} \text{minimize} & \frac{1}{2}x^T P x + q^T x \\ \text{subject to} & Ax = z \\ & \underline{l} \leq z \leq \underline{u} \end{array} \quad f$$

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How to split the QP?

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$$\begin{array}{ll} \text{minimize} & \frac{1}{2}\tilde{x}^T P \tilde{x} + q^T \tilde{x} + \mathcal{I}_{Ax=z}(\tilde{x}, \tilde{z}) + \mathcal{I}_{[l,u]}(z) \\ \text{subject to} & (\tilde{x}, \tilde{z}) = (x, z) \end{array} \quad \begin{array}{l} f \\ g \end{array}$$

OSQP

Problem

$$\begin{aligned} &\text{minimize} && \frac{1}{2}x^T P x + q^T x \\ &\text{subject to} && l \leq A x \leq u \end{aligned}$$

Algorithm

$$\begin{aligned} \mathbf{1} & \left\{ \begin{aligned} (x^{k+1}, \nu^{k+1}) &\leftarrow \text{solve} \begin{bmatrix} P + \sigma I & A^T \\ A & -\frac{1}{\rho} I \end{bmatrix} \begin{bmatrix} x^{k+1} \\ \nu^{k+1} \end{bmatrix} = \begin{bmatrix} \sigma x^k - q \\ z^k - \frac{1}{\rho} y^k \end{bmatrix} \\ \tilde{z}^{k+1} &\leftarrow z^k + \frac{1}{\rho} (\nu^{k+1} - y^k) \end{aligned} \right. \\ \mathbf{2} & \left\{ z^{k+1} \leftarrow \Pi \left(\tilde{z}^{k+1} + \frac{1}{\rho} y^k \right) \right. \\ \mathbf{3} & \left\{ y^{k+1} \leftarrow y^k + \rho (\tilde{z}^{k+1} - z^{k+1}) \right. \end{aligned}$$

OSQP

Problem

$$\begin{aligned} &\text{minimize} && \frac{1}{2}x^T P x + q^T x \\ &\text{subject to} && l \leq A x \leq u \end{aligned}$$

Algorithm

Linear system
solve

$$(x^{k+1}, \nu^{k+1}) \leftarrow \text{solve} \begin{bmatrix} P + \sigma I & A^T \\ A & -\frac{1}{\rho} I \end{bmatrix} \begin{bmatrix} x^{k+1} \\ \nu^{k+1} \end{bmatrix} = \begin{bmatrix} \sigma x^k - q \\ z^k - \frac{1}{\rho} y^k \end{bmatrix}$$

$$\tilde{z}^{k+1} \leftarrow z^k + \frac{1}{\rho} (\nu^{k+1} - y^k)$$

$$z^{k+1} \leftarrow \Pi \left(\tilde{z}^{k+1} + \frac{1}{\rho} y^k \right)$$

$$y^{k+1} \leftarrow y^k + \rho (\tilde{z}^{k+1} - z^{k+1})$$

OSQP

Problem

$$\begin{aligned} &\text{minimize} && \frac{1}{2}x^T P x + q^T x \\ &\text{subject to} && l \leq A x \leq u \end{aligned}$$

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Easy
operations

$$\tilde{z}^{k+1} \leftarrow z^k + \frac{1}{\rho} (\nu^{k+1} - y^k)$$

$$z^{k+1} \leftarrow \Pi \left(\tilde{z}^{k+1} + \frac{1}{\rho} y^k \right)$$

$$y^{k+1} \leftarrow y^k + \rho (\tilde{z}^{k+1} - z^{k+1})$$

OSQP

Problem

$$\begin{aligned} &\text{minimize} && \frac{1}{2}x^T P x + q^T x \\ &\text{subject to} && l \leq A x \leq u \end{aligned}$$

Algorithm

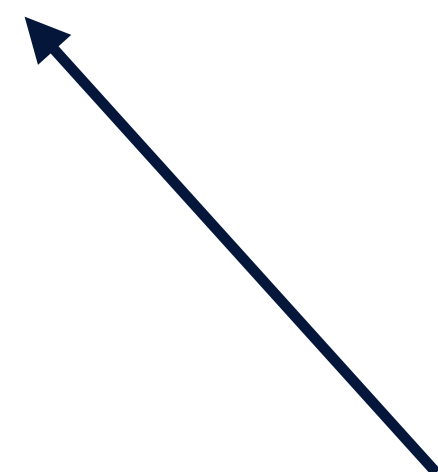
Linear system
solve

$$(x^{k+1}, \nu^{k+1}) \leftarrow \text{solve} \begin{bmatrix} P + \sigma I & A^T \\ A & -\frac{1}{\rho} I \end{bmatrix} \begin{bmatrix} x^{k+1} \\ \nu^{k+1} \end{bmatrix} = \begin{bmatrix} \sigma x^k - q \\ z^k - \frac{1}{\rho} y^k \end{bmatrix}$$

Easy
operations

$$\begin{aligned} \tilde{z}^{k+1} &\leftarrow z^k + \frac{1}{\rho} (\nu^{k+1} - y^k) \\ z^{k+1} &\leftarrow \Pi \left(\tilde{z}^{k+1} + \frac{1}{\rho} y^k \right) \\ y^{k+1} &\leftarrow y^k + \rho (\tilde{z}^{k+1} - z^{k+1}) \end{aligned}$$

Factorization
caching



OSQP

Problem

$$\begin{aligned} &\text{minimize} && \frac{1}{2}x^T P x + q^T x \\ &\text{subject to} && l \leq A x \leq u \end{aligned}$$

Algorithm

Linear system
solve

$$(x^{k+1}, \nu^{k+1}) \leftarrow \text{solve} \begin{bmatrix} P + \sigma I & A^T \\ A & -\frac{1}{\rho} I \end{bmatrix} \begin{bmatrix} x^{k+1} \\ \nu^{k+1} \end{bmatrix} = \begin{bmatrix} \sigma x^k - q \\ z^k - \frac{1}{\rho} y^k \end{bmatrix}$$

Easy
operations

$$\begin{aligned} \tilde{z}^{k+1} &\leftarrow z^k + \frac{1}{\rho} (\nu^{k+1} - y^k) \\ z^{k+1} &\leftarrow \Pi \left(\tilde{z}^{k+1} + \frac{1}{\rho} y^k \right) \\ y^{k+1} &\leftarrow y^k + \rho (\tilde{z}^{k+1} - z^{k+1}) \end{aligned}$$

Warm
starting

Factorization
caching

OSQP

Problem

$$\begin{aligned} &\text{minimize} && \frac{1}{2}x^T P x + q^T x \\ &\text{subject to} && l \leq A x \leq u \end{aligned}$$

Algorithm

Linear system
solve

$$(x^{k+1}, \nu^{k+1}) \leftarrow \text{solve} \begin{bmatrix} P + \sigma I & A^T \\ A & -\frac{1}{\rho} I \end{bmatrix} \begin{bmatrix} x^{k+1} \\ \nu^{k+1} \end{bmatrix} = \begin{bmatrix} \sigma x^k - q \\ z^k - \frac{1}{\rho} y^k \end{bmatrix}$$

Easy
operations

$$\begin{aligned} \tilde{z}^{k+1} &\leftarrow z^k + \frac{1}{\rho} (\nu^{k+1} - y^k) \\ z^{k+1} &\leftarrow \Pi \left(\tilde{z}^{k+1} + \frac{1}{\rho} y^k \right) \\ y^{k+1} &\leftarrow y^k + \rho (\tilde{z}^{k+1} - z^{k+1}) \end{aligned}$$

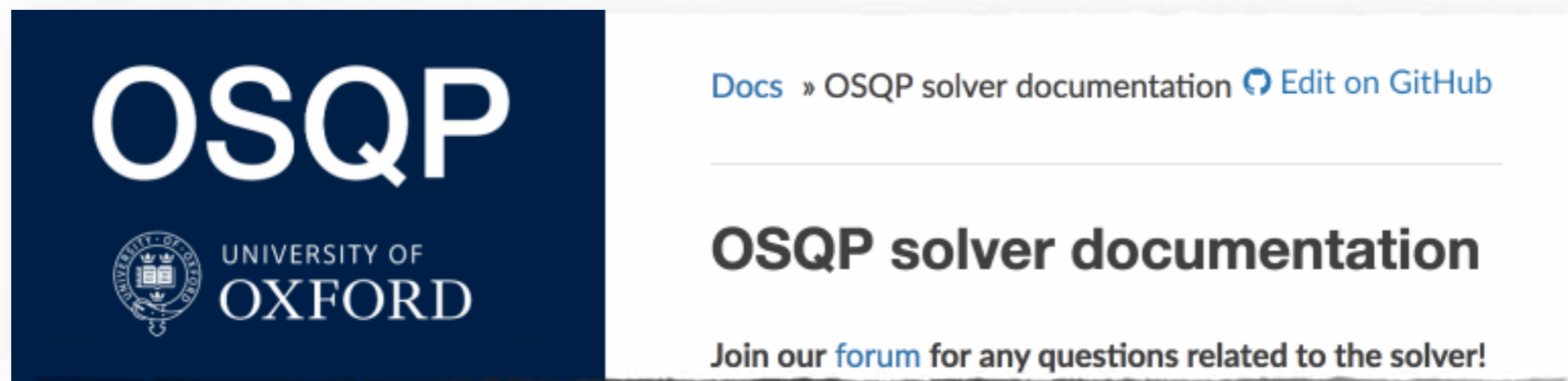
Solution
polishing

Warm
starting

Factorization
caching

OSQP

`osqp.readthedocs.io`



The screenshot shows the OSQP solver documentation page. On the left is a dark blue banner with the OSQP logo and the University of Oxford crest. The main content area is white and contains the following text: 'Docs » OSQP solver documentation' with a link to 'Edit on GitHub'. Below this is the title 'OSQP solver documentation' and a call to action: 'Join our forum for any questions related to the solver!'.

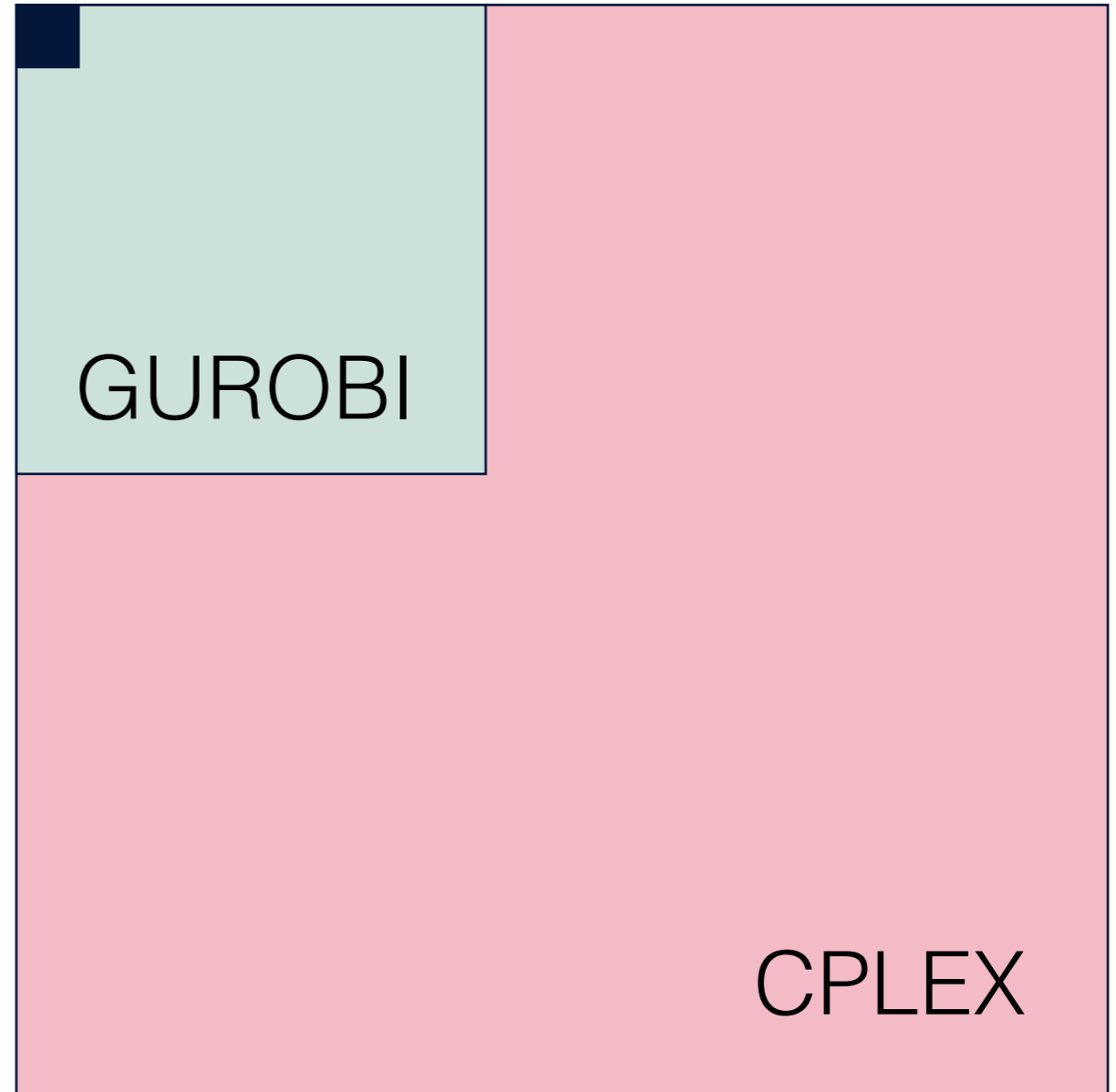
Library
free

Detects
Infeasibility

Embeddable

Compiled code size ~80kb

OSQP



300x
Reduction!

Interfaces

Languages



Parsers

JuMP

CVXPY

YALMIP

OSQP interface



```
# Create OSQP object
m = osqp.OSQP()

# Initialize solver
m.setup(P, q, A, l, u,
        settings)

# Solve
results = m.solve()

# Update cost with q_new
m.update(q=q_new)

# Solve again
results_new = m.solve()
```

```
% Create OSQP object
m = osqp();

% Initialize solver
m.setup(P, q, A, l, u,
        settings);

% Solve
results = m.solve();

% Update cost with q_new
m.update('q', q_new);

% Solve again
results_new = m.solve();
```

Code generation

Optimized
C code

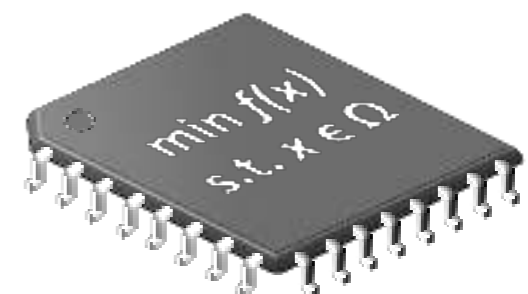
```
# Create OSQP object  
m = osqp.OSQP()  
  
# Initialize solver  
m.setup(P, q, A, l, u,  
        settings)  
  
# Generate C code  
m.codegen('folder_name')
```



```
/* Main ADMM algorithm  
for (iter = 1; iter <= work->settings->max_iter; iter++) {  
  // Main ADMM algorithm  
  // Update x, z, y (preallocated, no malloc)  
  // Compute x^(k+1), z^(k+1), y^(k+1)  
  // Compute x-tilde^(k+1), z-tilde^(k+1)  
  // Compute x^(k+1), z^(k+1), y^(k+1)  
  // End of ADMM Steps */  
#ifdef CTRLC  
  // Check the interrupt signal  
  if (isInterrupted()) {  
    update_status(work->info, OSQP_SIGINT);  
    c_print("Solver interrupted");  
    endInterruptListener();  
    return 1; // exitflag  
  }  
#endif  
}
```



Embedded
hardware



Numerical Example

Lasso

$$\text{minimize} \quad \|Ax - b\|_2^2 + \lambda \|x\|_1$$

Features
 n

Data points
 $m = 100n$

Lasso

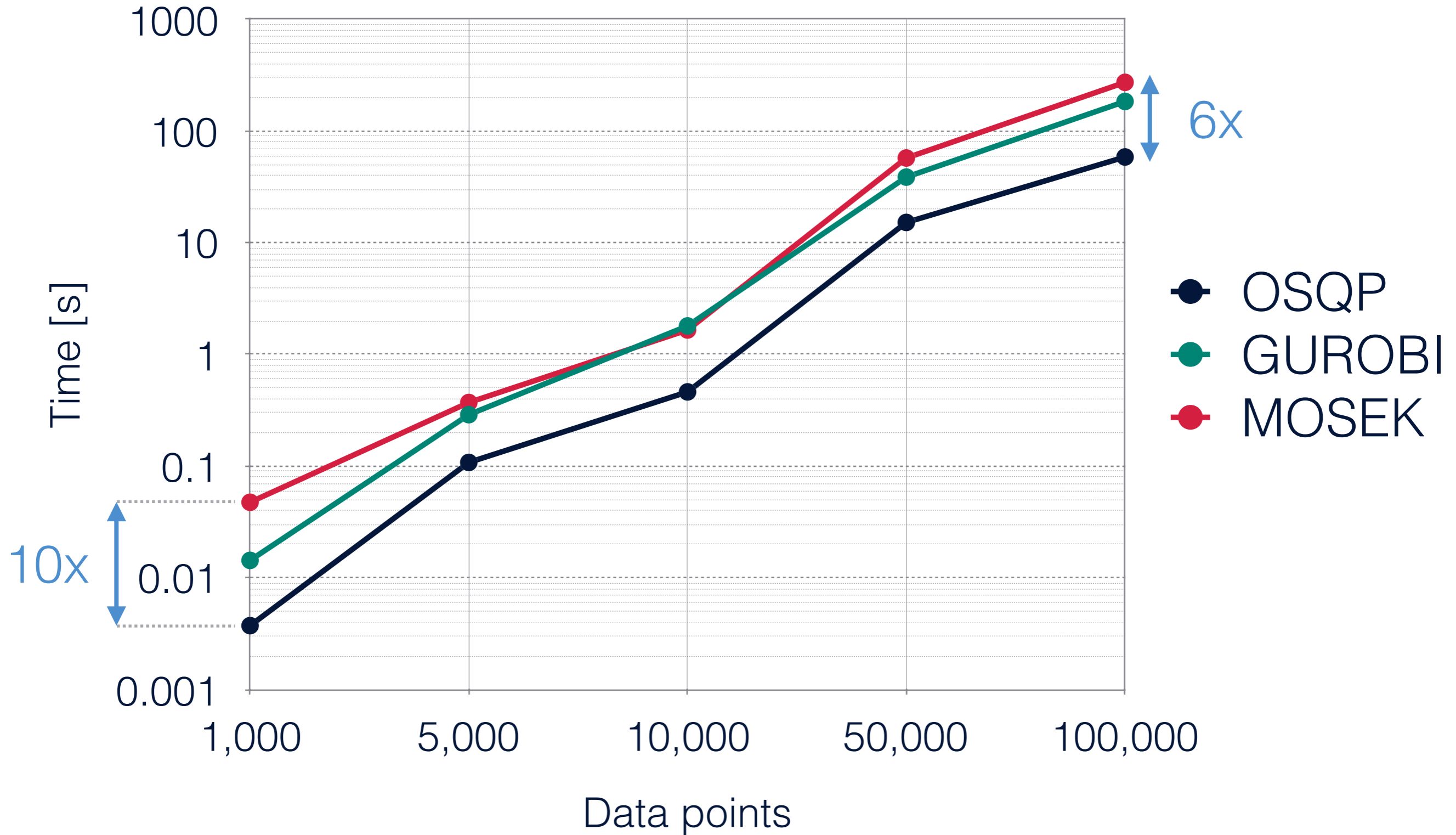
Weighting
parameter

minimize $\|Ax - b\|_2^2 + \lambda \|x\|_1$

Features
 n

Data points
 $m = 100n$

Lasso timings



OSQP
in
“meta-algorithms”

MIQP

$$\begin{array}{ll} \text{minimize} & \frac{1}{2}x^T P x + q^T x \\ \text{subject to} & \underline{l} \leq A x \leq \underline{u} \\ & x_i \in \mathbf{Z} \quad \forall i \in \mathbf{I} \end{array}$$

MIQP

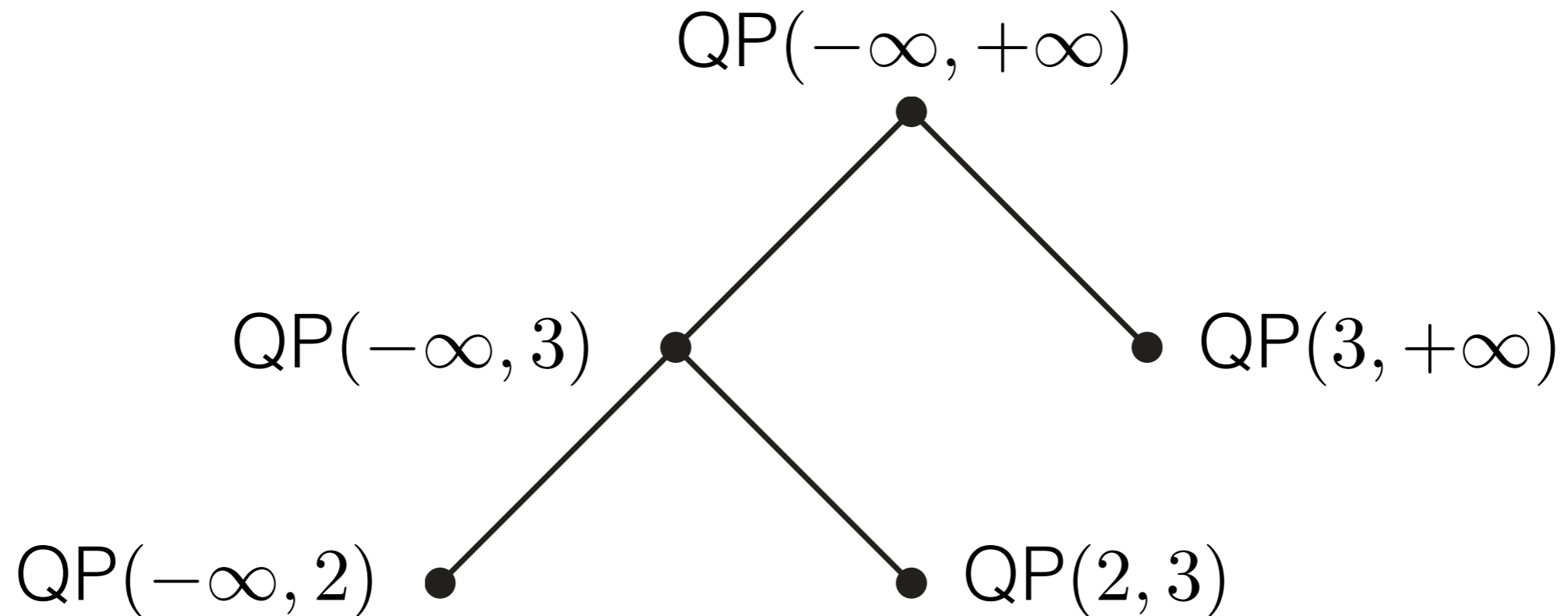
$$\begin{array}{ll} \text{minimize} & \frac{1}{2}x^T P x + q^T x \\ \text{subject to} & l \leq Ax \leq u \\ & x_i \in \mathbf{Z} \quad \forall i \in \mathbf{I} \end{array}$$

Integer
constraints

MIQP

$$\begin{aligned} &\text{minimize} && \frac{1}{2}x^T P x + q^T x \\ &\text{subject to} && l \leq Ax \leq u \\ &&& x_i \in \mathbf{Z} \quad \forall i \in \mathbf{I} \end{aligned}$$

Integer
constraints



Inner QPs

QP(\underline{x}, \bar{x})

$$\begin{array}{ll} \text{minimize} & \frac{1}{2}x^T P x + q^T x \\ \text{subject to} & \underline{l} \leq A x \leq \underline{u} \\ & \underline{x}_i \leq x_i \leq \bar{x}_i \quad \forall i \in \mathbf{I} \end{array}$$

Inner QPs

QP(\underline{x}, \bar{x})

minimize $\frac{1}{2}x^T P x + q^T x$
subject to $\underline{l} \leq A x \leq \underline{u}$

$$\underline{x}_i \leq x_i \leq \bar{x}_i \quad \forall i \in \mathbf{I}$$

Changing
bounds



Reusing
factorization

Saving computations

Factorization
caching

+

Warm
starting

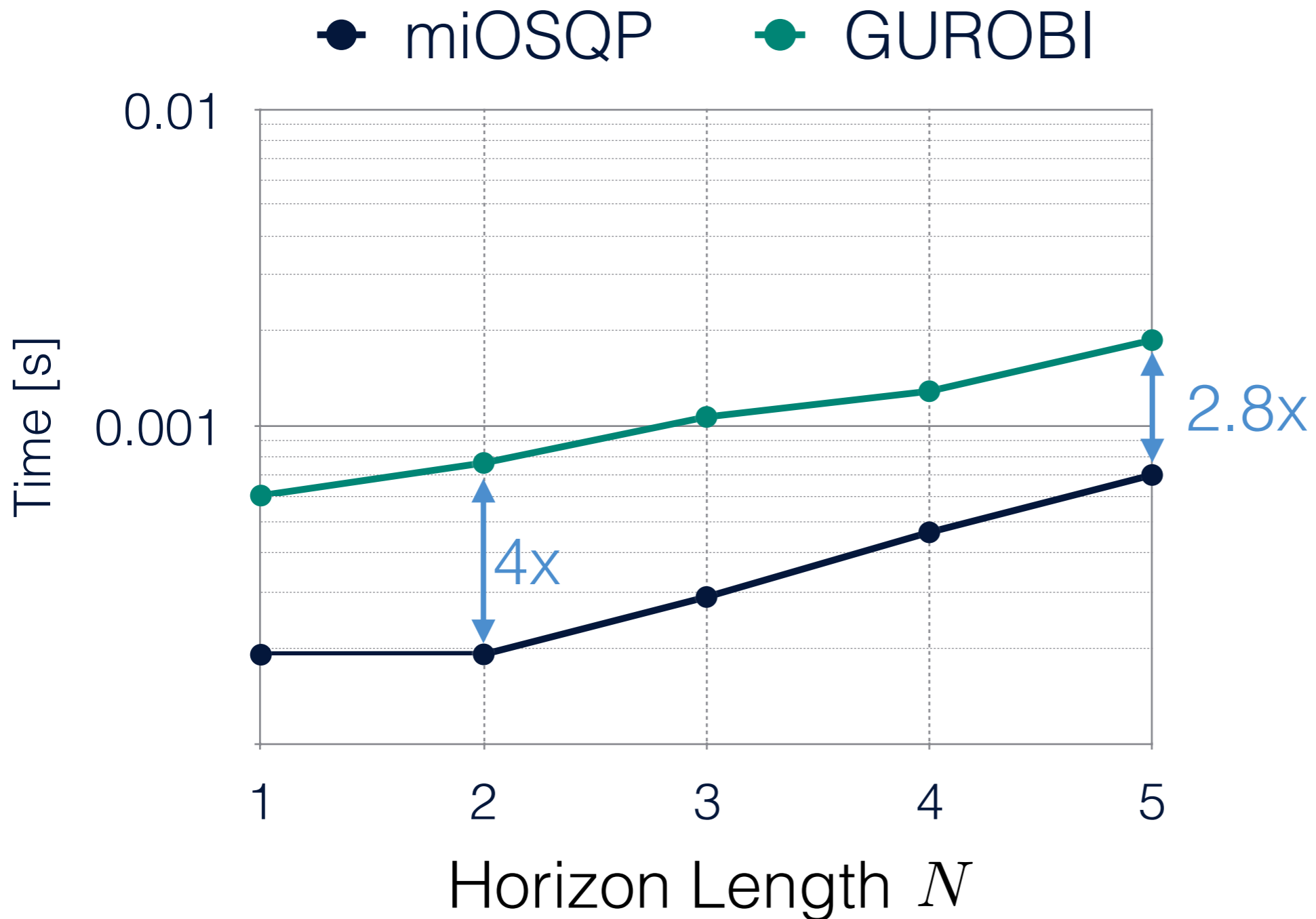
ADMM

QP(\underline{x}, \bar{x})

Repeated
MIQPs

MIQP Timings

Hybrid Model Predictive Control



Simple
Python
Implementation!

Conclusions

Acknowledgements



Goran
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Stanford



Paul
Goulart

Oxford



Alberto
Bemporad

IMT Lucca



Stephen
Boyd

Stanford

Final remarks

OSQP

Simple

Robust

Embeddable

Julia interface

Exploit
Initialization

C code
generation

New
high-level
algorithms

References

B. Stellato, G. Banjac, P. Goulart, A. Bemporad and S. Boyd. *OSQP: An Operator Splitting Solver for Quadratic Programs. (Coming soon!)*

G. Banjac, P. Goulart, B. Stellato, and S. Boyd. *Infeasibility detection in the alternating direction method of multipliers for convex optimization.* optimization-online.org, 2017

G. Banjac, B. Stellato, N. Moehle, P. Goulart, A. Bemporad and S. Boyd. *Embedded code generation using the OSQP solver.* IEEE Conference on Decision and Control (CDC) (submitted), 2017

B. Stellato, V. Naik, A. Bemporad, P. Goulart, and S. Boyd. *Embedded mixed-integer quadratic optimization using the OSQP solver.* IEEE Conference on Decision and Control (CDC) (submitted), 2017