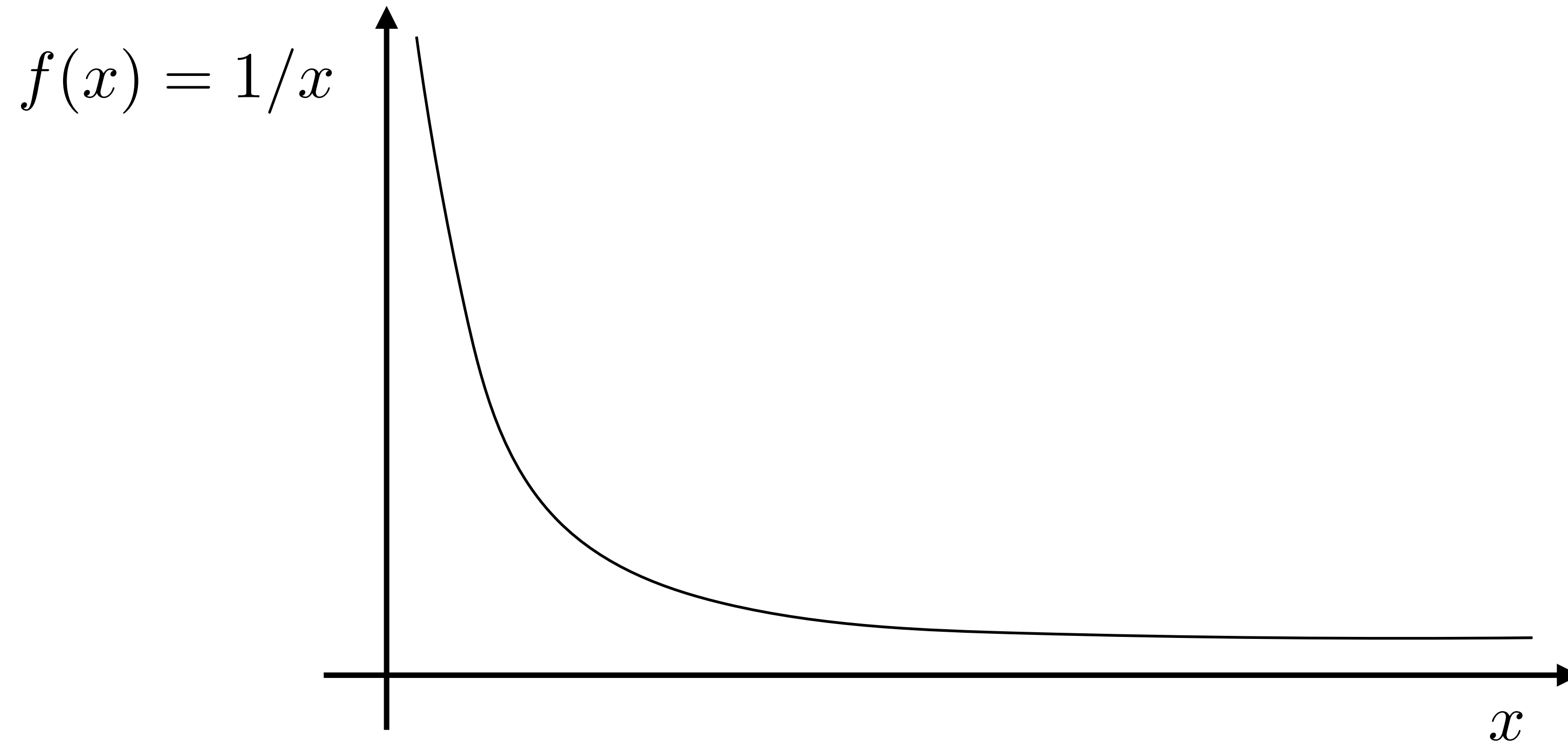


ORF522 – Linear and Nonlinear Optimization

15. Subgradient methods

Recap

Many functions are not strongly convex



Without strong convexity, the optimal solution might be very far ($x^* = \infty$) but the objective value very close

Focus on objective error $f(x^k) - f(x^*)$ instead of variable error $\|x^k - x^*\|_2$

Convergence for smooth functions

Theorem

Let f be L -smooth. If $t < 1/L$ then gradient descent satisfies

$$f(x^k) - f(x^*) \leq \frac{\|x^0 - x^*\|_2^2}{2tk}$$

Sublinear convergence rate $O(1/\epsilon)$ iterations (can be very slow!)

Issues with computing the optimal step size

Quadratic programs

The rule $t = 2/(\lambda_{\max}(P) + \lambda_{\min}(P))$ can be **very expensive to compute**

It relies on eigendecomposition of P (iterative factorizations...)

Smooth and strongly convex functions

Very hard to estimate μ and L in general

Can we select a good step-size as we go?

Line search

Exact line search

Choose the best step along the descent direction

$$t_k = \operatorname{argmin}_{t \geq 0} f(x^k - t \nabla f(x^k))$$

Used when

- computational cost very low or
- there exist closed-form solutions

In general, impractical to perform exactly

Backtracking line search

Condition

Armijo condition: for some $0 < \alpha \leq 1/2$

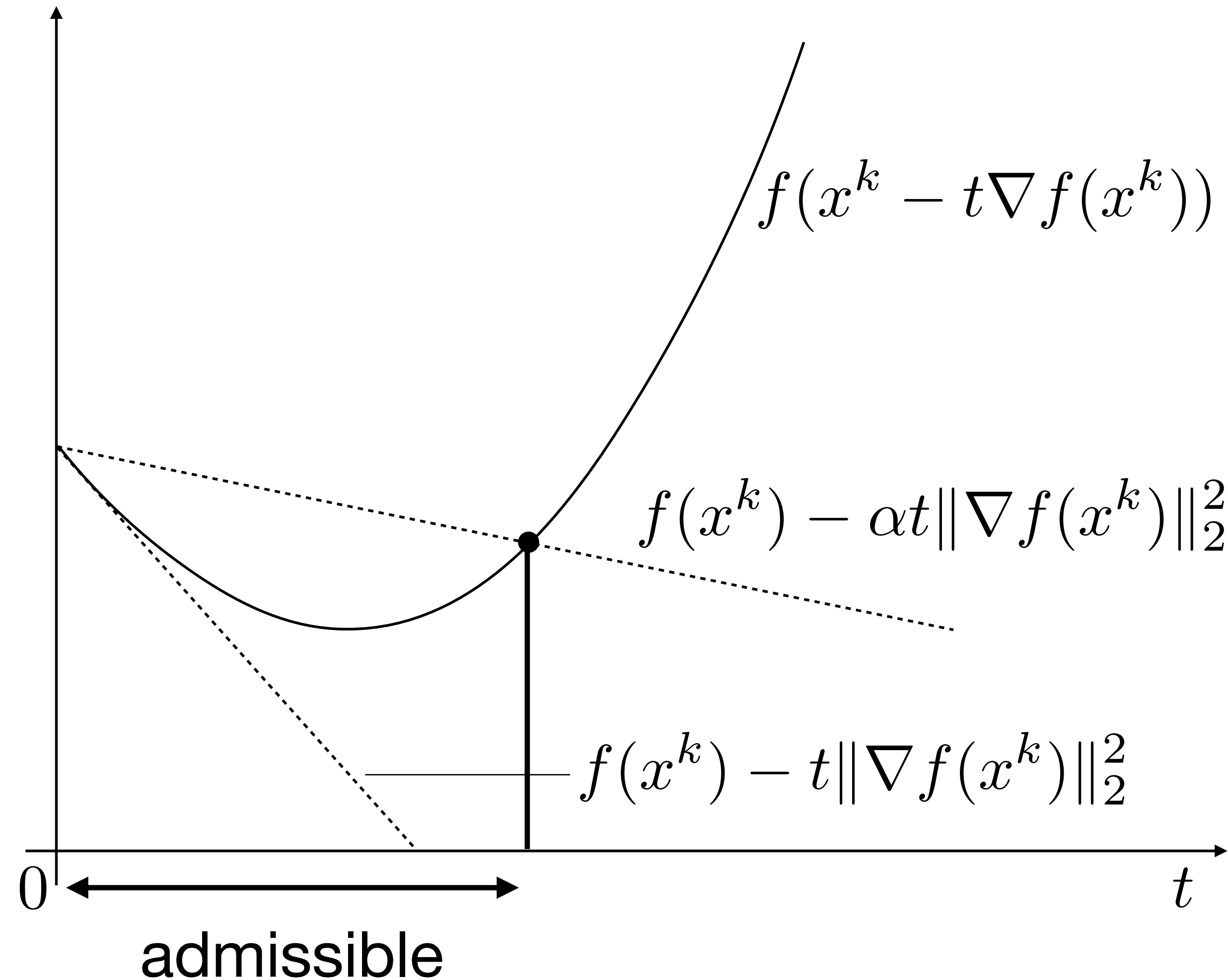
$$f(x^k + td^k) < f(x^k) + \alpha t \nabla f(x^k)^T d^k$$



where $d^k = -\nabla f(x^k)$

$$f(x^k - t\nabla f(x^k)) < f(x^k) - \alpha t \|\nabla f(x^k)\|_2^2$$

Guarantees
sufficient decrease
in objective value



Backtracking line search

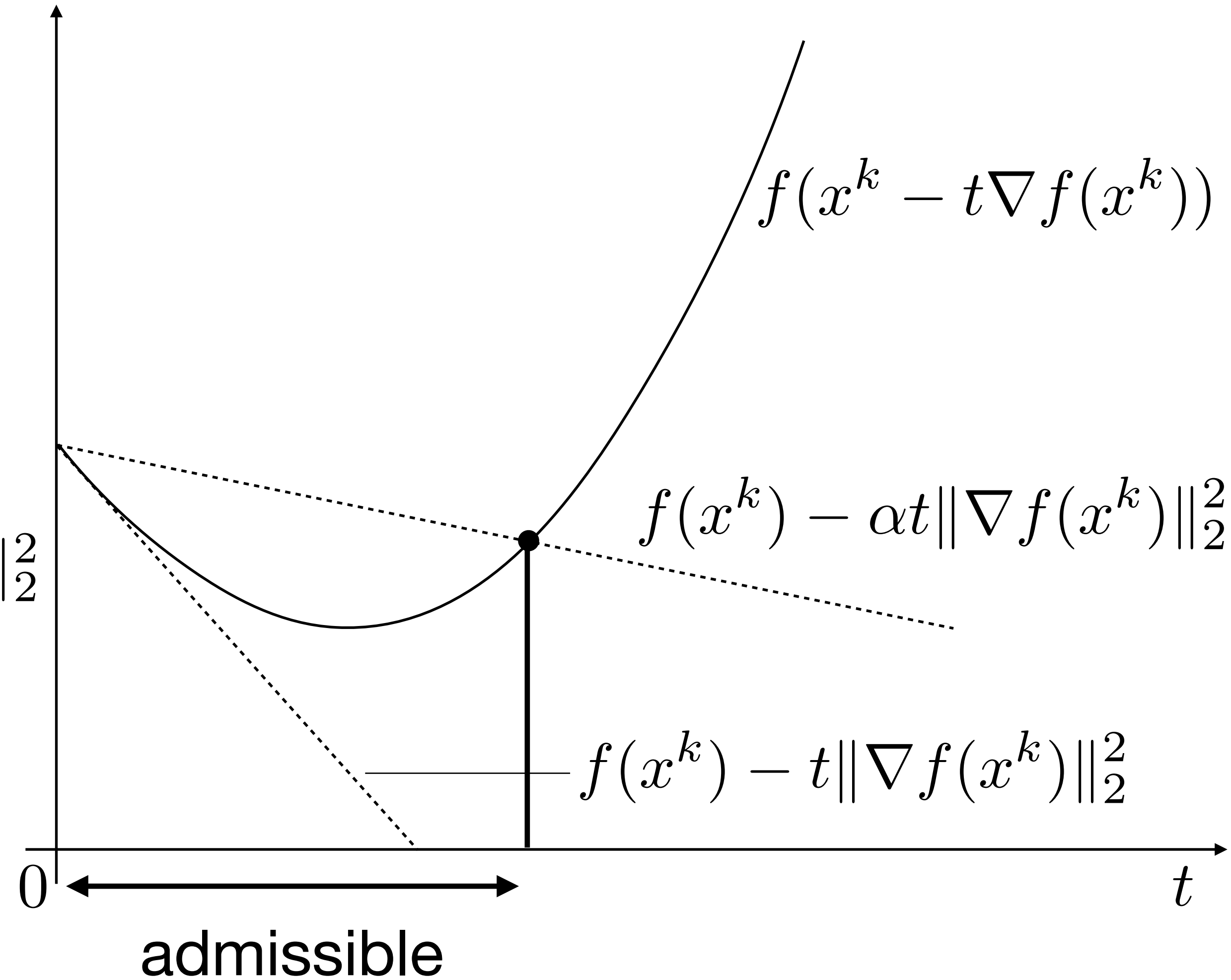
Iterations

initialization

$$t = 1, \quad 0 < \alpha \leq 1/2, \quad 0 < \beta < 1$$

$$\mathbf{while} \quad f(x^k - t\nabla f(x^k)) > f(x^k) - \alpha t \|\nabla f(x^k)\|_2^2$$

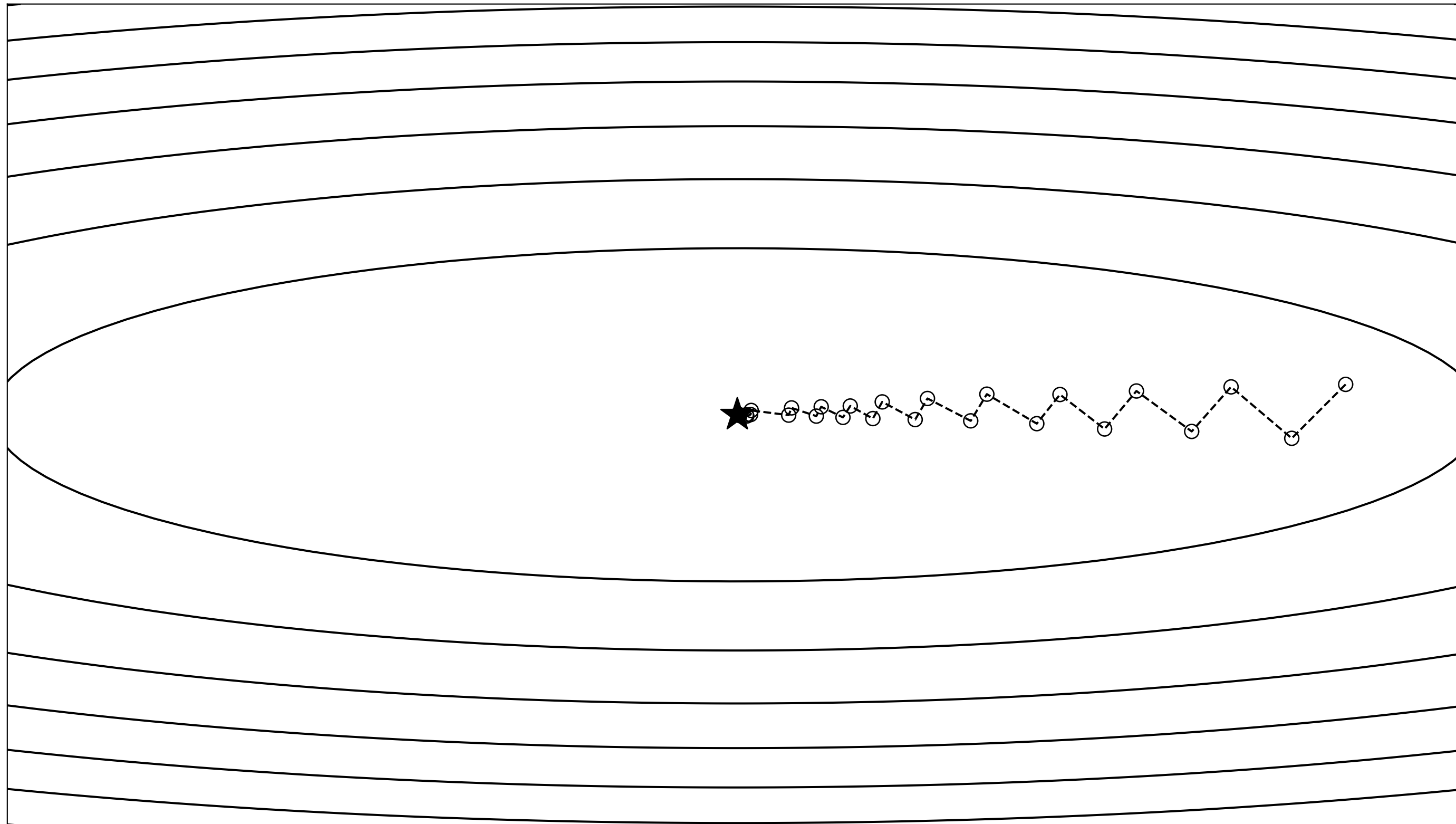
$$t \leftarrow \beta t$$



Backtracking line search

$$f(x) = (x_1^2 + 20x_2^2)/2$$

$$x^0 = (20, 1)$$



Backtracking line search

Converges in 31 iterations

Backtracking line search convergence

Theorem

Let f be L -smooth. Gradient descent with backtracking line search satisfies

$$f(x^k) - f(x^*) \leq \frac{\|x^0 - x^*\|_2^2}{2t_{\min}k}$$

where $t_{\min} = \min\{1, \beta/L\}$

Proof almost identical to fixed step case

Remarks

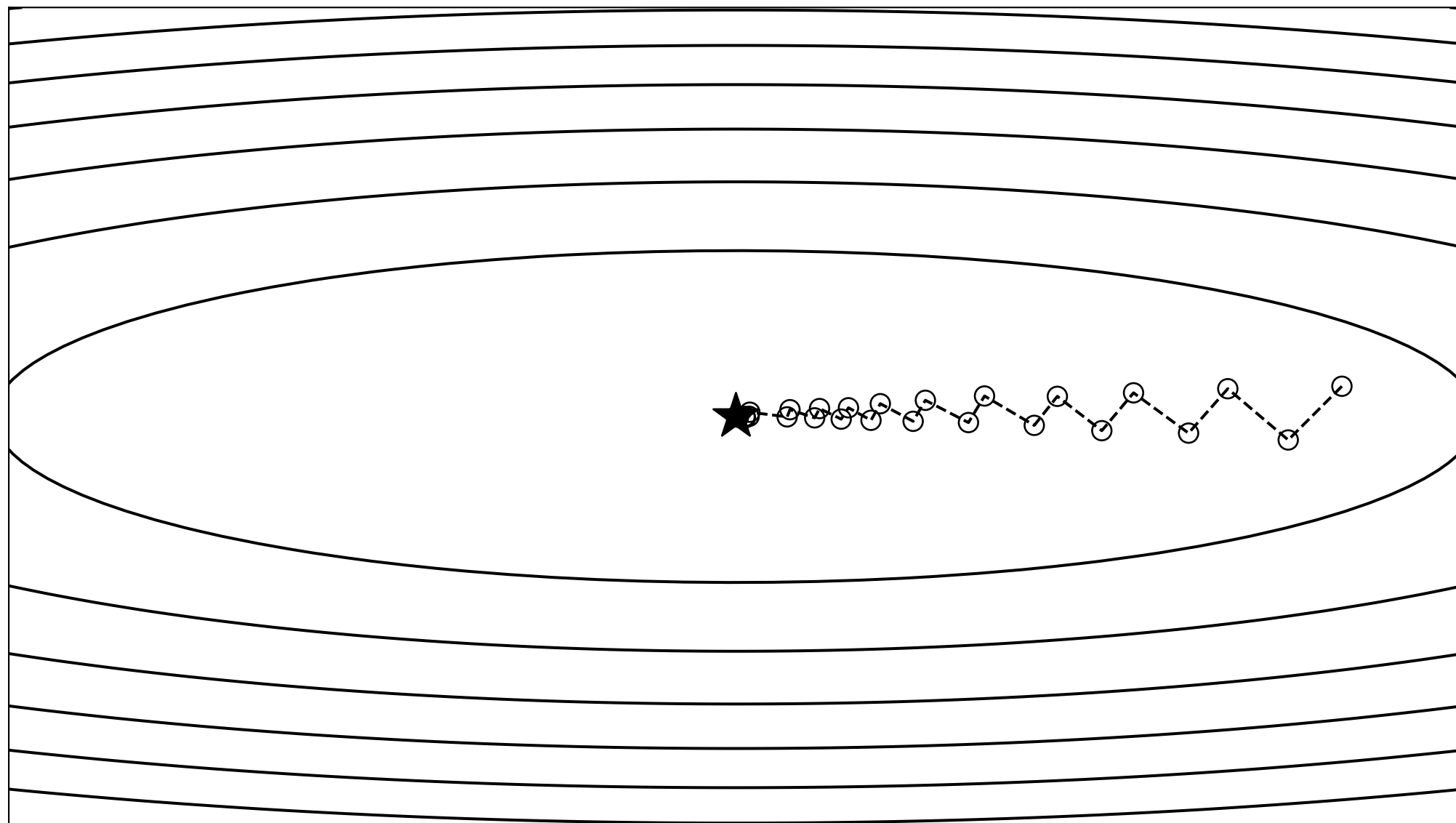
- If $\beta \approx 1$, similar to optimal step-size (β/L vs $1/L$)
- Still convergence rate $O(1/\epsilon)$ iterations (can be very slow!)

Gradient descent issues

Slow convergence

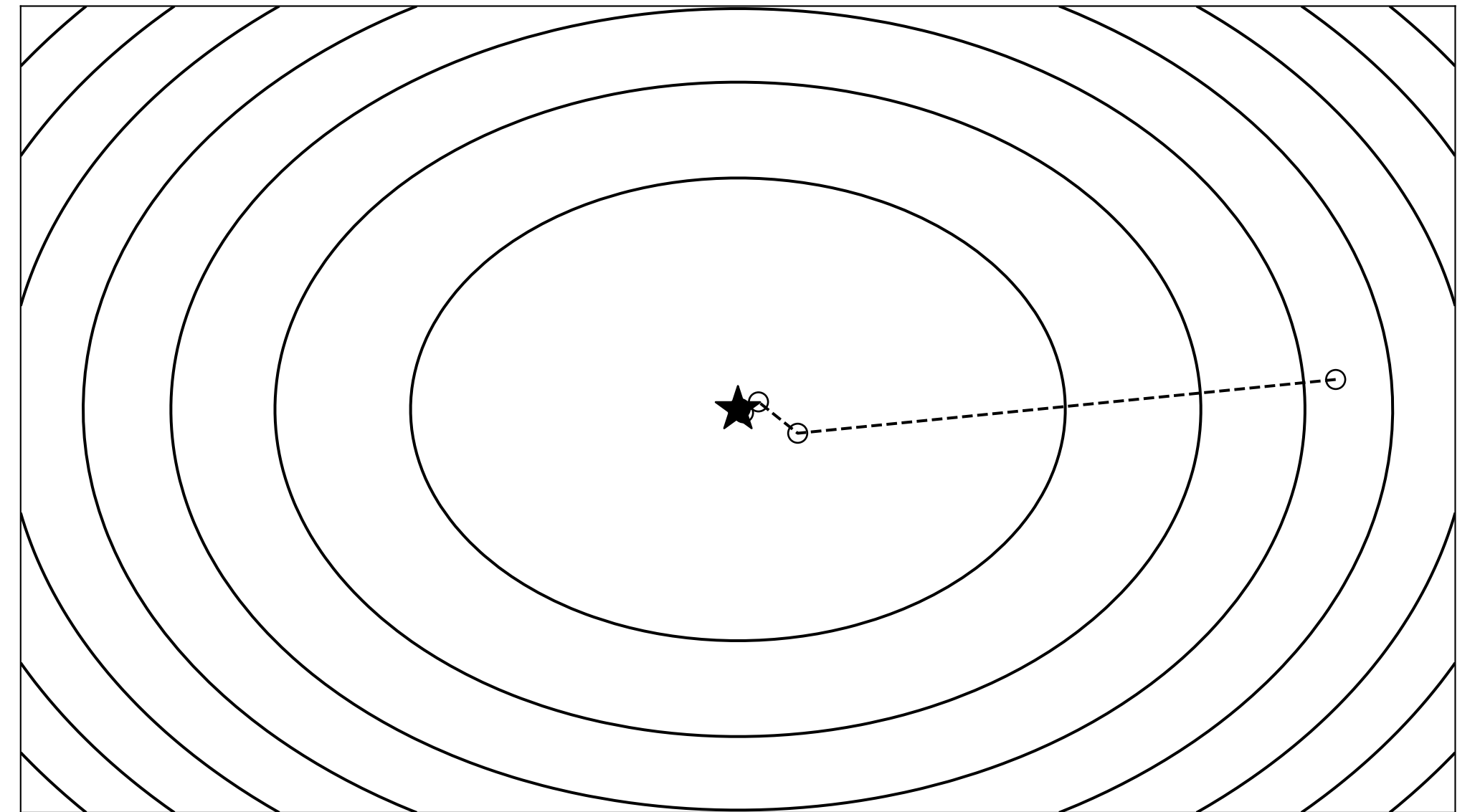
Very dependent on scaling

$$f(x) = (x_1^2 + 20x_2^2)/2$$



Slow convergence

$$f(x) = (x_1^2 + 2x_2^2)/2$$

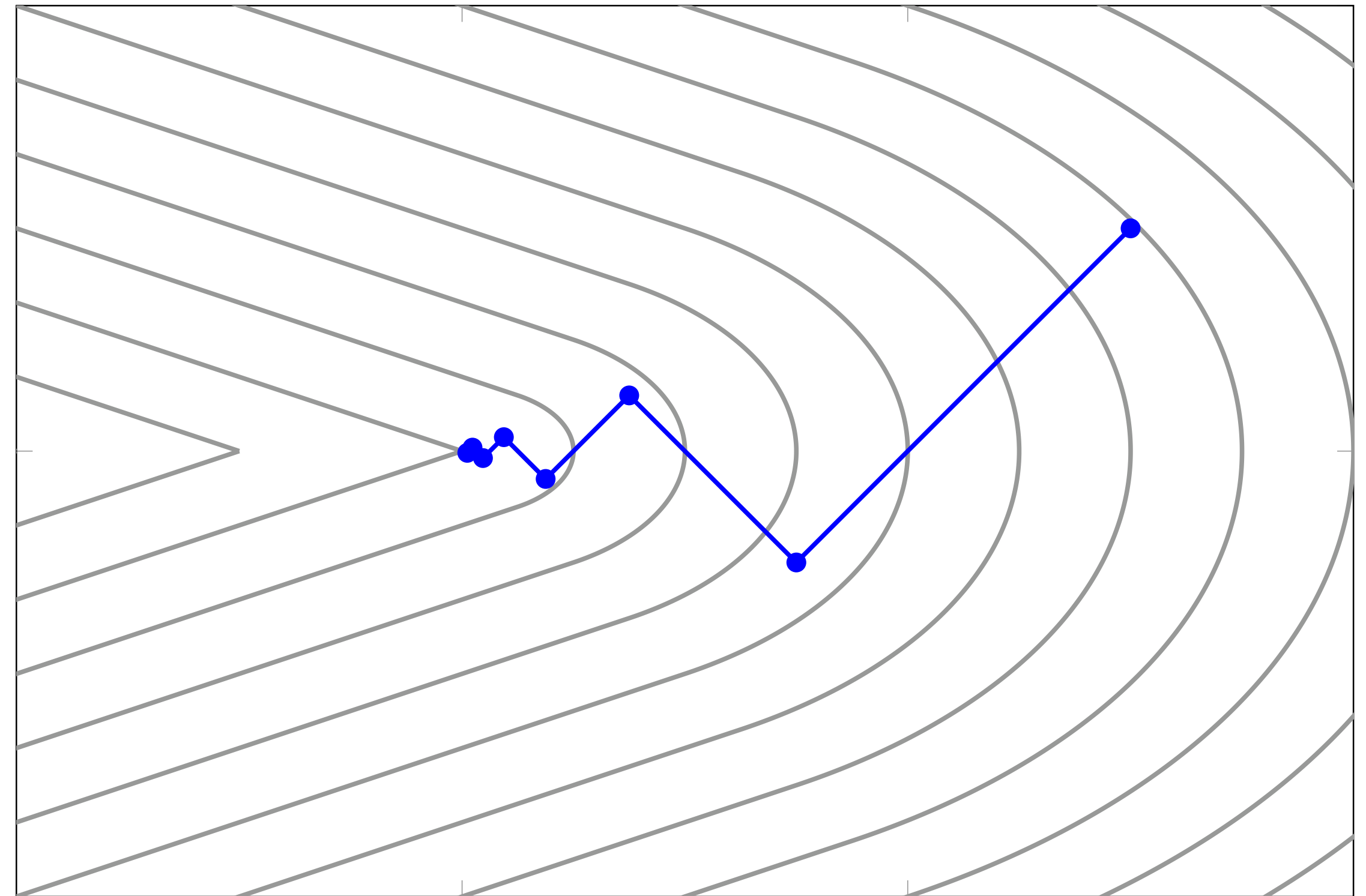


Faster

Non-differentiability

Wolfe's example

$$f(x) = \begin{cases} \sqrt{x_1^2 + \gamma x_2^2} & |x_2| \leq x_1 \\ \frac{x_1 + \gamma|x_2|}{\sqrt{1 + \gamma}} & |x_2| > x_1 \end{cases}$$



Gradient descent with *exact line search* gets stuck at $x = (0, 0)$

In general: gradient descent cannot handle non-differentiable functions and constraints

Today's lecture

[Chapter 3 and 8, FMO][ee364b][Chapter 3, ILCO]

Subgradient methods

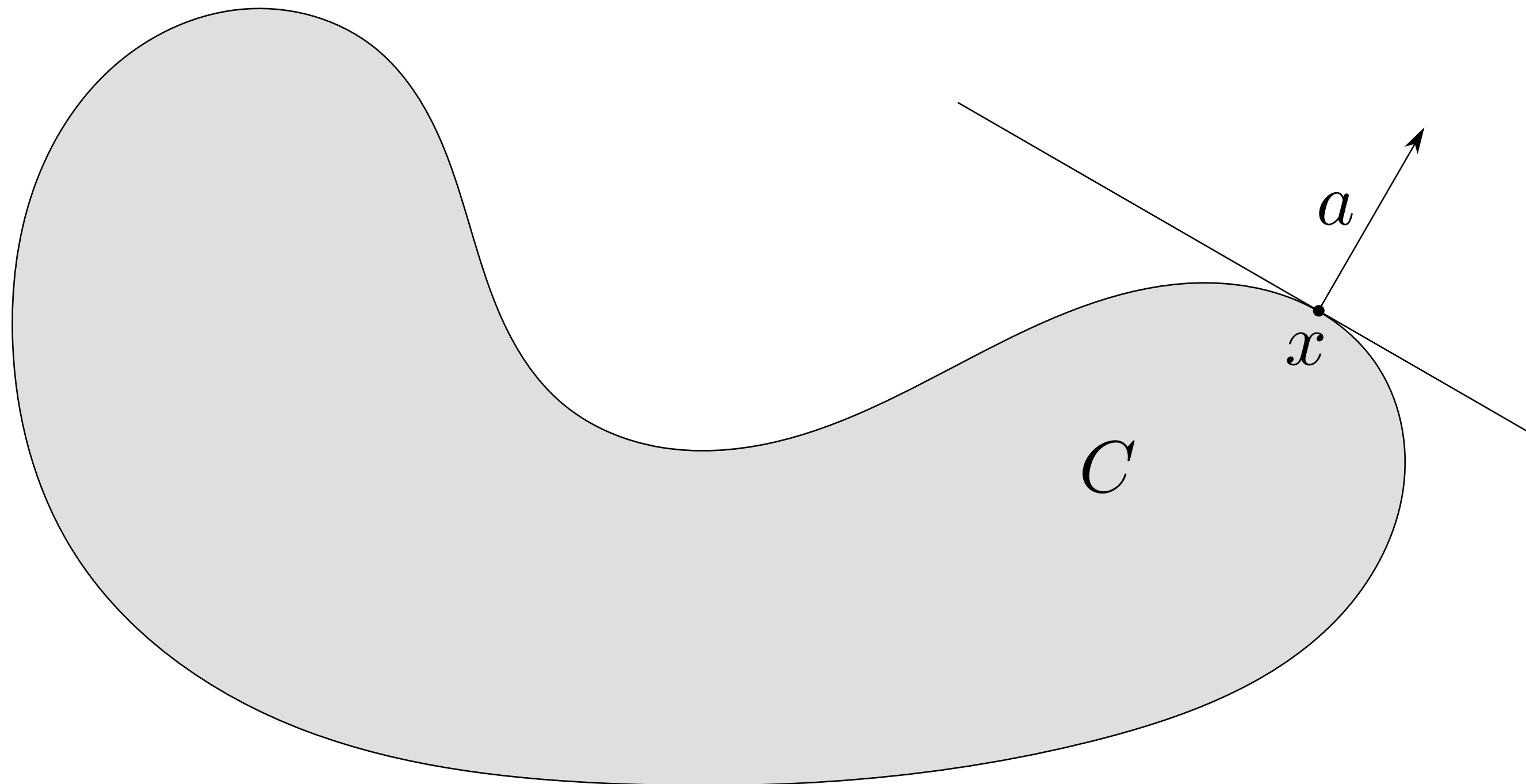
- Geometric definitions
- Subgradients
- Subgradient calculus
- Optimality conditions based on subgradients
- Subgradient methods

Geometric definitions

Supporting hyperplanes

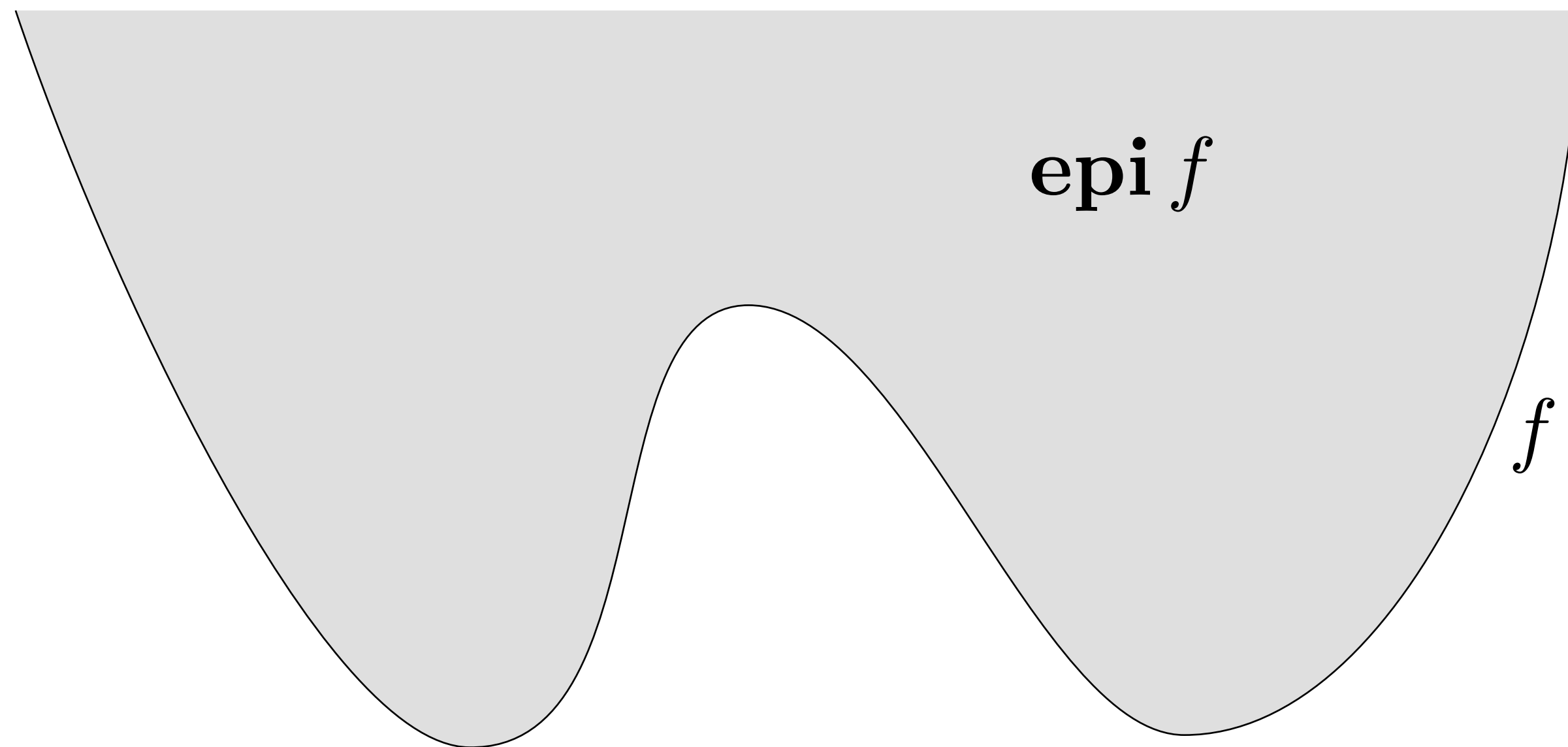
Given a set C point x at the boundary of C
a hyperplane $\{z \mid a^T z = a^T x\}$ is a **supporting hyperplane** if

$$a^T (y - x) \leq 0, \quad \forall y \in C$$



Function epigraph

$$\text{epi } f = \{(x, t) \mid x \in \text{dom } f, f(x) \leq t\}$$



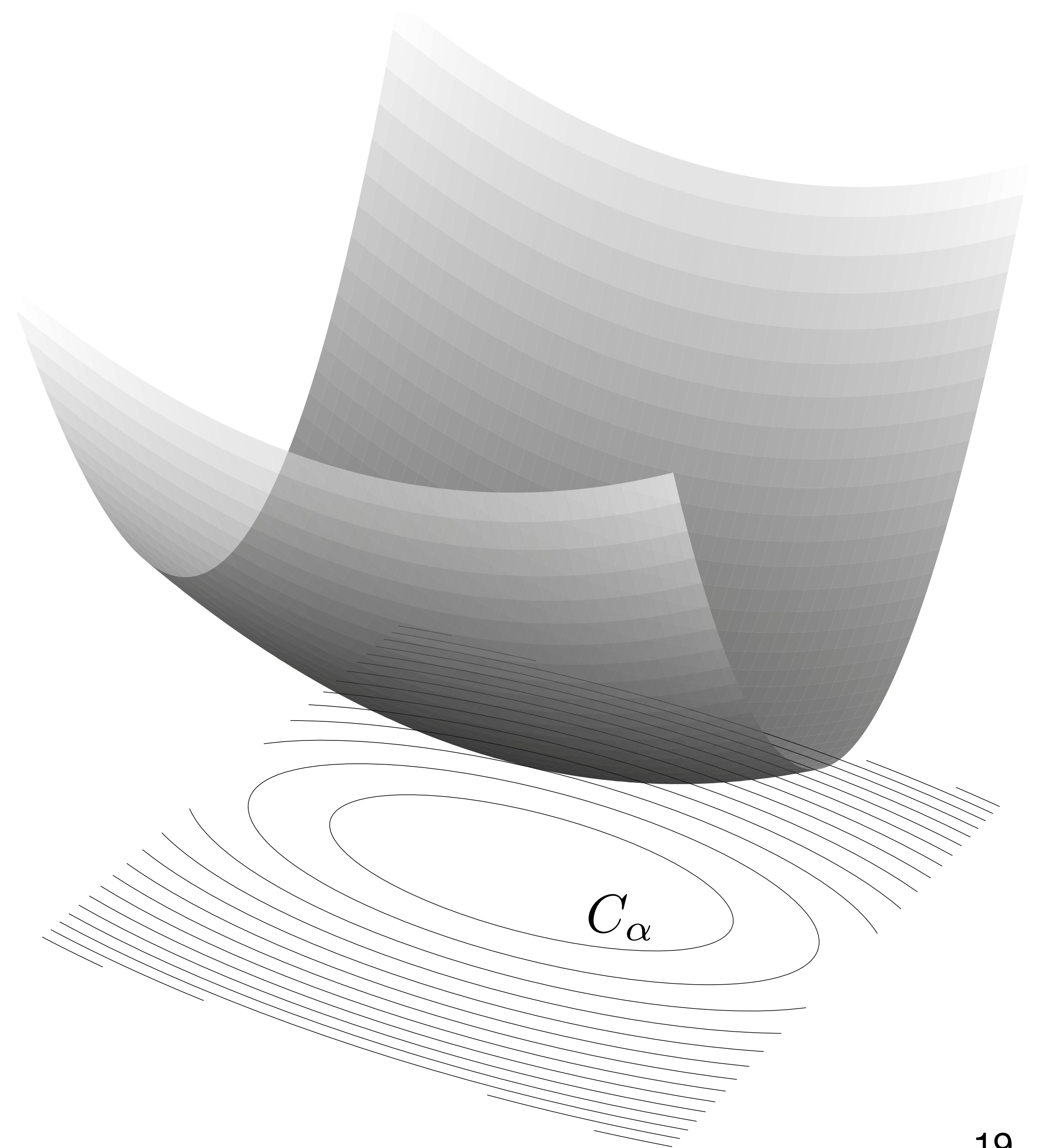
f is convex if and only if $\text{epi } f$ is a convex set

Sublevel sets

$$C_\alpha = \{x \in \text{dom } f \mid f(x) \leq \alpha\}$$

If f is convex, then C_α is convex $\forall \alpha$

Note converse not true, e.g., $f(x) = -e^x$



Subgradients

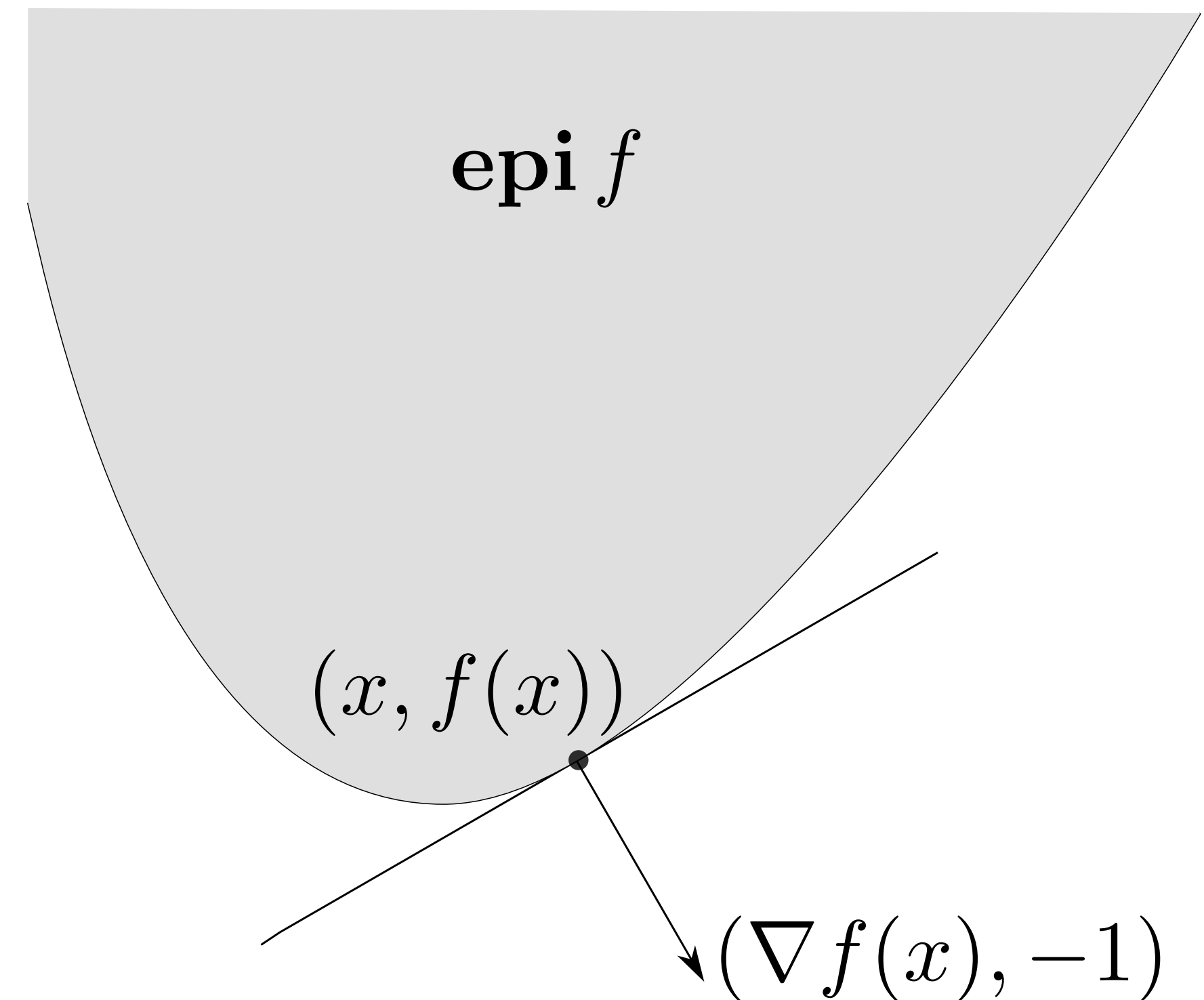
Gradients and epigraphs

For a convex differentiable function f , i.e.

$$f(y) \geq f(x) + \nabla f(x)^T (y - x), \quad \forall y \in \mathbf{dom} f$$

$(\nabla f(x), -1)$ defines a **supporting hyperplane** to epigraph of f at $(x, f(x))$

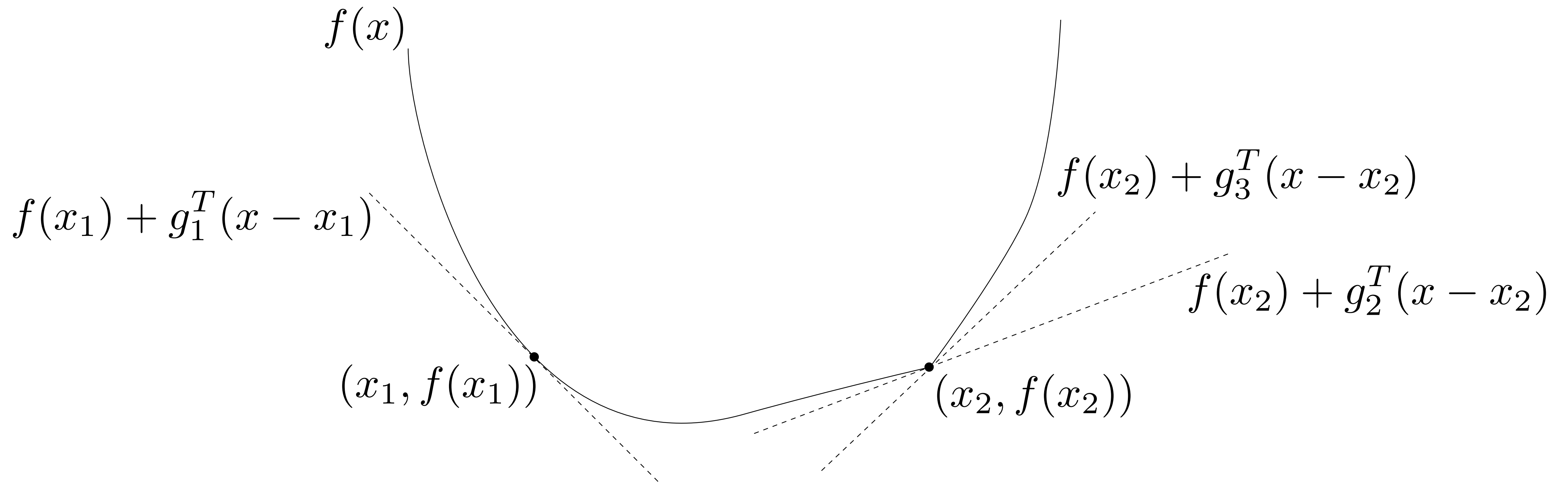
$$\begin{bmatrix} \nabla f(x) \\ -1 \end{bmatrix}^T \left(\begin{bmatrix} y \\ t \end{bmatrix} - \begin{bmatrix} x \\ f(x) \end{bmatrix} \right) \leq 0, \quad \forall (y, t) \in \mathbf{epi} f$$



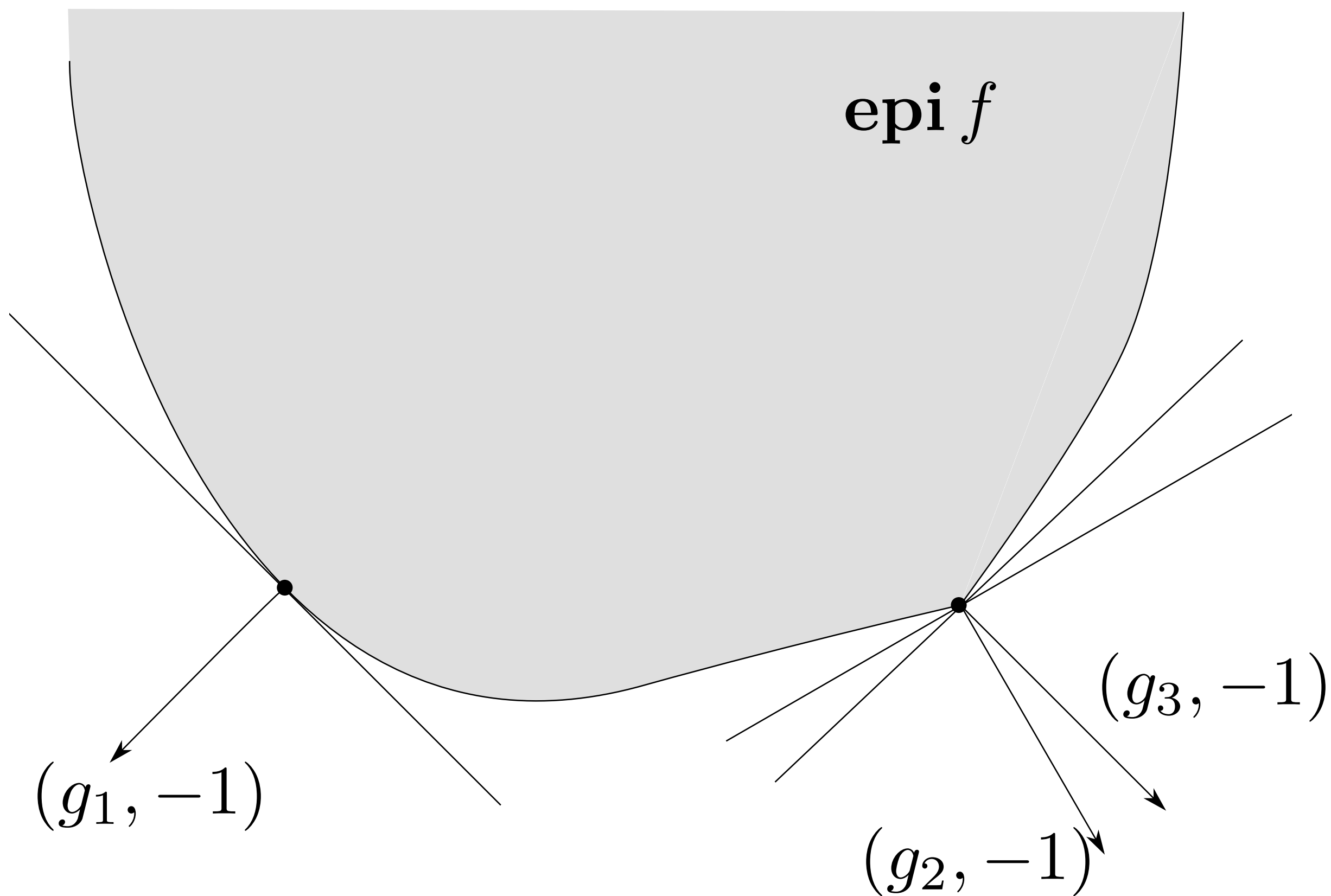
Subgradient

We say that g is a **subgradient** of function f at point x if

$$f(y) \geq f(x) + g^T(y - x), \quad \forall y$$



Subgradient properties



g is a subgradient of f at x iff $(g, -1)$ supports epi f at $(x, f(x))$

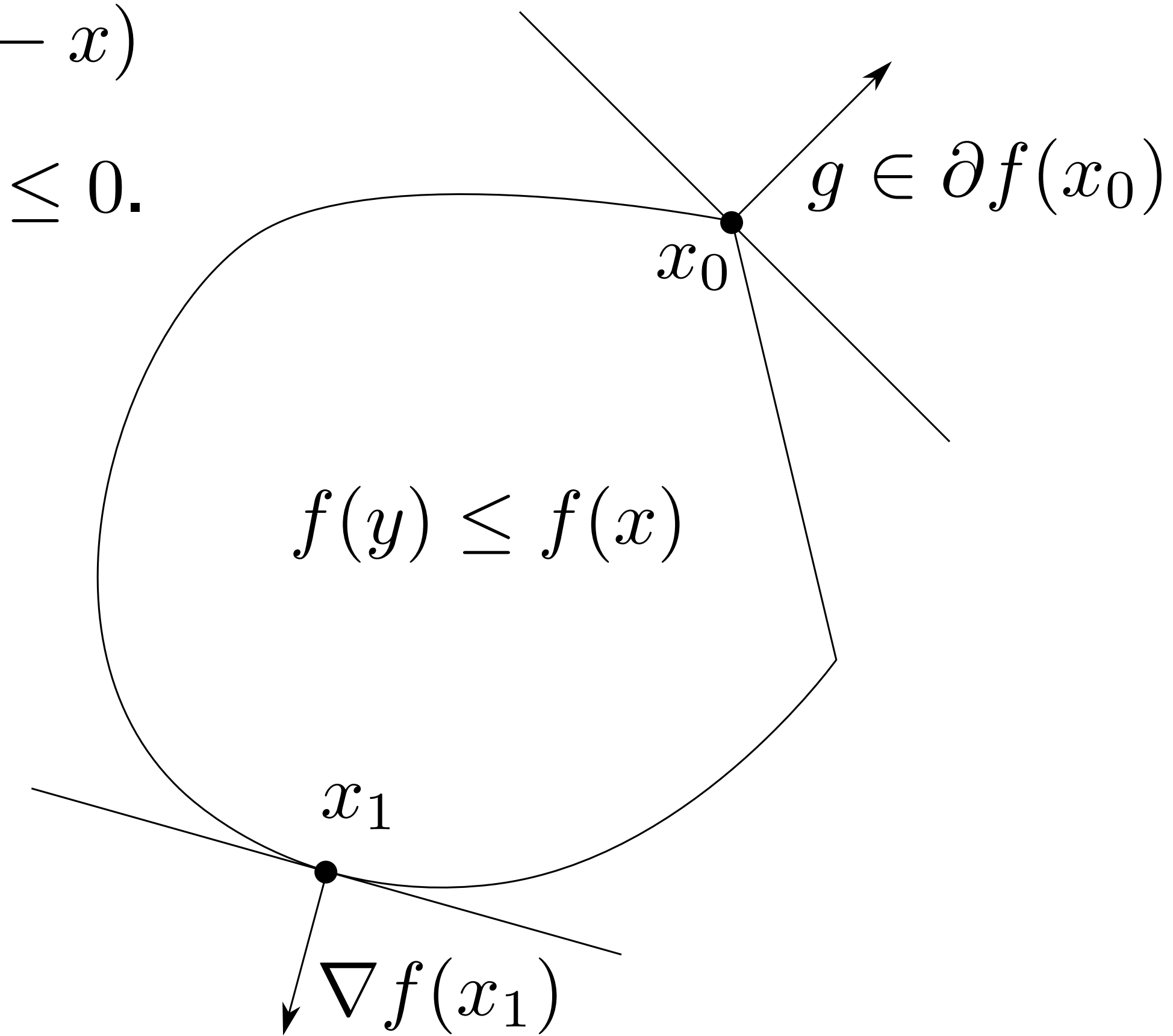
g is a subgradient of f iff $f(x) + g^T(y - x)$ is a global underestimator of f

If f is convex and differentiable, $\nabla f(x)$ is a subgradient of f at x

(Sub)gradients and sublevel sets

g being a subgradient of f means $f(y) \geq f(x) + g^T(y - x)$

Therefore, if $f(y) \leq f(x)$ (sublevel set), then $g^T(y - x) \leq 0$.



f differentiable at x

$\nabla f(x)$ is normal to the sublevel set $\{y \mid f(y) \leq f(x)\}$

f nondifferentiable at x

subgradients define supporting hyperplane to sublevel set through x

Subdifferential

The subdifferential $\partial f(x)$ of f at x is the **set of all subgradients**

$$\partial f(x) = \{g \mid g^T(y - x) \leq f(y) - f(x), \quad \forall y \in \text{dom } f\}$$

Properties

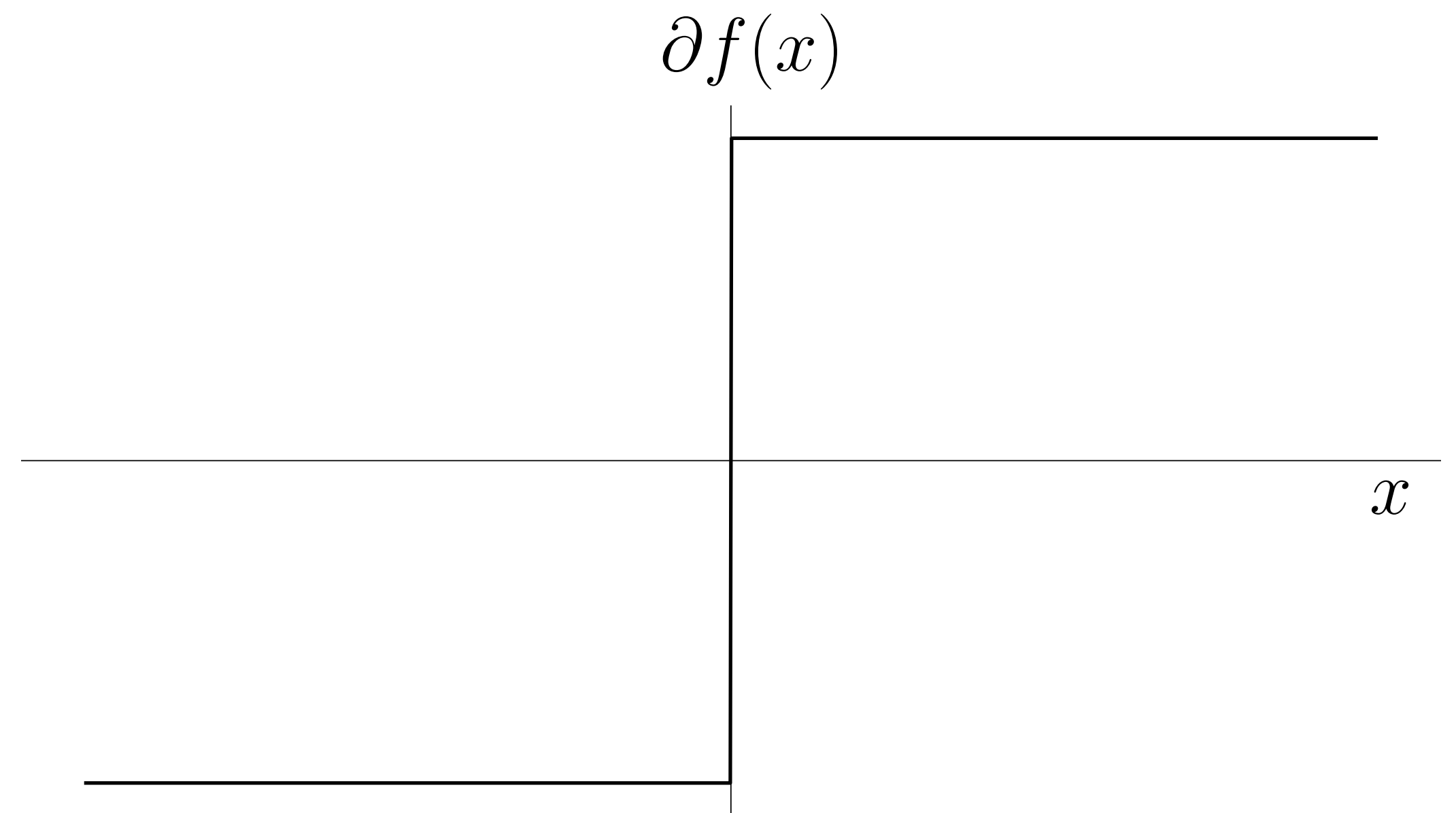
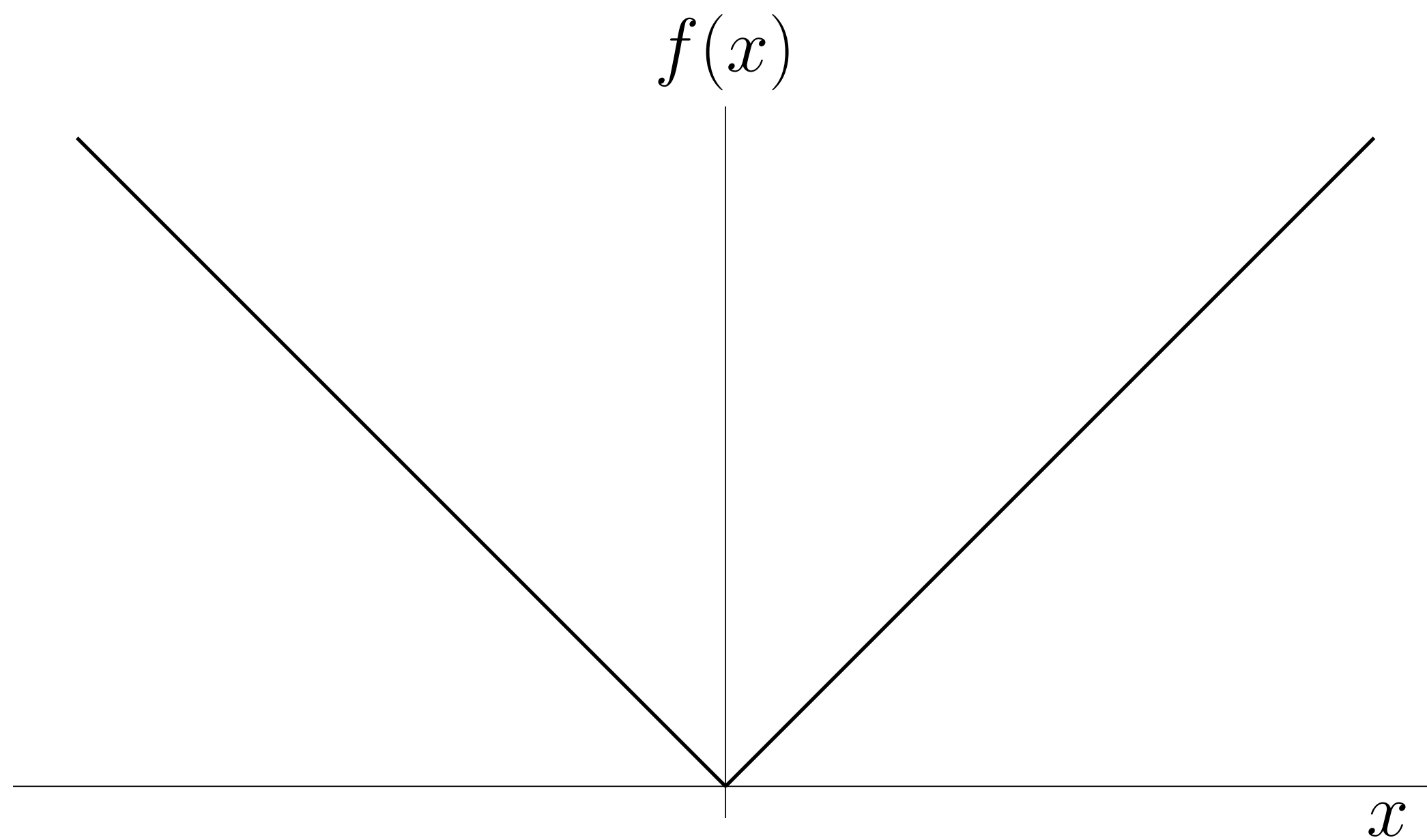
- $\partial f(x)$ is always closed and convex, also for nonconvex f .
(intersection of halfspaces)
- If f is convex and differentiable at x , then $\partial f(x) = \{\nabla f(x)\}$
- If f is convex and $\partial f(x) = \{g\}$, then f is differentiable at x and $g = \nabla f(x)$

Example

Absolute value

$$f(x) = |x|$$

$$\partial f(x) = \begin{cases} \{-1\} & x < 0 \\ [-1, 1] & x = 0 \\ \{1\} & x > 0 \end{cases} = \begin{cases} \mathbf{sign}(x) & x \neq 0 \\ [-1, 1] & x = 0 \end{cases}$$



Subgradient calculus

Subgradient calculus

Strong subgradient calculus

Formulas for finding the whole subdifferential $\partial f(x)$ \longrightarrow **Hard**

Weak subgradient calculus

Formulas for finding *one* subgradient $g \in \partial f(x)$ \longrightarrow **Easy**

In practice, most algorithms require only *one* subgradient g at point x

Basic rules

Nonnegative scaling: $\partial(\alpha f) = \alpha \partial f$ with $\alpha > 0$

Addition: $\partial(f_1 + f_2) = \partial f_1 + \partial f_2$

Affine transformation: $f(x) = h(Ax + b)$, then

$$\partial f(x) = A^T \partial h(Ax + b)$$

Basic rules

Pointwise maxima

Finite pointwise maximum $f(x) = \max_{i=1,\dots,m} f_i(x)$, then

$$\partial f(x) = \text{conv} \left(\bigcup \{ \partial f_i(x) \mid f_i(x) = f(x) \} \right) \quad (\text{convex hull of active functions})$$

General pointwise maximum (supremum) $f(x) = \max_{s \in S} f_s(x)$, then

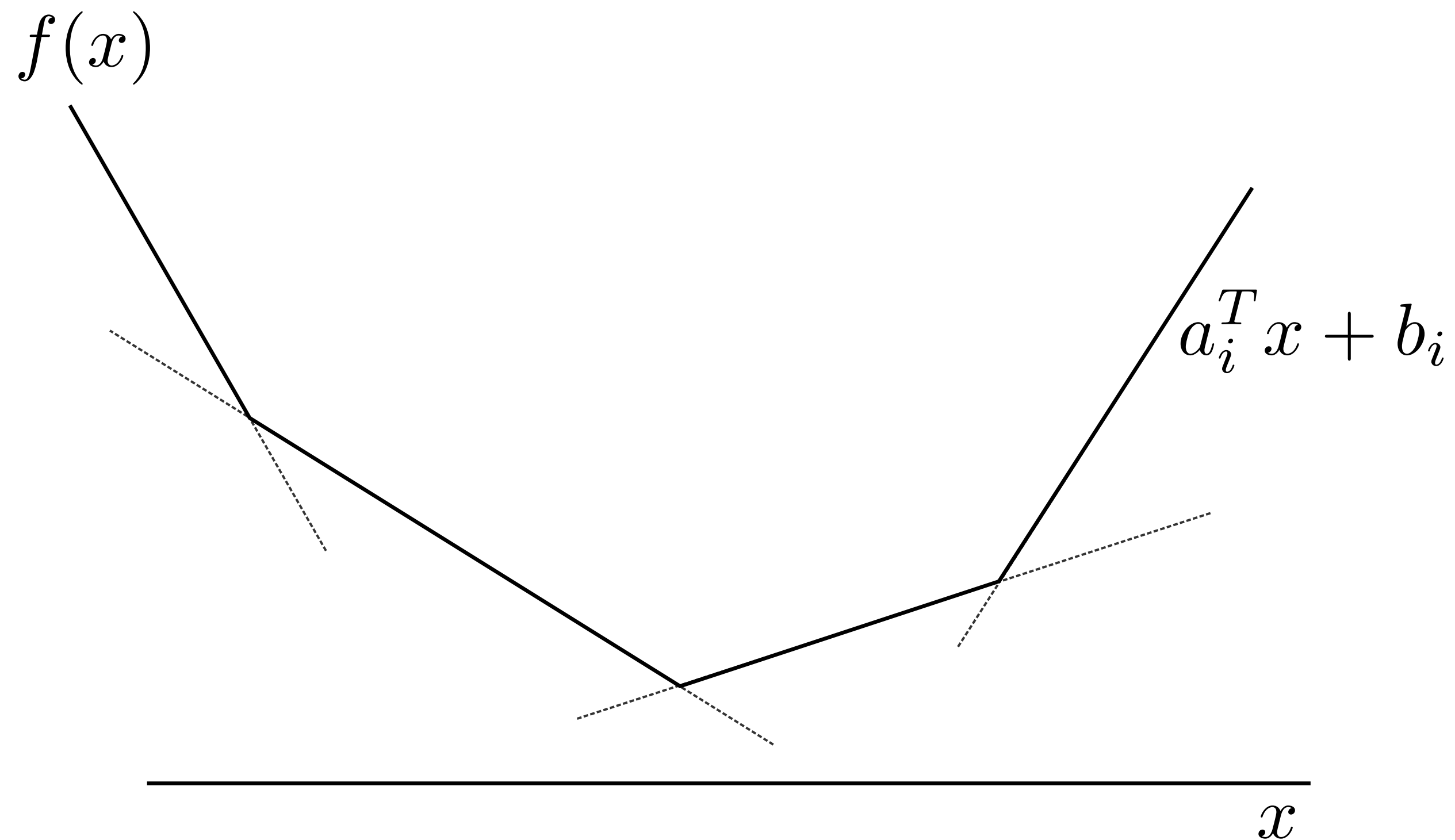
$$\partial f(x) \supseteq \text{conv} \left(\bigcup \{ \partial f_s(x) \mid f_s(x) = f(x) \} \right)$$

Note: Equality requires some regularity assumptions
(e.g. S compact and f_s is continuous in s)

Example

Piecewise linear function

$$f(x) = \max_{i=1, \dots, m} (a_i^T x + b_i)$$



Subdifferential is a polyhedron

$$\partial f(x) = \mathbf{conv}\{a_i \mid i \in I(x)\}$$

$$I(x) = \{i \mid a_i^T x + b_i = f(x)\}$$

Example

Norms

Given $f(x) = \|x\|_p$ we can express it as

$$\|x\|_p = \max_{\|z\|_q \leq 1} z^T x,$$

where q such that $1/p + 1/q = 1$ defines the **dual norm**. Therefore,

$$\partial f(x) = \operatorname{argmax}_{\|z\|_q \leq 1} z^T x$$

Example: $f(x) = \|x\|_1 = \max_{\|s\|_\infty \leq 1} s^T x$

$$\partial f(x) = J_1 \times \cdots \times J_n \quad \text{where} \quad J_i = \begin{cases} \{-1\} & x < 0 \\ [-1, 1] & x = 0 \\ \{1\} & x > 0 \end{cases}$$

weak result
 $\operatorname{sign}(x) \in \partial f(x)$

Basic rules

Composition

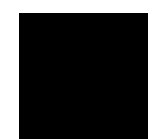
$f(x) = h(f_1(x), \dots, f_k(x))$, h convex nondecreasing, f_i convex

$$g = q_1 g_1 + \dots + q_k g_k \in \partial f(x)$$

where $q \in \partial h(f_1(x), \dots, f_k(x))$ and $g_i \in \partial f_i(x)$

Proof

$$\begin{aligned} f(y) &= h(f_1(y), \dots, f_k(y)) \\ &\geq h(f_1(x) + g_1^T(y - x), \dots, f_k(x) + g_k^T(y - x)) \\ &\geq h(f_1(x), \dots, f_k(x)) + q^T (g_1^T(y - x), \dots, g_k^T(y - x)) \\ &= f(x) + g^T(y - x) \end{aligned}$$



Optimality conditions

Fermat's optimality condition

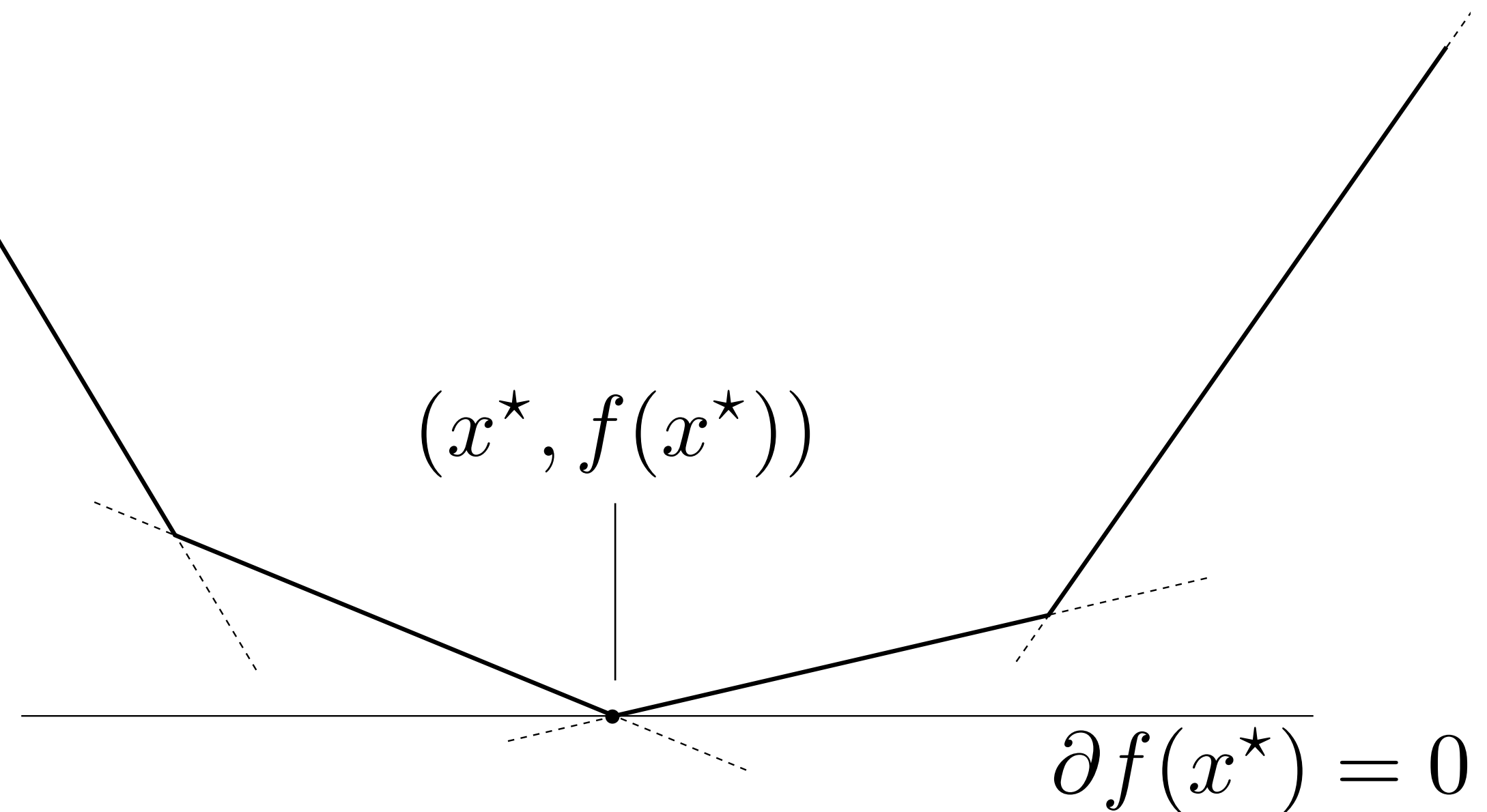
For a convex function f , then x^* is a global minimizer if and only if

$$0 \in \partial f(x^*)$$

Proof

A subgradient $g = 0$ means that, for all y

$$f(y) \geq f(x^*) + 0^T (y - x^*) = f(x^*) \quad \blacksquare$$



Note differentiable case with $\partial f(x) = \{\nabla f(x)\}$

Example: piecewise linear function

Optimality condition

$$f(x) = \max_{i=1,\dots,m} (a_i^T x + b_i) \longrightarrow 0 \in \partial f(x) = \text{conv}\{a_i \mid a_i^T x + b_i = f(x)\}$$

In other words, x^* is optimal if and only if $\exists \lambda$ such that

$$\lambda \geq 0, \quad \mathbf{1}^T \lambda = 1, \quad \sum_{i=1}^m \lambda_i a_i = 0 \quad \leftarrow (0 \in \partial f(x))$$

where $\lambda_i = 0$ if $a_i^T x^* + b_i < f(x^*)$

Same KKT optimality conditions as the primal-dual problems

$$\begin{array}{ll} \text{minimize} & t \\ \text{subject to} & Ax + b \leq t\mathbf{1} \end{array}$$

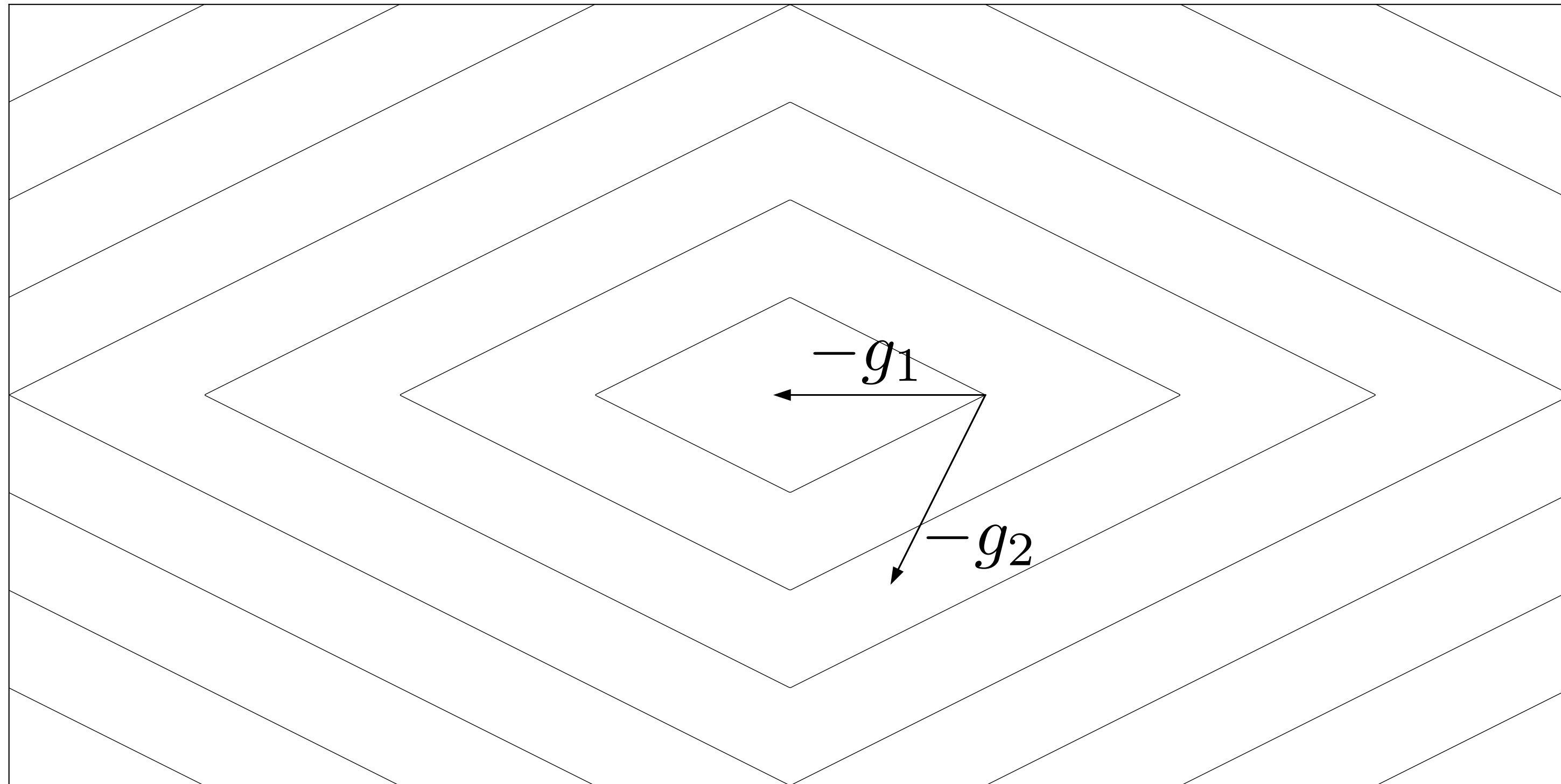
$$\begin{array}{ll} \text{maximize} & b^T \lambda \\ \text{subject to} & A^T \lambda = 0 \\ & \lambda \geq 0, \quad \mathbf{1}^T \lambda = 1 \end{array}$$

Subgradient method

Negative subgradients are not necessarily descent directions

$$f(x) = |x_1| + 2|x_2|$$

$$x = (1, 0)$$



$g_1 = (1, 0) \in \partial f(x)$ and
 $-g_1$ is a descent direction

$g_2 = (1, 2) \in \partial f(x)$ and
 $-g_2$ is not a descent direction

Subgradient method

Convex optimization problem

$$\text{minimize } f(x) \quad (\text{optimal cost } f^*)$$

Iterations

$$x^{k+1} = x^k - t_k g^k, \quad g^k \in \partial f(x^k)$$

g^k is **any subgradient** of f at x^k

Not a descent method, keep track of the best point

$$f_{\text{best}}^k = \min_{i=1, \dots, k} f(x^i)$$

Step sizes

Line search can lead to **suboptimal points**

Step sizes ***pre-specified***, not adaptively computed
(different than gradient descent)

Fixed: $t_k = t$ for $k = 0, \dots$

Diminishing: $\sum_{k=0}^{\infty} t_k^2 < \infty$, $\sum_{k=0}^{\infty} t_k = \infty$ Square summable but not summable
(goes to 0 but not too fast)
e.g., $t_k = O(1/k)$

Convergence

Assumptions

- f is convex with $\text{dom } f = \mathbf{R}^n$
- $f(x^*) > -\infty$ (finite optimal value)
- f is Lipschitz continuous with constant $G > 0$, i.e.

$$|f(x) - f(y)| \leq G\|x - y\|_2, \quad \forall x, y$$

which is equivalent to $\|g\|_2 \leq G, \quad \forall g \in \partial f(x), \forall x$

Convergence

Lipschitz continuity equivalence

f is Lipschitz continuous with constant $G > 0$, i.e.

$$|f(x) - f(y)| \leq G\|x - y\|_2, \quad \forall x, y$$

which is equivalent to $\|g\|_2 \leq G, \quad \forall g \in \partial f(x), \forall x$

Proof

If $\|g\| \leq G$ for all subgradients, pick $x, g_x \in \partial f(x)$ and $y, g_y \in \partial f(y)$. Then,

$$\begin{aligned} g_x^T(x - y) &\geq f(x) - f(y) \geq g_y^T(x - y) \\ \implies G\|x - y\|_2 &\geq f(x) - f(y) \geq -G\|x - y\|_2 \end{aligned}$$

If $\|g\|_2 > G$ for some $g \in \partial f(x)$. Take $y = x + g/\|g\|_2$ such that $\|x - y\|_2 = 1$:

$$f(y) \geq f(x) + g^T(y - x) = f(x) + \|g\|_2 > f(x) + G \quad \blacksquare$$

Convergence

Theorem

Given a convex, G -Lipschitz continuous f with finite optimal value, the subgradient method obeys

$$f_{\text{best}}^k - f^* \leq \frac{R^2 + G^2 \sum_{i=0}^k t_i^2}{2 \sum_{i=0}^k t_i}$$

where $\|x^0 - x^*\|_2 \leq R$

Convergence

Proof

Key quantity: euclidean distance to optimal set

(not function value since it can go up and down)

$$\begin{aligned}\|x^{k+1} - x^*\|_2^2 &= \|x^k - t_k g^k - x^*\|_2^2 \\ &= \|x^k - x^*\|_2^2 - 2t_k (g^k)^T (x^k - x^*) + t_k^2 \|g^k\|_2^2 \\ &\leq \|x^k - x^*\|_2^2 - 2t_k (f(x^k) - f^*) + t_k^2 \|g^k\|_2^2\end{aligned}$$

using subgradient definition $f^* = f(x^*) \geq f(x^k) + (g^k)^T (x^* - x^k)$

Convergence

Proof (continued)

Combine inequalities for $i = 0, \dots, k$

$$\begin{aligned}\|x^{k+1} - x^*\|_2^2 &\leq \|x^0 - x^*\|_2^2 - 2 \sum_{i=0}^k t_i (f(x^i) - f^*) + \sum_{i=0}^k t_i^2 \|g^i\|_2^2 \\ &\leq R^2 - 2 \sum_{i=0}^k t_i (f(x^i) - f^*) + G^2 \sum_{i=0}^k t_i^2\end{aligned}$$

Using $\|x^{k+1} - x^*\|_2^2 \geq 0$ we get

$$2 \sum_{i=0}^k t_i (f(x^i) - f^*) \leq R^2 + G^2 \sum_{i=0}^k t_i^2$$

Convergence

Proof (continued)

$$2 \sum_{i=0}^k t_i (f(x^i) - f^*) \leq R^2 + G^2 \sum_{i=0}^k t_i^2$$

Combine it with

$$\sum_{i=0}^k t_i (f(x^i) - f(x^*)) \geq \left(\sum_{i=0}^k t_i \right) \min_{i=0, \dots, k} (f(x^i) - f^*) = \left(\sum_{i=0}^k t_i \right) (f_{\text{best}}^k - f^*)$$

to get

$$f_{\text{best}}^k - f^* \leq \frac{R^2 + G^2 \sum_{i=0}^k t_i^2}{2 \sum_{i=0}^k t_i}$$



Implications for step size rules

$$f_{\text{best}}^k - f^* \leq \frac{R^2 + G^2 \sum_{i=0}^k t_i^2}{2 \sum_{i=0}^k t_i}$$

Fixed: $t_k = t$ for $k = 0, \dots$

$$f_{\text{best}}^k - f^* \leq \frac{R^2 + G^2(k+1)t^2}{2(k+1)t}$$

May be suboptimal

$$\lim_{k \rightarrow \infty} f_{\text{best}}^k \leq f^* + \frac{G^2 t}{2}$$

Diminishing:

$$\sum_{k=0}^{\infty} t_k^2 < \infty, \quad \sum_{k=0}^{\infty} t_k = \infty$$

Optimal

$$\lim_{k \rightarrow \infty} f_{\text{best}}^k = f^*$$

e.g., $t_k = \tau/(k+1)$ or $t_k = \tau/\sqrt{k+1}$

Optimal step size and convergence rate

For a tolerance $\epsilon > 0$, let's find the optimal t_k for a fixed k :

$$\frac{R^2 + G^2 \sum_{i=0}^k t_i^2}{2 \sum_{i=0}^k t_i} \leq \epsilon$$

Convex and symmetric in (t_0, \dots, t_k)

Hence, minimum when $t_i = t$

$$\longrightarrow \frac{R^2 + G^2(k+1)t^2}{2(k+1)t}$$

Optimal choice $t = \frac{R}{G\sqrt{k+1}}$

Convergence rate

$$f_{\text{best}}^k - f^* \leq \frac{RG}{\sqrt{k+1}}$$

Iterations required

$$k = O(1/\epsilon^2)$$

(gradient descent $k = O(1/\epsilon)$)

Stopping criterion

Terminating when

$$\frac{R^2 + G^2 \sum_{i=0}^k t_i^2}{2 \sum_{i=0}^k t_i} \leq \epsilon$$

is really, really slow.

Bad news

There is not really a good stopping criterion for the subgradient method

Optimal step size when f^* is known

Polyak step size

$$t_k = \frac{f(x^k) - f^*}{\|g^k\|_2^2}$$

Motivation: minimize righthand side of

$$\|x^{k+1} - x^*\|_2^2 \leq \|x^k - x^*\|_2^2 - 2t_k(f(x^k) - f^*) + t_k^2 \|g^k\|_2^2$$

Obtaining $(f(x^k) - f^*)^2 \leq (\|x^{k+1} - x^*\|_2^2 - \|x^k - x^*\|_2^2) G^2$

Applying recursively, $f_{\text{best}}^k - f^* \leq \frac{GR}{\sqrt{k+1}}$

Iterations required

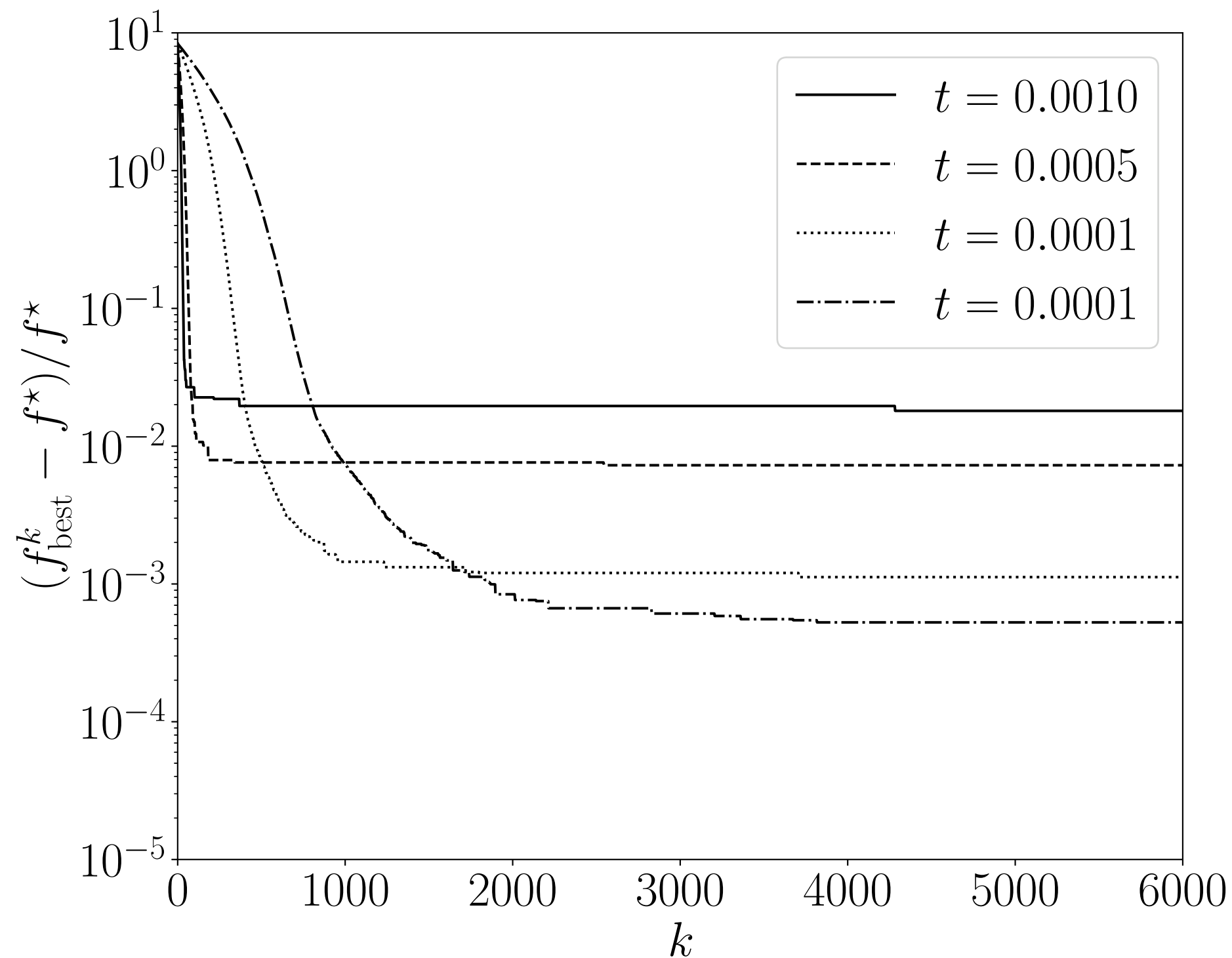
$$k = O(1/\epsilon^2)$$

still slow

Example: 1-norm minimization

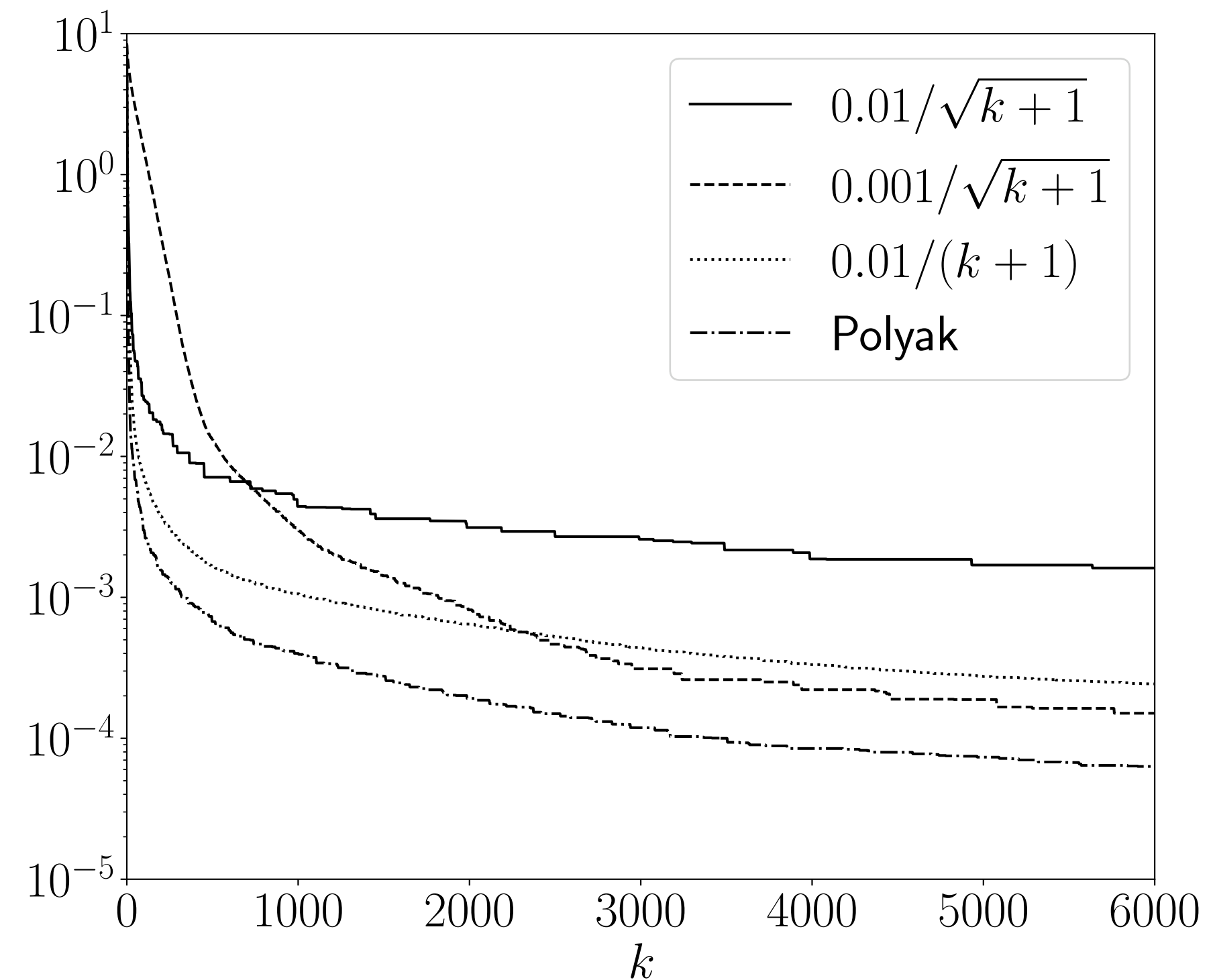
minimize $f(x) = \|Ax - b\|_1$

Fixed step size



$g = A^T \text{sign}(Ax - b) \in \partial f(x)$

Diminishing step size



Efficient packages to automatically compute (sub)gradients:

Python: JAX, PyTorch

Julia: Zygote.jl, ForwardDiff.jl, ReverseDiff.jl

Summary subgradient method

- Simple
- Handles general nondifferentiable convex functions
- Very slow convergence $O(1/\epsilon^2)$
- No good stopping criterion

Can we do better?

Can we incorporate constraints?

Subgradient methods

Today, we learned to:

- **Define** subgradients
- **Apply** subgradient calculus
- **Derive** optimality conditions from subgradients
- **Define** subgradient method and **analyze** its convergence

Next lecture

- Proximal algorithms