

# **ORF522 – Linear and Nonlinear Optimization**

## **9. Sensitivity analysis for linear optimization**

# Ed Forum

$X_B$  in BASIS  
 $X_N$  NOT in the BASIS  $\Rightarrow \bar{c}$ ?

- Dual simplex applications?

- In the dual simplex part, we talked about in primal simplex,  $X_B > 0$  and  $X_N = 0$ . However, I can confused about why in dual problem  $\bar{c}_B = 0$  and  $\bar{c}_N > 0$ . Is there any intuition behind this?

- In the illustration of dual simplex method, we used the fact that if  $y = -A_b^{-T}c_b$ , then  $A^T y + c \geq 0$  is equivalent to reduced cost  $\geq 0$ . From there (page 32), we seem to be constantly using  $A^T y + c$  as the vector of reduced cost. However, I'm wondering why we can use this previous assumption in all our steps during the dual simplex. Why is  $y = -A_b^{-T}c_b$  satisfied at all such middle steps or is this something only satisfied at the optimal solution?

**Recap**

# Reduced costs

## Interpretation

Change in objective/marginal cost of adding  $x_j$  to the basis

$$\bar{c}_j = c_j - c_B^T A_B^{-1} A_j$$

- $\bar{c}_j > 0$ : adding  $x_j$  will increase the objective (bad)
- $\bar{c}_j < 0$ : adding  $x_j$  will decrease the objective (good)

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Cost per-unit increase  
of variable  $x_j$

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Cost to change other variables  
compensating for  $x_j$   
to enforce  $Ax = b$

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# Reduced costs

$$\begin{aligned} \min \quad & c^T x \\ \text{st.} \quad & Ax = b \\ & x \geq 0 \quad x \in \mathbb{R}^n \end{aligned}$$

$$\times \begin{cases} x_B \geq 0 \\ x_N = 0 \end{cases} \rightarrow \bar{c}_B = 0$$

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## Reduced costs for basic variables is 0

$$\begin{aligned} \bar{c}_{B(i)} &= c_{B(i)} - c_B^T A_B^{-1} A_{B(i)} = c_{B(i)} - c_B^T (A_B^{-1} A_B) e_i \\ &= c_{B(i)} - c_B^T e_i = c_{B(i)} - c_{B(i)} = 0 \end{aligned}$$

# BASIC FEASIBLE SOLUTION

$$x \quad x_B > 0 \quad \rightarrow \quad \bar{c}_B = 0$$

$$\downarrow \quad x_N = 0 \\ g = -A_B^{-T} c_B$$

1) Primal Simplex

$$\exists i \quad \bar{c}_i < 0$$

BEFORE  
CONVERGENCE

Dual Problem

$$\max -b^T y$$

$$\text{st. } A^T y + c \geq 0$$

2) Dual Simplex

$$\bar{c} = A^T y + c \geq 0$$

$$L \rightarrow \bar{c}_N > 0$$

ALWAYS  
DUAL  
FEASIBLE



# Vector of reduced costs

## Reduced costs

$$\bar{c}_j = c_j - c_B^T A_B^{-1} A_j$$

## Full vector in one shot?

$$\bar{c} = (\bar{c}_1, \dots, \bar{c}_n)$$

# Vector of reduced costs

## Reduced costs

$$\bar{c}_j = c_j - c_B^T A_B^{-1} A_j$$

Isolate basis  $B$ -related components  $p$   
(they are the same across  $j$ )

$$\bar{c}_j = c_j - A_j^T (A_B^{-1})^T c_B = c_j - A_j^T p$$

## Full vector in one shot?

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# Vector of reduced costs

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## Full vector in one shot?

$$\bar{c} = (\bar{c}_1, \dots, \bar{c}_n)$$

Obtain  $p$  by solving linear system

$$p = (A_B^{-1})^T c_B \quad \Rightarrow \quad A_B^T p = c_B$$

Note:  $(M^{-1})^T = (M^T)^{-1}$   
for any square invertible  $M$

# Vector of reduced costs

## Reduced costs

$$\bar{c}_j = c_j - c_B^T A_B^{-1} A_j$$

Isolate basis  $B$ -related components  $p$   
(they are the same across  $j$ )

$$\bar{c}_j = c_j - A_j^T (A_B^{-1})^T c_B = c_j - A_j^T p$$

$$p = -y$$

## Full vector in one shot?

$$\bar{c} = (\bar{c}_1, \dots, \bar{c}_n)$$

Obtain  $p$  by solving linear system

$$p = (A_B^{-1})^T c_B \Rightarrow A_B^T p = c_B$$

Note:  $(M^{-1})^T = (M^T)^{-1}$   
for any square invertible  $M$

## Computing reduced cost vector

1. Solve  $A_B^T p = c_B$
2.  $\bar{c} = c - A^T p$

# Primal and dual basic feasible solutions

## Primal problem

$$\begin{aligned} &\text{minimize} && c^T x \\ &\text{subject to} && Ax = b \\ &&& x \geq 0 \end{aligned}$$

## Dual problem

$$\begin{aligned} &\text{maximize} && -b^T y \\ &\text{subject to} && A^T y + c \geq 0 \end{aligned}$$

Given a **basis** matrix  $A_B$

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Given a **basis** matrix  $A_B$

$$\text{Primal feasible: } Ax = b, x \geq 0 \quad \Rightarrow \quad x_B = A_B^{-1} b \geq 0$$

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**Reduced costs**



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Given a **basis** matrix  $A_B$

**Primal feasible:**  $Ax = b, x \geq 0 \Rightarrow x_B = A_B^{-1} b \geq 0$

**Dual feasible:**  $A^T y + c \geq 0$ . If  $y = -A_B^{-T} c_B \Rightarrow c - A^T A_B^{-T} c_B \geq 0$

**Reduced costs**



**Zero duality gap:**  $c^T x + b^T y = c_B^T x_B - b^T A_B^{-T} c_B = c_B^T x_B - c_B^T A_B^{-1} b = 0$

# Primal and dual basic feasible solutions

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Given a **basis** matrix  $A_B$

**Primal feasible:**  $Ax = b, x \geq 0 \Rightarrow x_B = A_B^{-1} b \geq 0$

**Dual feasible:**  $A^T y + c \geq 0$ . If  $y = -A_B^{-T} c_B \Rightarrow c - A^T A_B^{-T} c_B \geq 0$

**Zero duality gap:**  $c^T x + b^T y = c_B^T x_B - b^T A_B^{-T} c_B = c_B^T x_B - c_B^T A_B^{-1} b = 0$

↑  
(by construction)

**Reduced costs**

# Today's lecture

[Chapter 5, LO]

## Sensitivity analysis in linear optimization

- Adding new constraints and variables
- Change problem data
- Differentiable optimization

**Adding new constraints and  
variables**

# Adding new variables

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \geq 0 \end{array}$$

Solution  $x^*, y^*$

# Adding new variables

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \geq 0 \end{array} \longrightarrow \begin{array}{ll} \text{minimize} & c^T x + c_{n+1} x_{n+1} \\ \text{subject to} & Ax + A_{n+1} x_{n+1} = b \\ & x, x_{n+1} \geq 0 \end{array}$$

Solution  $x^*, y^*$

# Adding new variables

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Solution  $x^*, y^*$

Solution  $(x^*, 0), y^*$  **optimal** for the new problem?



# Adding new variables

## Optimality conditions

minimize  $c^T x + c_{n+1} x_{n+1}$

subject to  $Ax + \cancel{A_{n+1} x_{n+1}} = b \longrightarrow$  Solution  $(x^*, 0)$  is still **primal feasible**

$x, \cancel{x_{n+1}} \geq 0$

# Adding new variables

## Optimality conditions

$$\begin{aligned} \max. \quad & -b^T y \\ \text{st.} \quad & A^T y + c \geq 0 \\ & A_{n+1}^T y + c_{n+1} \geq 0 \end{aligned}$$

$$\text{minimize} \quad c^T x + c_{n+1} x_{n+1}$$

$$\text{subject to} \quad Ax + A_{n+1} x_{n+1} = b \longrightarrow \text{Solution } (x^*, 0) \text{ is still } \mathbf{\text{primal feasible}}$$

$$x, x_{n+1} \geq 0$$

Is  $y^*$  still dual feasible?

$$A_{n+1}^T y^* + c_{n+1} \geq 0$$

# Adding new variables

## Optimality conditions

minimize  $c^T x + c_{n+1} x_{n+1}$

subject to  $Ax + A_{n+1} x_{n+1} = b \longrightarrow$  Solution  $(x^*, 0)$  is still **primal feasible**

$$x, x_{n+1} \geq 0$$

Is  $y^*$  still **dual feasible**?

$$A_{n+1}^T y^* + c_{n+1} \geq 0$$

**Yes**

$(x^*, 0)$  still **optimal** for new problem

**Otherwise**

Primal simplex

# Adding new variables

## Example

$$\text{minimize} \quad -60x_1 - 30x_2 - 20x_3$$

$$\text{subject to} \quad 8x_1 + 6x_2 + x_3 \leq 48$$

$$4x_1 + 2x_2 + 1.5x_3 \leq 20$$

$$2x_1 + 1.5x_2 + 0.5x_3 \leq 8$$

$$x \geq 0$$

# Adding new variables

## Example

$$\begin{array}{ll} \text{minimize} & -60x_1 - 30x_2 - 20x_3 \\ \text{subject to} & 8x_1 + 6x_2 + x_3 \leq 48 \\ & 4x_1 + 2x_2 + 1.5x_3 \leq 20 \\ & 2x_1 + 1.5x_2 + 0.5x_3 \leq 8 \\ & x \geq 0 \end{array} \quad \text{-profit}$$

# Adding new variables

## Example

minimize

$$-60x_1 - 30x_2 - 20x_3$$

-profit

subject to

$$8x_1 + 6x_2 + x_3 \leq 48$$

material

$$4x_1 + 2x_2 + 1.5x_3 \leq 20$$

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# Adding new variables

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$$-60x_1 - 30x_2 - 20x_3$$

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subject to

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# Adding new variables

## Example

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$$x \geq 0$$

-profit

material

production

quality control



# Adding new variables

## Example

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-profit  
material  
production  
quality control

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \geq 0 \end{array}$$

$$c = (-60, -30, -20, 0, 0, 0)$$

$$A = \begin{bmatrix} 8 & 6 & 1 & 1 & 0 & 0 \\ 4 & 2 & 1.5 & 0 & 1 & 0 \\ 2 & 1.5 & 0.5 & 0 & 0 & 1 \end{bmatrix}$$

$$b = (48, 20, 8)$$

# Adding new variables

## Example

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$$b = (48, 20, 8)$$

$$x^* = (2, 0, 8, 24, 0, 0), \quad y^* = (0, 10, 10), \quad c^T x^* = -280, \quad \text{basis } \{1, 3, 4\}$$

# Adding new variables

Example: add new product?

minimize  $c^T x + c_{n+1} x_{n+1}$   
subject to  $Ax + A_{n+1} x_{n+1} = b$   
 $x, x_{n+1} \geq 0$

$$c = (-60, -30, -20, 0, 0, 0, -15)$$

$$A = \begin{bmatrix} 8 & 6 & 1 & 1 & 0 & 0 & 1 \\ 4 & 2 & 1.5 & 0 & 1 & 0 & 1 \\ 2 & 1.5 & 0.5 & 0 & 0 & 1 & 1 \end{bmatrix}$$

$$b = (48, 20, 8)$$

# Adding new variables

## Example: add new product?

$$\begin{aligned} &\text{minimize} && c^T x + c_{n+1} x_{n+1} \\ &\text{subject to} && Ax + A_{n+1} x_{n+1} = b \\ &&& x, x_{n+1} \geq 0 \end{aligned}$$

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$$b = (48, 20, 8)$$

## Previous solution

$$x^* = (2, 0, 8, 24, 0, 0), \quad y^* = (0, 10, 10), \quad c^T x^* = -280, \quad \text{basis } \{1, 3, 4\}$$

# Adding new variables

## Example: add new product?

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$$b = (48, 20, 8)$$

## Previous solution

$$x^* = (2, 0, 8, 24, 0, 0), \quad y^* = (0, 10, 10), \quad c^T x^* = -280, \quad \text{basis } \{1, 3, 4\}$$

## Still optimal

$$A_{n+1}^T y^* + c_{n+1} = \begin{bmatrix} 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} 0 \\ 10 \\ 10 \end{bmatrix} - 15 = 5 \geq 0$$

# Adding new variables

## Example: add new product?

minimize  $c^T x + c_{n+1} x_{n+1}$   
subject to  $Ax + A_{n+1} x_{n+1} = b$   
 $x, x_{n+1} \geq 0$

$$c = (-60, -30, -20, 0, 0, 0, -15)$$

$$A = \begin{bmatrix} 8 & 6 & 1 & 1 & 0 & 0 & 1 \\ 4 & 2 & 1.5 & 0 & 1 & 0 & 1 \\ 2 & 1.5 & 0.5 & 0 & 0 & 1 & 1 \end{bmatrix}$$

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## Still optimal

$$A_{n+1}^T y^* + c_{n+1} = \begin{bmatrix} 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} 0 \\ 10 \\ 10 \end{bmatrix} - 15 = 5 \geq 0$$

**Shall we add a new product?**

$$x = (x^*, 0), y^*$$

$$c^T x + b^T y = 0$$

ADDED TO PREVIOUS COST

# Adding new constraints

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \geq 0 \end{array}$$

Solution  $x^*, y^*$

# Adding new constraints

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \geq 0 \end{array}$$

Solution  $x^*, y^*$



$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & a_{m+1}^T x = b_{m+1} \\ & x \geq 0 \end{array}$$



# Adding new constraints

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \geq 0 \end{array} \longrightarrow \begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & a_{m+1}^T x = b_{m+1} \\ & x \geq 0 \end{array}$$

Solution  $x^*, y^*$

## Dual

$$\begin{array}{ll} \text{maximize} & -b^T y - b_{m+1} y_{m+1} \\ \text{subject to} & A^T y + a_{m+1} y_{m+1} + c \geq 0 \end{array}$$

# Adding new constraints

$$\begin{array}{ll}
 \text{minimize} & c^T x \\
 \text{subject to} & Ax = b \\
 & x \geq 0
 \end{array}
 \longrightarrow
 \begin{array}{ll}
 \text{minimize} & c^T x \\
 \text{subject to} & Ax = b \\
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Solution  $x^*, y^*$

## Dual

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 \end{array}$$

Solution  $x^*, (y^*, 0)$  **optimal** for the new problem?

# Adding new constraints

## Optimality conditions

maximize  $-b^T y - \cancel{b_{m+1} y_{m+1}}$   
subject to  $A^T y + \cancel{a_{m+1} y_{m+1}} + c \geq 0 \longrightarrow$  Solution  $(y^*, 0)$  is still **dual feasible**

# Adding new constraints

## Optimality conditions

maximize  $-b^T y$   
subject to  $A^T y + a_{m+1} y_{m+1} + c \geq 0 \longrightarrow$  Solution  $(y^*, 0)$  is still **dual feasible**

Is  $x^*$  still **primal feasible**?

$$Ax = b$$

$$a_{m+1}^T x = b_{m+1}$$

$$x \geq 0$$

# Adding new constraints

## Optimality conditions

maximize  $-b^T y$   
subject to  $A^T y + a_{m+1} y_{m+1} + c \geq 0 \longrightarrow$  Solution  $(y^*, 0)$  is still **dual feasible**

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$$Ax = b$$

$$a_{m+1}^T x = b_{m+1}$$

$$x \geq 0$$

**Yes**

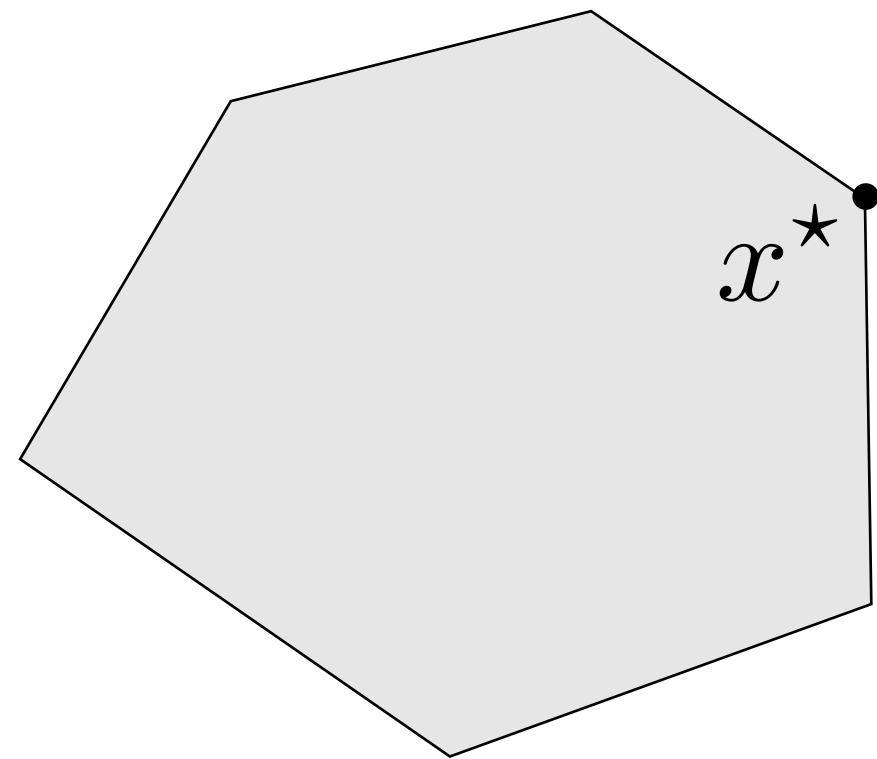
$x^*$  still **optimal** for new problem

**Otherwise**

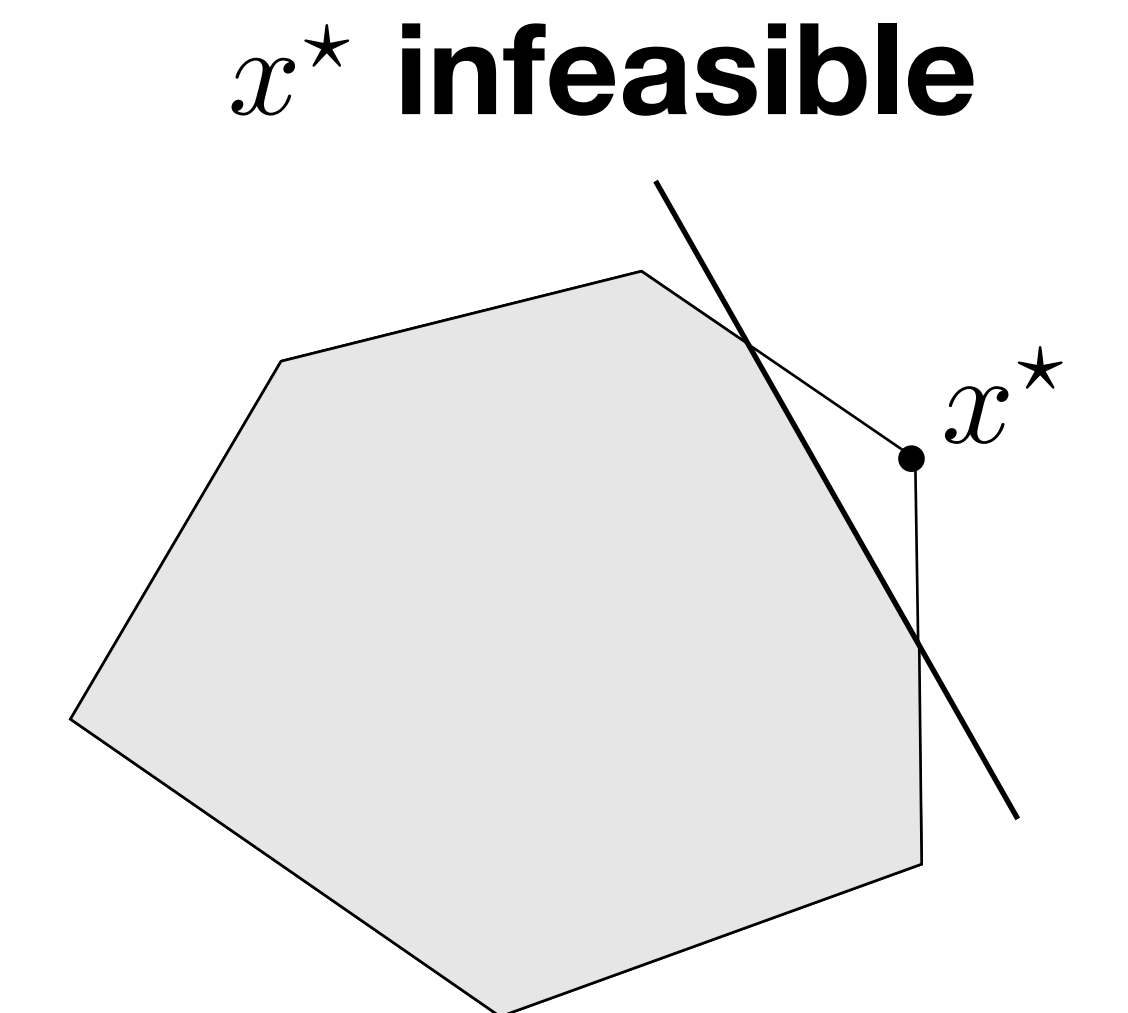
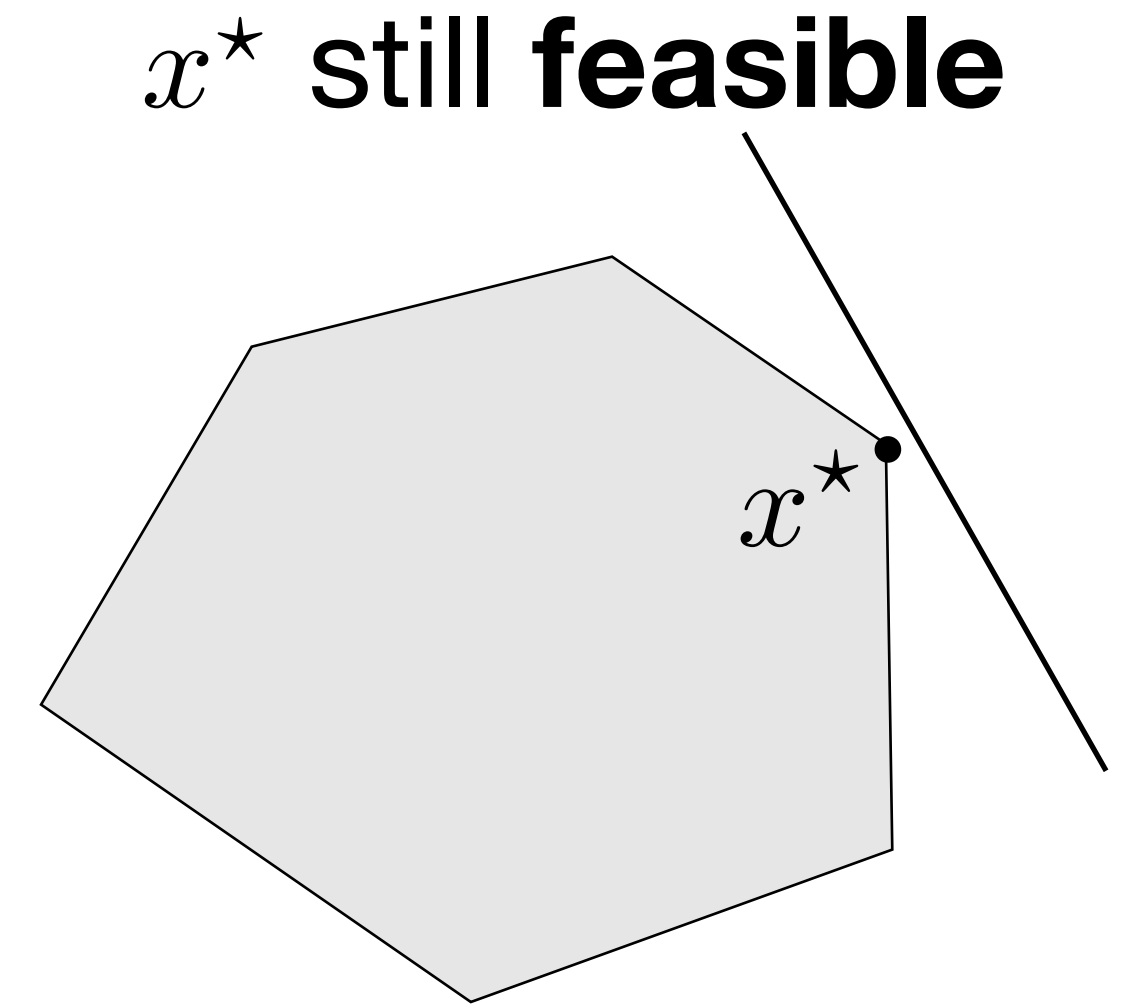
Dual simplex

# Adding new constraints

## Example



Add new constraint



# Global sensitivity analysis

# Information from primal-dual solution

**Goal:** extract information from  $x^*, y^*$  about their sensitivity with respect to changes in problem data

## Modified LP

$$\begin{aligned} &\text{minimize} && c^T x \\ &\text{subject to} && Ax = b + u \\ & && x \geq 0 \end{aligned}$$

**Optimal cost**  $p^*(u)$



# Global sensitivity

## Dual of modified LP

$$\begin{array}{ll} \text{maximize} & -(b + u)^T y \\ \text{subject to} & A^T y + c \geq 0 \end{array}$$

# Global sensitivity

## Dual of modified LP

maximize  $-(b + u)^T y$

subject to  $A^T y + c \geq 0$

*DOES NOT DEPEND ON u*

## Global lower bound

Given  $y^*$  a dual optimal solution for  $u = 0$ , then

$$\begin{aligned} p^*(u) &\geq -(b + u)^T y^* && \text{(from weak duality and} \\ &= p^*(0) - u^T y^* && \text{dual feasibility)} \end{aligned}$$

# Global sensitivity

## Dual of modified LP

$$\begin{aligned} &\text{maximize} && -(b + u)^T y \\ &\text{subject to} && A^T y + c \geq 0 \end{aligned}$$

## Global lower bound

Given  $y^*$  a dual optimal solution for  $u = 0$ , then

$$\begin{aligned} p^*(u) &\geq -(b + u)^T y^* && \text{(from weak duality and} \\ &= p^*(0) - u^T y^* && \text{dual feasibility)} \end{aligned}$$

It holds for any  $u$

# Global sensitivity

## Example

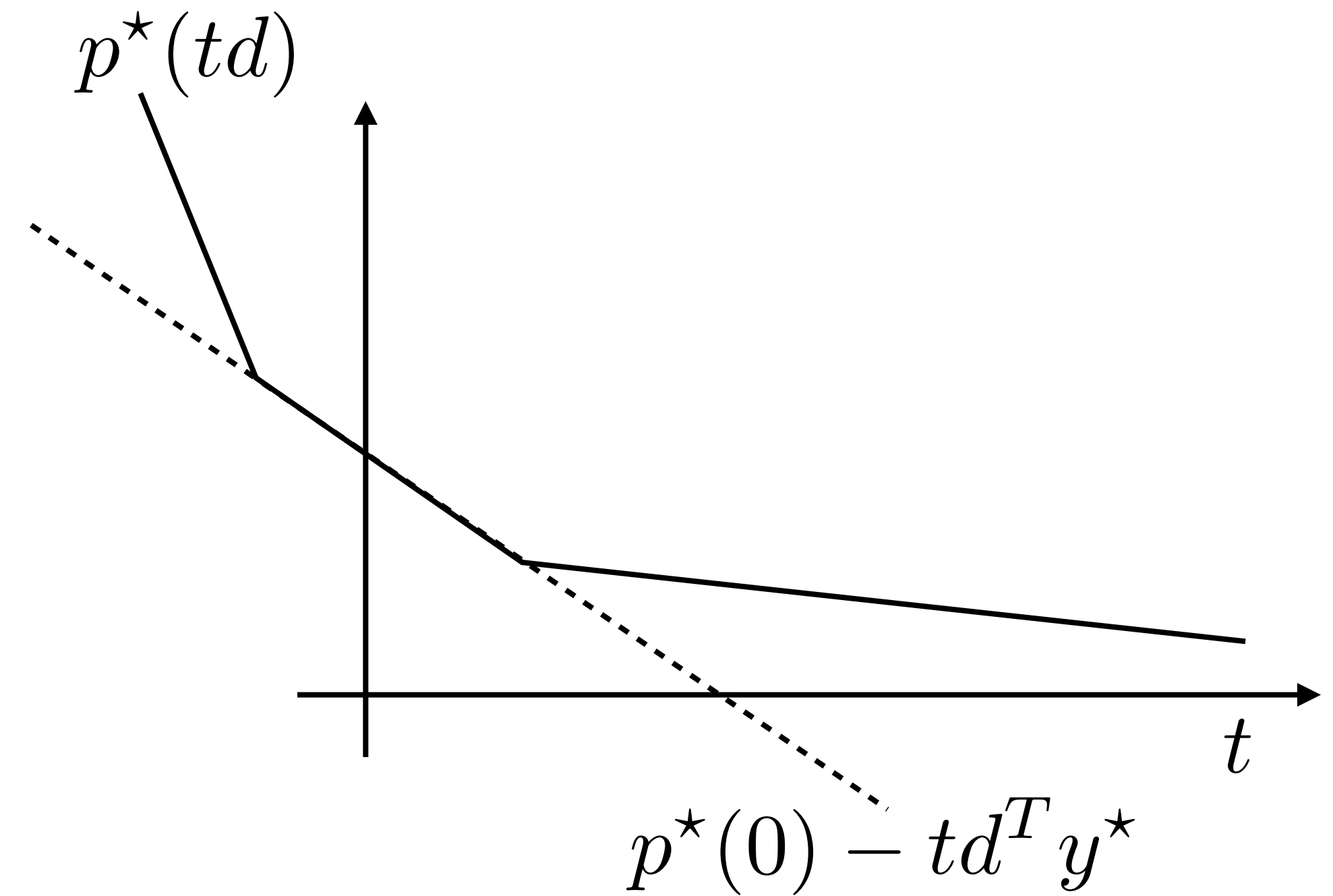
Take  $u = td$  with  $d \in \mathbf{R}^m$  fixed

$$\text{minimize } c^T x$$

$$\text{subject to } Ax = b + td$$

$$x \geq 0$$

$p^*(td)$  is the optimal value as a function of  $t$



# Global sensitivity

## Example

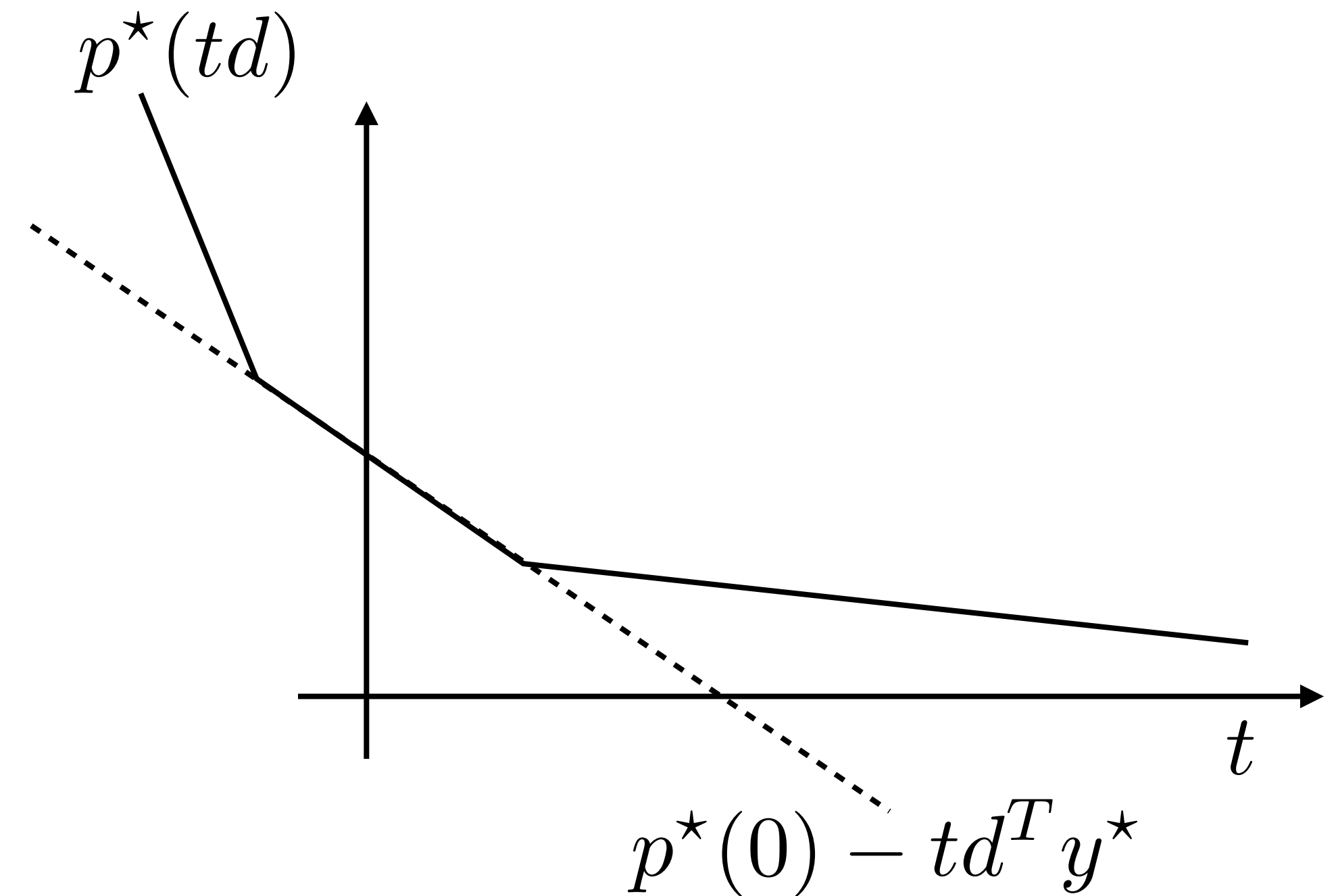
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$$x \geq 0$$

$p^*(td)$  is the optimal value as a function of  $t$



**Sensitivity information** (assuming  $d^T y^* \geq 0$ )

- $t < 0$  the optimal value increases
- $t > 0$  the optimal value decreases (not so much if  $t$  is small)

# Optimal value function

$$p^*(u) = \min\{c^T x \mid Ax = b + u, x \geq 0\}$$

**Assumption:**  $p^*(0)$  is finite

## Properties

- $p^*(u) > -\infty$  everywhere (from global lower bound)
- the domain  $\{u \mid p^*(u) < +\infty\}$  is a polyhedron
- $p^*(u)$  is piecewise-linear on its domain

# Optimal value function is piecewise linear

## Proof

$$p^*(u) = \min\{c^T x \mid Ax = b + u, x \geq 0\}$$

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**Dual feasible set**

$$D = \{y \mid A^T y + c \geq 0\}$$

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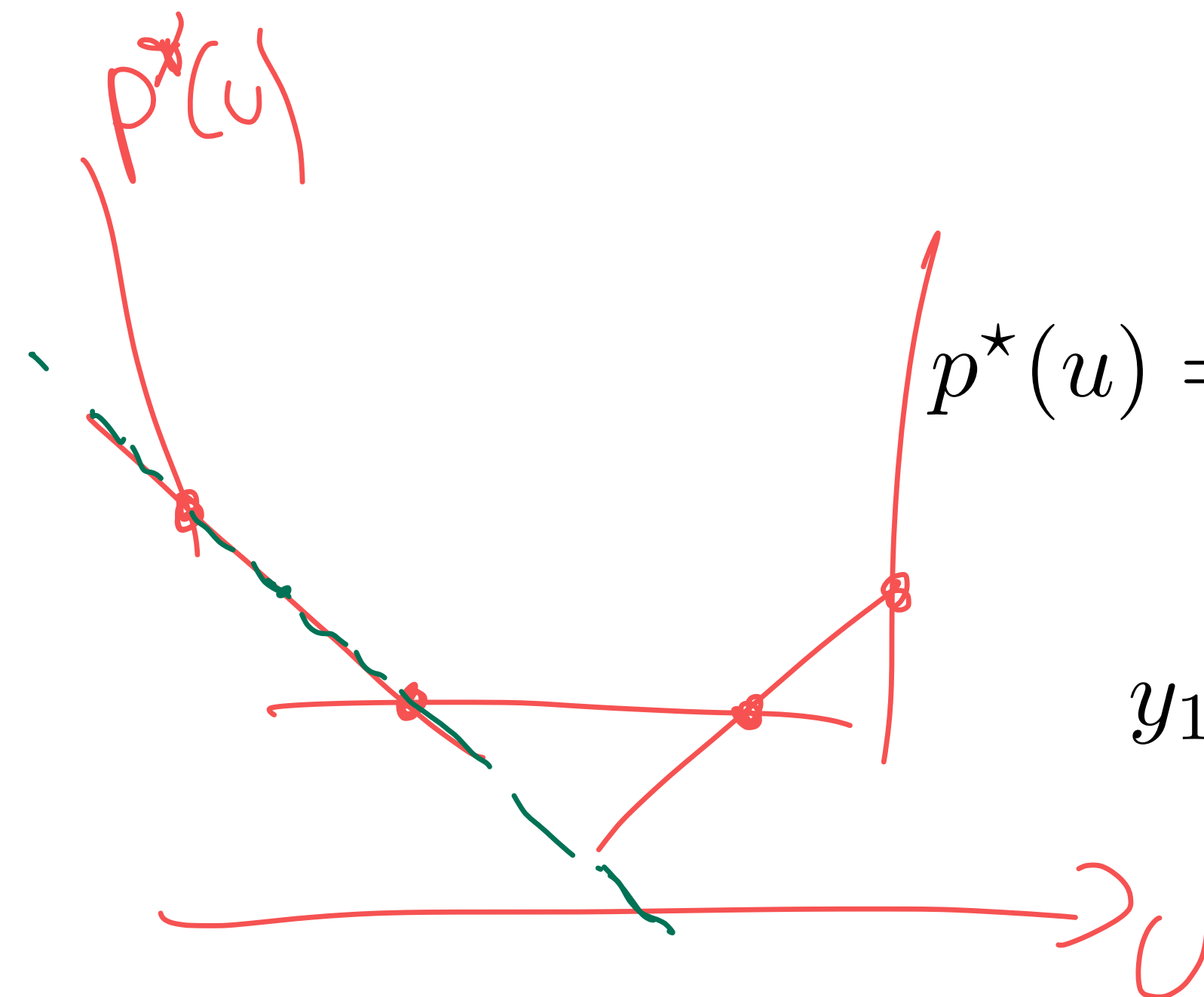
$$D = \{y \mid A^T y + c \geq 0\}$$

**Assumption:**  $p^*(0)$  is finite

If  $p^*(u)$  finite

$$p^*(u) = \max_{y \in D} -(b + u)^T y = \max_{k=1, \dots, r} -y_k^T u - b^T y_k$$

$y_1, \dots, y_r$  are the extreme points of  $D$



# Local sensitivity analysis

# Local sensitivity

$u$  in neighborhood of the origin

## Original LP

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \geq 0 \end{array} \longrightarrow$$

## Optimal solution

$$\begin{array}{ll} \text{Primal} & x_i = 0, \quad i \notin B \\ & x_B^* = A_B^{-1} b \\ \text{Dual} & y^* = -A_B^{-T} c_B \end{array}$$

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## Modified LP

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b + u \\ & x \geq 0 \end{array}$$

## Modified dual

$$\begin{array}{ll} \text{maximize} & -(b + u)^T y \\ \text{subject to} & A^T y + c \geq 0 \end{array}$$

**Optimal basis  
does not change**

# Local sensitivity

$u$  in neighborhood of the origin

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$$\begin{aligned} &\text{maximize} && -(b + u)^T y \\ &\text{subject to} && A^T y + c \geq 0 \end{aligned}$$

**Optimal basis does not change**

## Modified optimal solution

$$\begin{aligned} x_B^*(u) &= A_B^{-1} (b + u) = x_B^* + A_B^{-1} u \\ y^*(u) &= y^* \end{aligned}$$

# Derivative of the optimal value function

## Modified optimal solution

$$x_B^*(u) = A_B^{-1}(b + u) = x_B^* + A_B^{-1}u$$

$$y^*(u) = y^*$$

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## Optimal value function

$$p^*(u) = c^T x^*(u)$$

$$= c^T x^* + c_B^T A_B^{-1} u$$

$$= p^*(0) - y^{*T} u \quad (\text{affine for small } u)$$

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## Optimal value function

$$p^*(u) = c^T x^*(u)$$

$$= c^T x^* + c_B^T A_B^{-1}u$$

$$= p^*(0) - y^{*T}u \quad (\text{affine for small } u)$$

## Local derivative

$$\frac{\partial p^*(u)}{\partial u} = -y^*$$

$(y^*$  are the **shadow prices**)



# Sensitivity example

$$\begin{array}{ll} \text{minimize} & -60x_1 - 30x_2 - 20x_3 \\ \text{subject to} & 8x_1 + 6x_2 + x_3 \leq 48 \\ & 4x_1 + 2x_2 + 1.5x_3 \leq 20 \\ & 2x_1 + 1.5x_2 + 0.5x_3 \leq 8 \\ & x \geq 0 \end{array}$$

# Sensitivity example

minimize  $-60x_1 - 30x_2 - 20x_3$  -profit

subject to  $8x_1 + 6x_2 + x_3 \leq 48$

$4x_1 + 2x_2 + 1.5x_3 \leq 20$

$2x_1 + 1.5x_2 + 0.5x_3 \leq 8$

$x \geq 0$

# Sensitivity example

minimize

$$-60x_1 - 30x_2 - 20x_3$$

-profit

subject to

$$8x_1 + 6x_2 + x_3 \leq 48$$

material

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quality control

$$x \geq 0$$

$$x^* = (2, 0, 8, 24, 0, 0), \quad y^* = (0, 10, 10), \quad c^T x^* = -280, \quad \text{basis } \{1, 3, 4\}$$

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What does  $y_3^* = 10$  mean?

# Sensitivity example

$$b = \begin{bmatrix} 48 \\ 20 \\ 8 \end{bmatrix}$$

minimize  $-60x_1 - 30x_2 - 20x_3$

-profit

subject to  $8x_1 + 6x_2 + x_3 \leq 48$

material

$4x_1 + 2x_2 + 1.5x_3 \leq 20$

production

$2x_1 + 1.5x_2 + 0.5x_3 \leq 8$

quality control

$x \geq 0$

$x^* = (2, 0, 8, 24, 0, 0), \quad y^* = (0, 10, 10), \quad c^T x^* = -280, \quad \text{basis } \{1, 3, 4\}$

What does  $y_3^* = 10$  mean?

Let's increase the quality control budget by 1, i.e.,  $u = (0, 0, 1)$

$b + u$

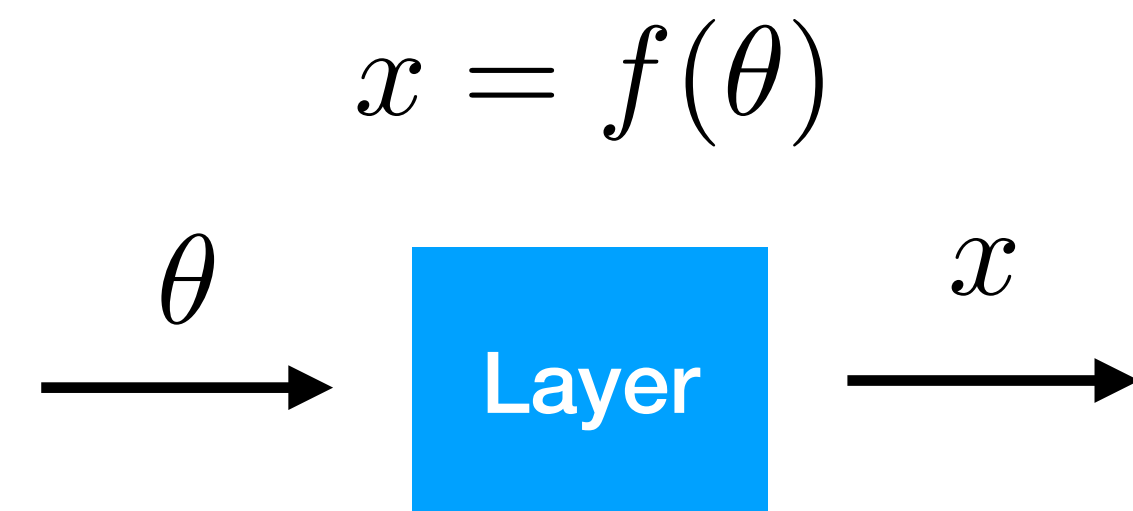
$p^*(\overset{u}{\cancel{b}}) = p^*(0) - y^{*T} u = -280 - 10 = -290$



# Differentiable optimization

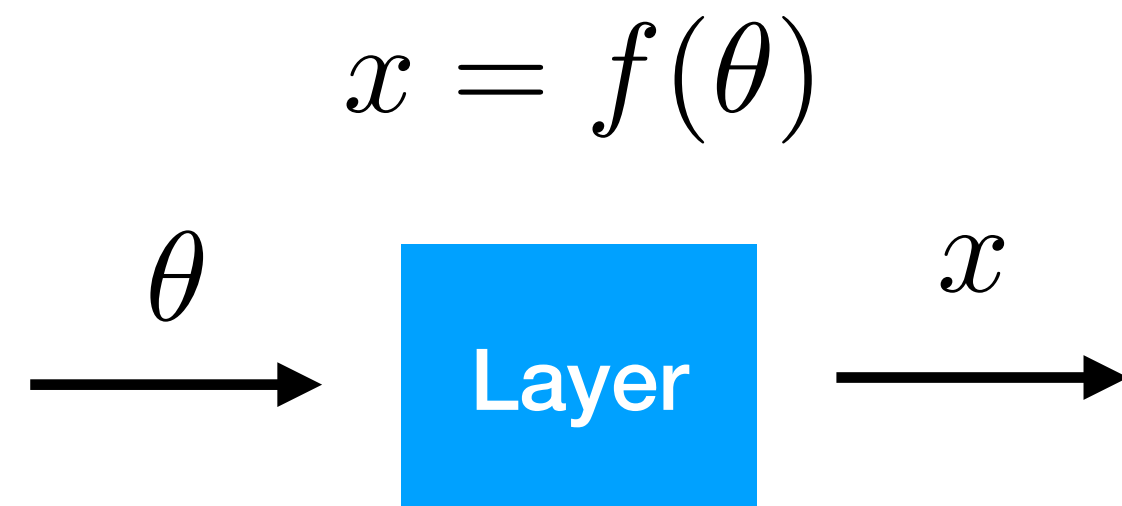
# Training a neural network

## Single layer model



# Training a neural network

## Single layer model



## Training

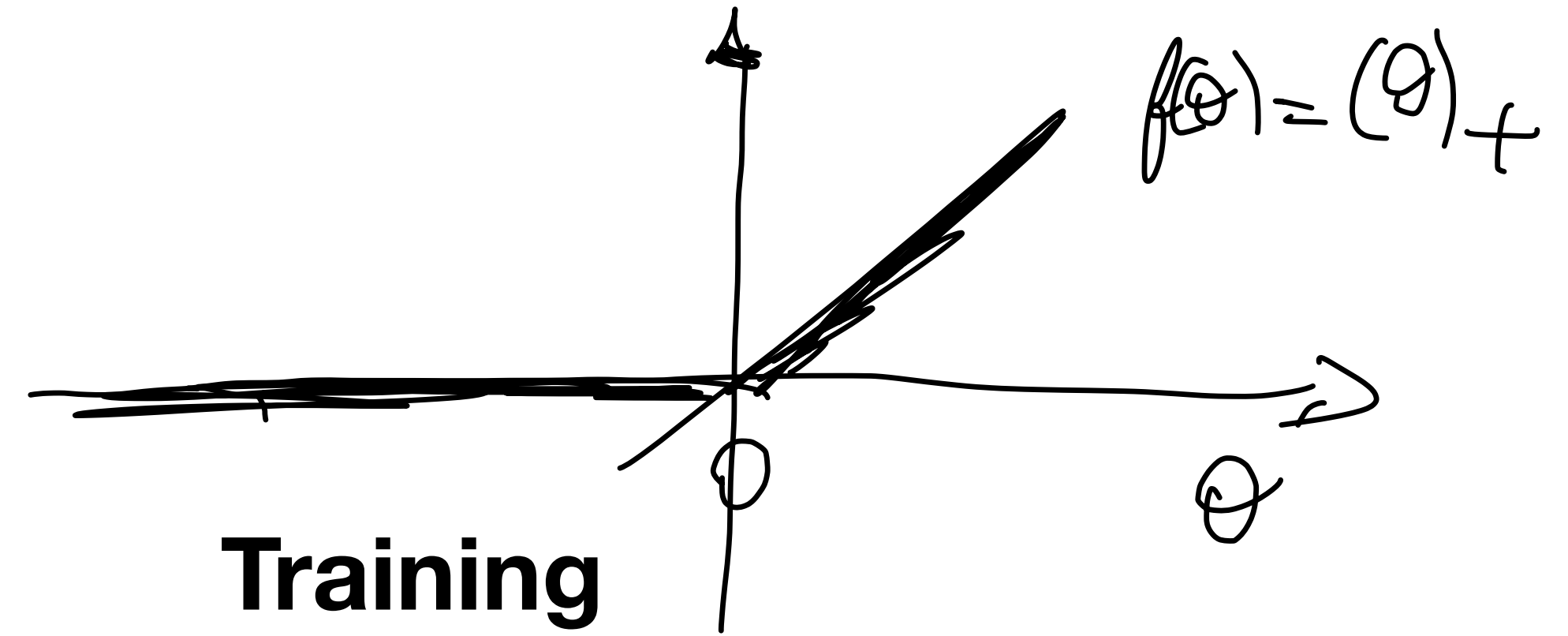
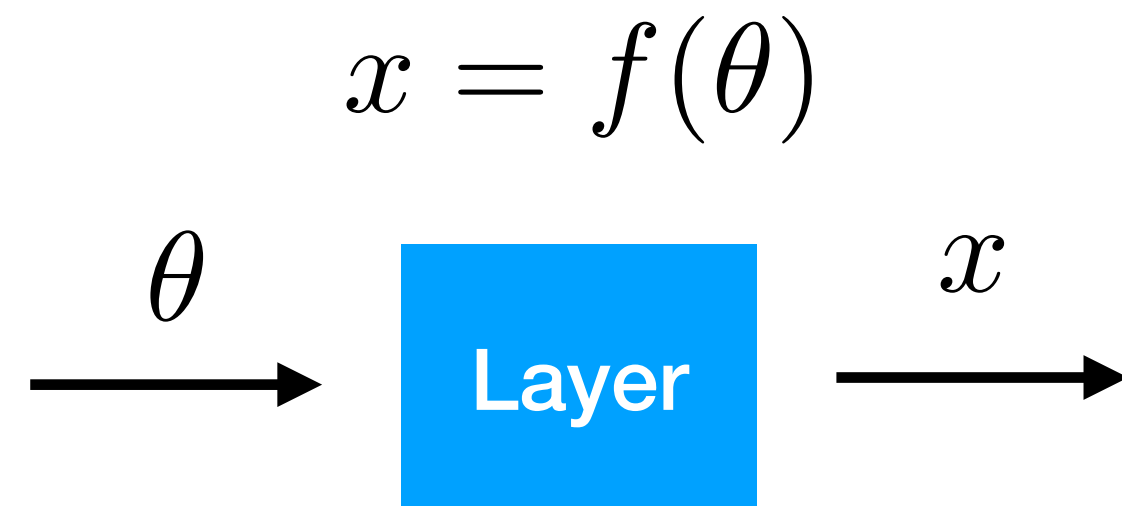
minimize  $\mathcal{L}(\theta)$

Gradient descent (more on this later)

$$\theta \leftarrow \theta - t \nabla_{\theta} \mathcal{L}(\theta)$$

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## Single layer model



**Training**

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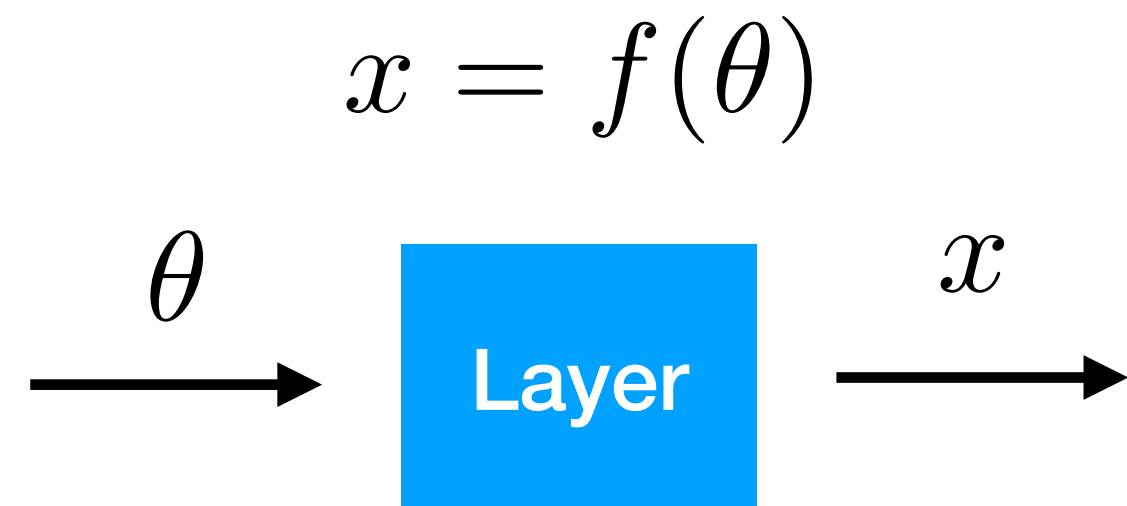
$$\theta \leftarrow \theta - t \nabla_{\theta} \mathcal{L}(\theta)$$

Sensitivity

$$\nabla_{\theta} \mathcal{L} = \left( \frac{\partial \mathcal{L}}{\partial \theta} \right)^T = \left( \frac{\partial \mathcal{L}}{\partial x} \frac{\partial x}{\partial \theta} \right)^T = \left( \frac{\partial x}{\partial \theta} \right)^T \nabla_x \mathcal{L}$$

# Training a neural network

## Single layer model



## Training

minimize  $\mathcal{L}(\theta)$

Gradient descent (more on this later)

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Can  $f$  be an **optimization problem**?

# Implicit layers

<https://implicit-layers-tutorial.org/>

$$\begin{aligned} &\text{find} && x(\theta) \\ &\text{subject to} && r(\theta, x(\theta)) = 0 \end{aligned}$$

$(x(\theta))$  is implicitly defined by  $r$

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find  $x(\theta)$   
subject to  $r(\theta, x(\theta)) = 0$   $(x(\theta)$  is implicitly defined by  $r$ )

**How do we compute derivatives?**

$$\frac{\partial x(\theta)}{\partial \theta}$$

# Implicit layers

<https://implicit-layers-tutorial.org/>

find  $x(\theta)$   
subject to  $r(\theta, x(\theta)) = 0$  ( $x(\theta)$  is implicitly defined by  $r$ )

**How do we compute derivatives?**

$$\frac{\partial x(\theta)}{\partial \theta}$$

**Implicit function theorem**

Under mild assumptions (non-singularity),

$$\frac{\partial r(\theta, x(\theta))}{\partial x} \frac{\partial x(\theta)}{\partial \theta} + \frac{\partial r(\theta, x(\theta))}{\partial \theta} = 0 \longrightarrow \frac{\partial x(\theta)}{\partial \theta} = - \left( \frac{\partial r(\theta, x(\theta))}{\partial x} \right)^{-1} \frac{\partial r(\theta, x(\theta))}{\partial \theta}$$



# Optimization layers

$$x^*(\theta) = \underset{x}{\operatorname{argmin}} \quad c^T x$$

subject to  $Ax \leq b$

Parameters:  $\theta = \{c, A, b\}$

Solution  $x^*(\theta)$

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## Features

- Add **domain knowledge** and **hard constraints**
- **End-to-end** training and optimization
- Nice theory and algorithms for general **convex optimization**
- **Applications** in RL, control, meta-learning, game theory, etc.

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- Add **domain knowledge** and **hard constraints**
- **End-to-end** training and optimization
- Nice theory and algorithms for general **convex optimization**
- **Applications** in RL, control, meta-learning, game theory, etc.

## Goal

Compute  $\frac{\partial x^*(\theta)}{\partial \theta}$

# Optimality conditions

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax \leq b \end{array}$$

Parameters:  $\theta = \{c, A, b\}$   
Solution  $x^*(\theta)$

**Solve** and obtain primal-dual pair  $x^*, y^*$  (forward-pass)

## Optimality conditions

$$\begin{aligned} A^T y + c &= 0 \\ \text{diag}(y)(Ax - b) &= 0 \\ y &\geq 0, \quad b - Ax \geq 0 \end{aligned}$$

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Mapping  $r(\theta, x(\theta)) = 0$

# Computing derivatives

**Take differentials**

$$\begin{array}{l} A^T y^* + c = 0 \\ \mathbf{diag}(y^*)(Ax - b) = 0 \end{array} \longrightarrow \begin{array}{l} \mathbf{diag}(Ax^* - b)dy + \mathbf{diag}(y^*)(dAx^* + Adx - db) + dc = 0 \end{array}$$
$$dA^T y^* + A^T dy = 0$$

# Computing derivatives

**Take differentials**

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**Linear system**

$$\begin{bmatrix} 0 & A^T \\ \mathbf{diag}(y^*)A & \mathbf{diag}(Ax^* - b) \end{bmatrix} \begin{bmatrix} dx \\ dy \end{bmatrix} = - \begin{bmatrix} dA^T y^* + dc \\ \mathbf{diag}(y^*)(dAx^* - db) \end{bmatrix}$$

# Computing derivatives

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**Linear system**

$$\begin{bmatrix} 0 & A^T \\ \mathbf{diag}(y^*)A & \mathbf{diag}(Ax^* - b) \end{bmatrix} \begin{bmatrix} dx \\ dy \end{bmatrix} = - \begin{bmatrix} dA^T y^* + dc \\ \mathbf{diag}(y^*)(dAx^* - db) \end{bmatrix}$$

**Example:** How does  $x^*$  change with  $b_1$ ?

Set  $db = e_1$ ,  $dA = 0$ ,  $dc = 0$  and solve the linear system.

The solution  $dx$  will correspond to  $\frac{\partial x}{\partial b_1}$



# Is it always differentiable?

The linear system matrix must be invertible  
(the problem must have unique solution)

$$\begin{bmatrix} 0 & A^T \\ \mathbf{diag}(y^*)A & \mathbf{diag}(Ax^* - b) \end{bmatrix} \begin{bmatrix} dx \\ dy \end{bmatrix} = - \begin{bmatrix} dA^T y^* + dc \\ \mathbf{diag}(y^*)(dAx^* - db) \end{bmatrix}$$

$M$   $q$

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$M$   $q$

**Remember.** implicit function theorem

$$\frac{\partial x(\theta)}{\partial \theta} = - \left( \frac{\partial r(\theta, x(\theta))}{\partial x} \right)^{-1} \frac{\partial r(\theta, x(\theta))}{\partial \theta}$$

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$M$   $q$

**Remember.** implicit function theorem

$$\frac{\partial x(\theta)}{\partial \theta} = - \left( \frac{\partial r(\theta, x(\theta))}{\partial x} \right)^{-1} \frac{\partial r(\theta, x(\theta))}{\partial \theta}$$

If not, **least squares** “subdifferential”

$$\text{minimize} \left\| M \begin{bmatrix} dx \\ dy \end{bmatrix} + q \right\|_2^2$$

# Example

## Learning to play Sudoku

			3
1			
		4	
4			1

2	4	1	3
1	3	2	4
3	1	4	2
4	2	3	1

### Sudoku constraint satisfaction problem

minimize  $0$

subject to  $Ax = b$

$x \geq 0, x \in \mathbf{Z}^d$

# Example

## Learning to play Sudoku

			3
1			
		4	
4			1

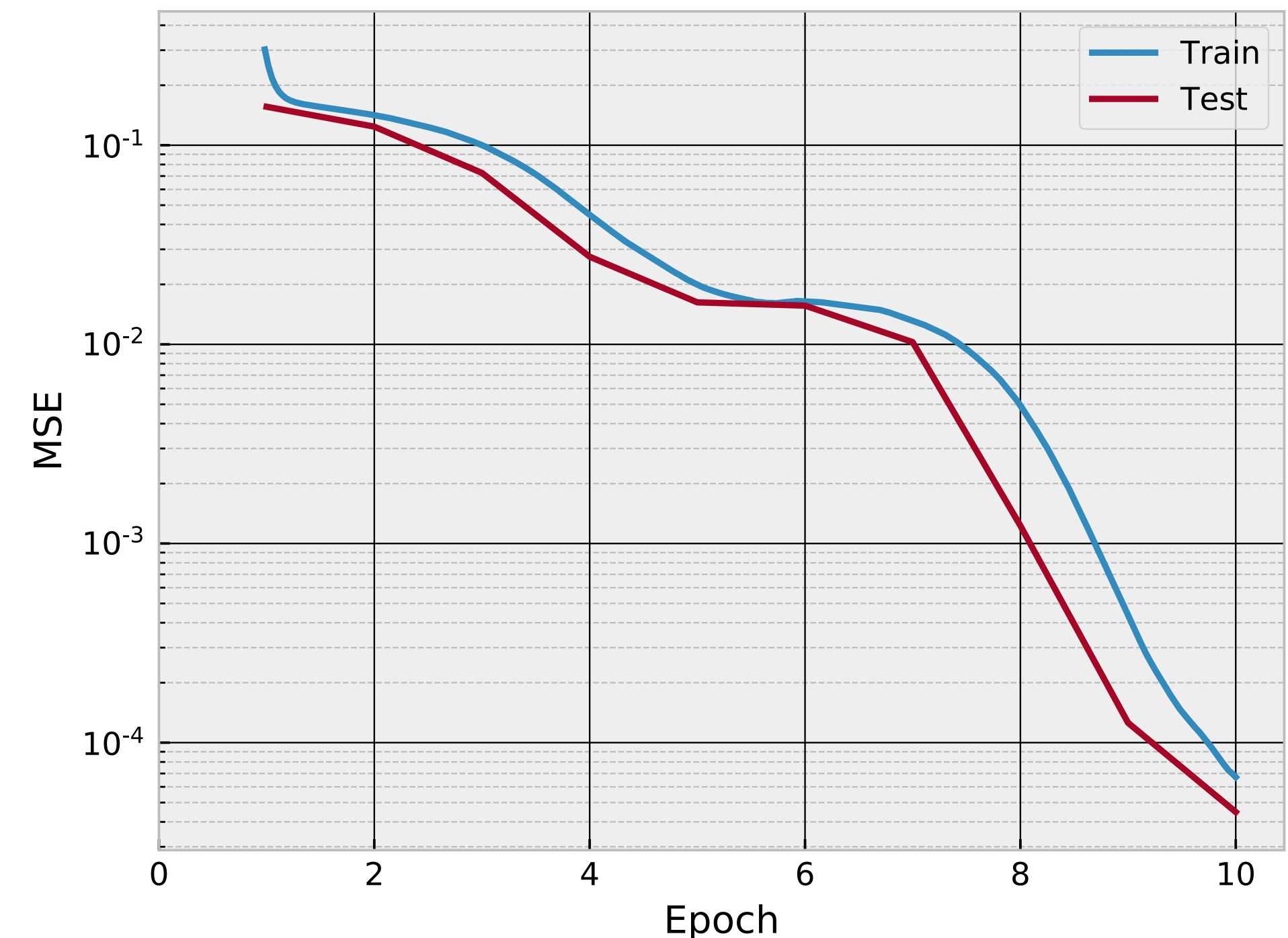
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1	3	2	4
3	1	4	2
4	2	3	1

### Sudoku constraint satisfaction problem

$$\begin{aligned} & \text{minimize} && 0 \\ & \text{subject to} && Ax = b \\ & && x \geq 0, x \in \mathbf{Z}^d \end{aligned}$$

### Linear optimization layer (parameters $\theta = \{A, b\}$ )

$$\begin{aligned} x^* = & \underset{x}{\operatorname{argmin}} && 0 \\ & \text{subject to} && Ax = b \\ & && x \geq 0 \end{aligned}$$



# Sensitivity analysis in linear optimization

Today, we learned to:

- **Use** the most appropriate primal/dual simplex algorithm when variables and/or constraints are added
- **Analyze** sensitivity of the cost with respect to change in the data
- **Apply** sensitivity analysis to differentiable linear optimization layers

# Next lecture

- Barrier methods for linear optimization