

ORF522 – Linear and Nonlinear Optimization

9. Sensitivity analysis for linear optimization

Ed Forum

- Dual simplex applications?
- In the dual simplex part, we talked about in primal simplex, $X_B > 0$ and $X_N = 0$. However, I can be confused about why in dual problem $\bar{c}_B = 0$ and $\bar{c}_N > 0$. Is there any intuition behind this?
- In the illustration of dual simplex method, we used the fact that if $y = -A_B^{-T}c_B$, then $A^T y + c >= 0$ is equivalent to reduced cost $>= 0$. From there (page 32), we seem to be constantly using $A^T y + c$ as the vector of reduced cost. However, I'm wondering why we can use this previous assumption in all our steps during the dual simplex. Why is $y = -A_B^{-T}c_B$ satisfied at all such middle steps or is this something only satisfied at the optimal solution?

Recap

Reduced costs

Interpretation

Change in objective/marginal cost of adding x_j to the basis

$$\bar{c}_j = c_j - c_B^T A_B^{-1} A_j$$

Cost per-unit increase
of variable x_j

Cost to change other variables
compensating for x_j
to enforce $Ax = b$

- $\bar{c}_j > 0$: adding x_j will increase the objective (bad)
- $\bar{c}_j < 0$: adding x_j will decrease the objective (good)

Reduced costs for basic variables is 0

$$\begin{aligned}\bar{c}_{B(i)} &= c_{B(i)} - c_B^T A_B^{-1} A_{B(i)} = c_{B(i)} - c_B^T (A_B^{-1} A_B) e_i \\ &= c_{B(i)} - c_B^T e_i = c_{B(i)} - c_{B(i)} = 0\end{aligned}$$

Vector of reduced costs

Reduced costs

$$\bar{c}_j = c_j - c_B^T A_B^{-1} A_j$$

Full vector in one shot?

$$\bar{c} = (\bar{c}_1, \dots, \bar{c}_n)$$

Isolate basis B -related components p
(they are the same across j)

$$\bar{c}_j = c_j - A_j^T (A_B^{-1})^T c_B = c_j - A_j^T p$$

Obtain p by solving linear system

$$p = (A_B^{-1})^T c_B \quad \Rightarrow \quad A_B^T p = c_B$$

Note: $(M^{-1})^T = (M^T)^{-1}$
for any square invertible M

Computing reduced cost vector

1. Solve $A_B^T p = c_B$
2. $\bar{c} = c - A^T p$

Primal and dual basic feasible solutions

Primal problem

$$\begin{aligned} &\text{minimize} && c^T x \\ &\text{subject to} && Ax = b \\ &&& x \geq 0 \end{aligned}$$

Dual problem

$$\begin{aligned} &\text{maximize} && -b^T y \\ &\text{subject to} && A^T y + c \geq 0 \end{aligned}$$

Given a **basis** matrix A_B

Primal feasible: $Ax = b, x \geq 0 \Rightarrow x_B = A_B^{-1} b \geq 0$

Dual feasible: $A^T y + c \geq 0$. If $y = -A_B^{-T} c_B \Rightarrow c - A^T A_B^{-T} c_B \geq 0$

Zero duality gap: $c^T x + b^T y = c_B^T x_B - b^T A_B^{-T} c_B = c_B^T x_B - c_B^T A_B^{-1} b = 0$

(by construction)

Today's lecture

[Chapter 5, LO]

Sensitivity analysis in linear optimization

- Adding new constraints and variables
- Change problem data
- Differentiable optimization

**Adding new constraints and
variables**

Adding new variables

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \geq 0 \end{array} \longrightarrow \begin{array}{ll} \text{minimize} & c^T x + c_{n+1} x_{n+1} \\ \text{subject to} & Ax + A_{n+1} x_{n+1} = b \\ & x, x_{n+1} \geq 0 \end{array}$$

Solution x^*, y^*

Solution $(x^*, 0), y^*$ **optimal** for the new problem?

Adding new variables

Optimality conditions

minimize $c^T x + c_{n+1} x_{n+1}$

subject to $Ax + A_{n+1} x_{n+1} = b \longrightarrow$ Solution $(x^*, 0)$ is still **primal feasible**

$$x, x_{n+1} \geq 0$$

Is y^* still **dual feasible**?

$$A_{n+1}^T y^* + c_{n+1} \geq 0$$

Yes

$(x^*, 0)$ still **optimal** for new problem

Otherwise

Primal simplex

Adding new variables

Example

$$\begin{array}{ll} \text{minimize} & -60x_1 - 30x_2 - 20x_3 \\ \text{subject to} & 8x_1 + 6x_2 + x_3 \leq 48 \\ & 4x_1 + 2x_2 + 1.5x_3 \leq 20 \\ & 2x_1 + 1.5x_2 + 0.5x_3 \leq 8 \\ & x \geq 0 \end{array}$$

-profit
material
production
quality control

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \geq 0 \end{array}$$
$$c = (-60, -30, -20, 0, 0, 0)$$
$$A = \begin{bmatrix} 8 & 6 & 1 & 1 & 0 & 0 \\ 4 & 2 & 1.5 & 0 & 1 & 0 \\ 2 & 1.5 & 0.5 & 0 & 0 & 1 \end{bmatrix}$$
$$b = (48, 20, 8)$$

$$x^* = (2, 0, 8, 24, 0, 0), \quad y^* = (0, 10, 10), \quad c^T x^* = -280, \quad \text{basis } \{1, 3, 4\}$$

Adding new variables

Example: add new product?

$$\begin{aligned} \text{minimize} \quad & c^T x + c_{n+1} x_{n+1} \\ \text{subject to} \quad & Ax + A_{n+1} x_{n+1} = b \\ & x, x_{n+1} \geq 0 \end{aligned}$$

$$c = (-60, -30, -20, 0, 0, 0, -15)$$

$$A = \begin{bmatrix} 8 & 6 & 1 & 1 & 0 & 0 & 1 \\ 4 & 2 & 1.5 & 0 & 1 & 0 & 1 \\ 2 & 1.5 & 0.5 & 0 & 0 & 1 & 1 \end{bmatrix}$$

$$b = (48, 20, 8)$$

Previous solution

$$x^* = (2, 0, 8, 24, 0, 0), \quad y^* = (0, 10, 10), \quad c^T x^* = -280, \quad \text{basis } \{1, 3, 4\}$$

Still optimal

$$A_{n+1}^T y^* + c_{n+1} = \begin{bmatrix} 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} 0 \\ 10 \\ 10 \end{bmatrix} - 15 = 5 \geq 0$$

Shall we add a new product?

Adding new constraints

$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & x \geq 0 \end{array} \longrightarrow \begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & a_{m+1}^T x = b_{m+1} \\ & x \geq 0 \end{array}$$

Solution x^*, y^*

Dual

$$\begin{array}{ll} \text{maximize} & -b^T y - b_{m+1} y_{m+1} \\ \text{subject to} & A^T y + a_{m+1} y_{m+1} + c \geq 0 \end{array}$$

Solution $x^*, (y^*, 0)$ **optimal** for the new problem?

Adding new constraints

Optimality conditions

maximize $-b^T y - b_{m+1} y_{m+1}$
subject to $A^T y + a_{m+1} y_{m+1} + c \geq 0$ \longrightarrow Solution $(y^*, 0)$ is still **dual feasible**

Is x^* still **primal feasible**?

$$Ax = b$$

$$a_{m+1}^T x = b_{m+1}$$

$$x \geq 0$$

Yes

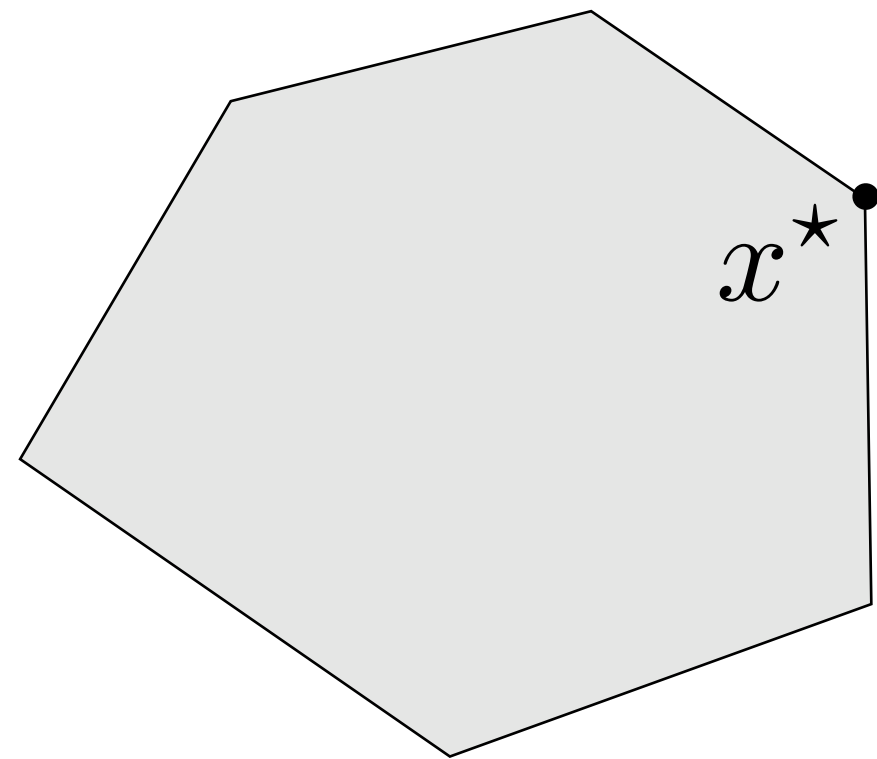
x^* still **optimal** for new problem

Otherwise

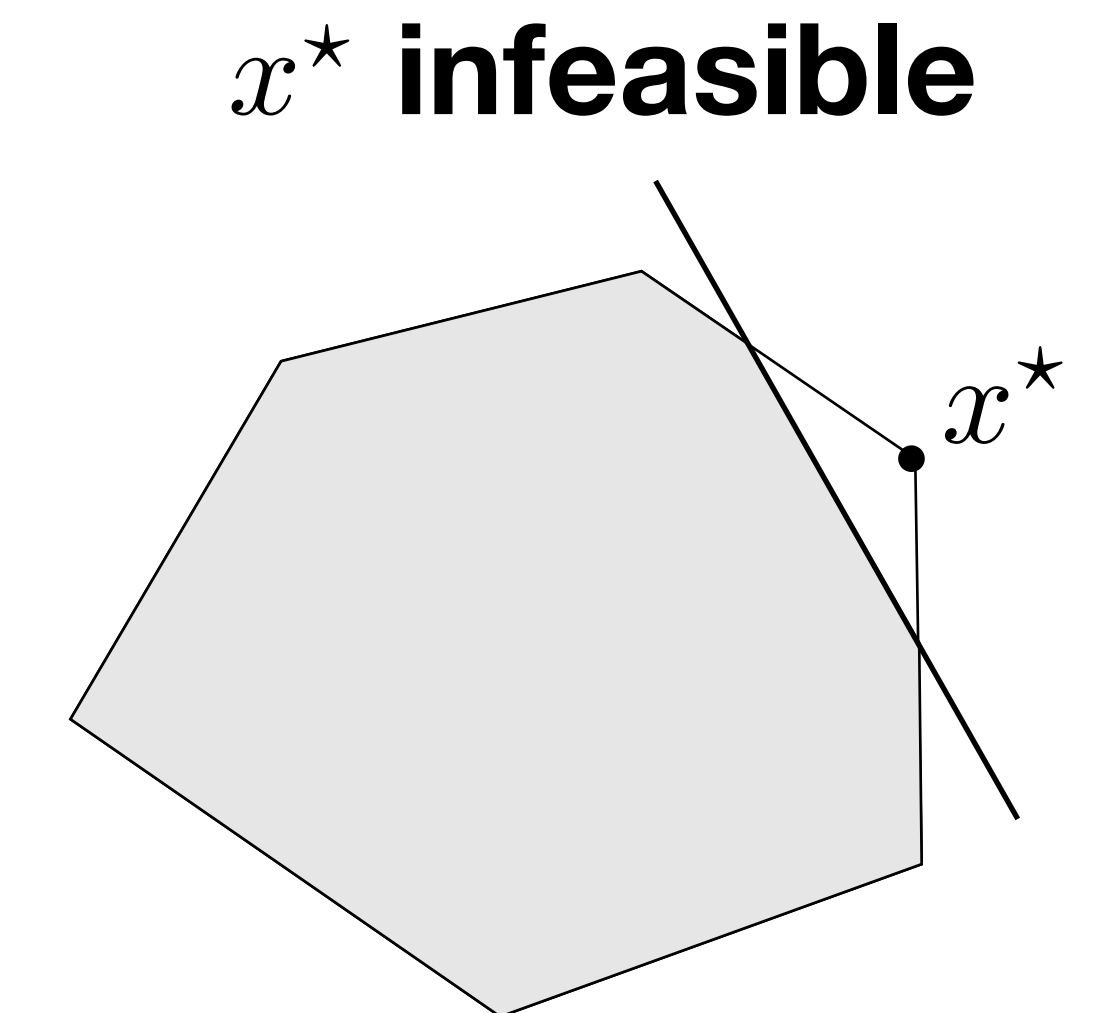
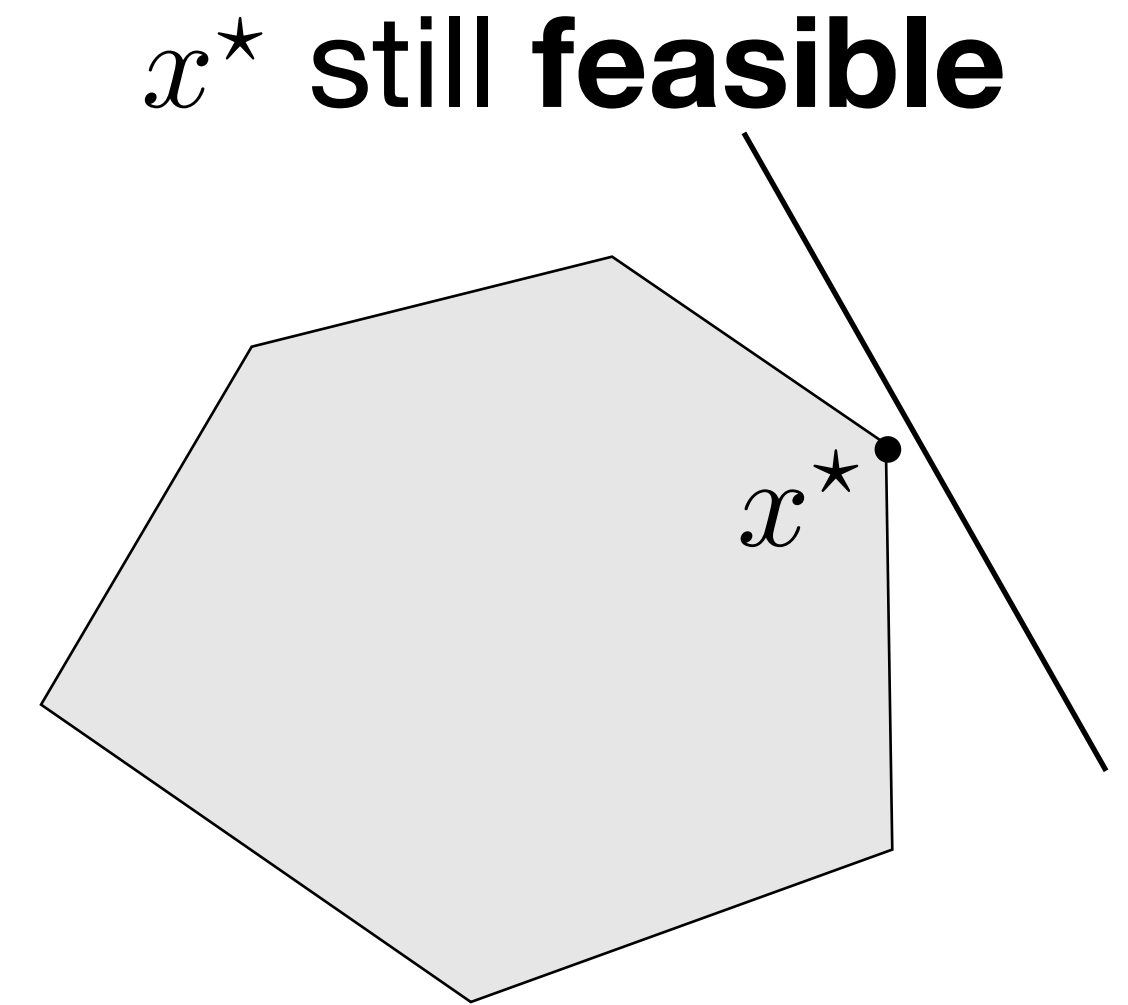
Dual simplex

Adding new constraints

Example



Add new constraint



Global sensitivity analysis

Information from primal-dual solution

Goal: extract information from x^*, y^* about their sensitivity with respect to changes in problem data

Modified LP

$$\begin{aligned} &\text{minimize} && c^T x \\ &\text{subject to} && Ax = b + u \\ & && x \geq 0 \end{aligned}$$

Optimal cost $p^*(u)$

Global sensitivity

Dual of modified LP

$$\begin{aligned} &\text{maximize} && -(b + u)^T y \\ &\text{subject to} && A^T y + c \geq 0 \end{aligned}$$

Global lower bound

Given y^* a dual optimal solution for $u = 0$, then

$$\begin{aligned} p^*(u) &\geq -(b + u)^T y^* && \text{(from weak duality and} \\ &= p^*(0) - u^T y^* && \text{dual feasibility)} \end{aligned}$$

It holds for any u

Global sensitivity

Example

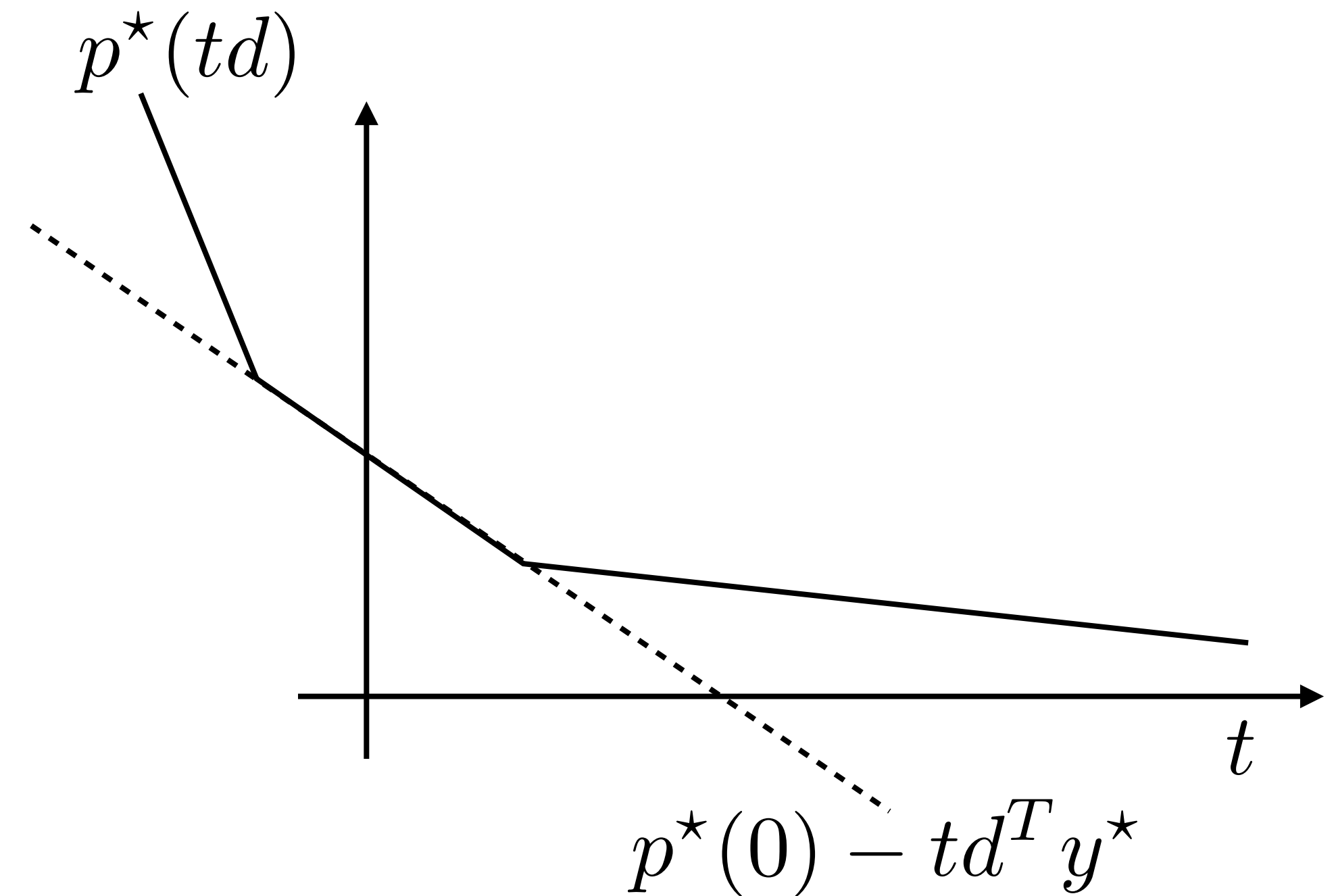
Take $u = td$ with $d \in \mathbf{R}^m$ fixed

$$\text{minimize } c^T x$$

$$\text{subject to } Ax = b + td$$

$$x \geq 0$$

$p^*(td)$ is the optimal value as a function of t



Sensitivity information (assuming $d^T y^* \geq 0$)

- $t < 0$ the optimal value increases
- $t > 0$ the optimal value decreases (not so much if t is small)

Optimal value function

$$p^*(u) = \min\{c^T x \mid Ax = b + u, x \geq 0\}$$

Assumption: $p^*(0)$ is finite

Properties

- $p^*(u) > -\infty$ everywhere (from global lower bound)
- the domain $\{u \mid p^*(u) < +\infty\}$ is a polyhedron
- $p^*(u)$ is piecewise-linear on its domain

Optimal value function is piecewise linear

Proof

Dual feasible set

$$p^*(u) = \min\{c^T x \mid Ax = b + u, x \geq 0\}$$

$$D = \{y \mid A^T y + c \geq 0\}$$

Assumption: $p^*(0)$ is finite

If $p^*(u)$ finite

$$p^*(u) = \max_{y \in D} -(b + u)^T y = \max_{k=1, \dots, r} -y_k^T u - b^T y_k$$

y_1, \dots, y_r are the extreme points of D

Local sensitivity analysis

Local sensitivity

u in neighborhood of the origin

Original LP

$$\begin{aligned} \text{minimize} \quad & c^T x \\ \text{subject to} \quad & Ax = b \\ & x \geq 0 \end{aligned}$$



Optimal solution

$$\begin{aligned} \text{Primal} \quad & x_i = 0, \quad i \notin B \\ & x_B^* = A_B^{-1} b \\ \text{Dual} \quad & y^* = -A_B^{-T} c_B \end{aligned}$$

Modified LP

$$\begin{aligned} \text{minimize} \quad & c^T x \\ \text{subject to} \quad & Ax = b + u \\ & x \geq 0 \end{aligned}$$

Modified dual

$$\begin{aligned} \text{maximize} \quad & -(b + u)^T y \\ \text{subject to} \quad & A^T y + c \geq 0 \end{aligned}$$

**Optimal basis
does not change**

Modified optimal solution

$$\begin{aligned} x_B^*(u) &= A_B^{-1} (b + u) = x_B^* + A_B^{-1} u \\ y^*(u) &= y^* \end{aligned}$$

Derivative of the optimal value function

Modified optimal solution

$$x_B^*(u) = A_B^{-1}(b + u) = x_B^* + A_B^{-1}u$$

$$y^*(u) = y^*$$

Optimal value function

$$p^*(u) = c^T x^*(u)$$

$$= c^T x^* + c_B^T A_B^{-1}u$$

$$= p^*(0) - y^{*T}u \quad (\text{affine for small } u)$$

Local derivative

$$\frac{\partial p^*(u)}{\partial u} = -y^*$$

$(y^*$ are the **shadow prices**)

Sensitivity example

minimize	$-60x_1 - 30x_2 - 20x_3$	-profit
subject to	$8x_1 + 6x_2 + x_3 \leq 48$	material
	$4x_1 + 2x_2 + 1.5x_3 \leq 20$	production
	$2x_1 + 1.5x_2 + 0.5x_3 \leq 8$	quality control
	$x \geq 0$	

$$x^* = (2, 0, 8, 24, 0, 0), \quad y^* = (0, 10, 10), \quad c^T x^* = -280, \quad \text{basis } \{1, 3, 4\}$$

What does $y_3^* = 10$ mean?

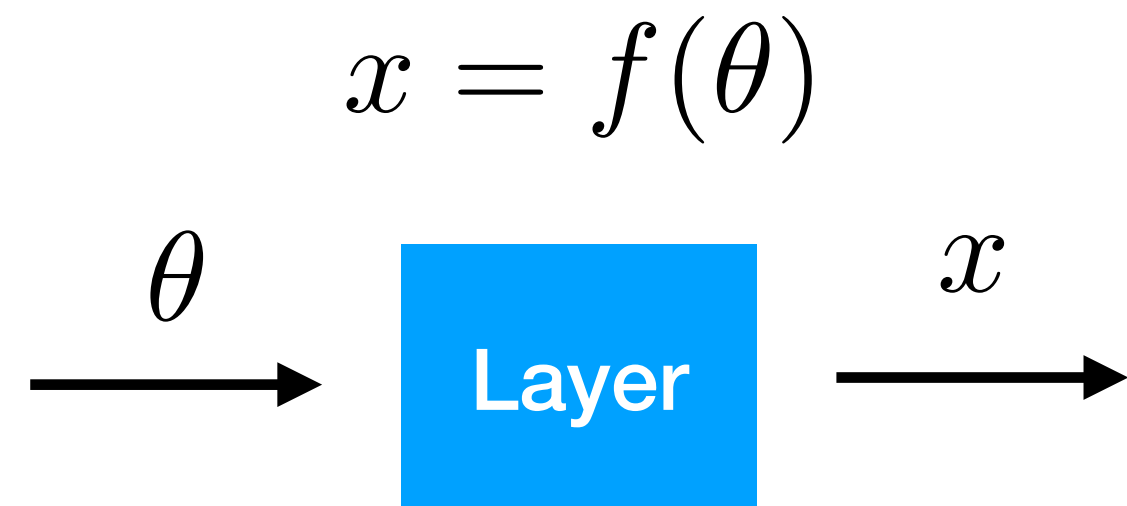
Let's increase the quality control budget by 1, i.e., $u = (0, 0, 1)$

$$p^*(10) = p^*(0) - y^{*T} u = -280 - 10 = -290$$

Differentiable optimization

Training a neural network

Single layer model



Training

minimize $\mathcal{L}(\theta)$

Gradient descent (more on this later)

$$\theta \leftarrow \theta - t \nabla_{\theta} \mathcal{L}(\theta)$$

Sensitivity

$$\nabla_{\theta} \mathcal{L} = \left(\frac{\partial \mathcal{L}}{\partial \theta} \right)^T = \left(\frac{\partial \mathcal{L}}{\partial x} \frac{\partial x}{\partial \theta} \right)^T = \left(\frac{\partial x}{\partial \theta} \right)^T \nabla_x \mathcal{L}$$

Can f be an **optimization problem**?

Implicit layers

<https://implicit-layers-tutorial.org/>

find $x(\theta)$
subject to $r(\theta, x(\theta)) = 0$ ($x(\theta)$ is implicitly defined by r)

How do we compute derivatives?

$$\frac{\partial x(\theta)}{\partial \theta}$$

Implicit function theorem

Under mild assumptions (non-singularity),

$$\frac{\partial r(\theta, x(\theta))}{\partial x} \frac{\partial x(\theta)}{\partial \theta} + \frac{\partial r(\theta, x(\theta))}{\partial \theta} = 0 \longrightarrow \frac{\partial x(\theta)}{\partial \theta} = - \left(\frac{\partial r(\theta, x(\theta))}{\partial x} \right)^{-1} \frac{\partial r(\theta, x(\theta))}{\partial \theta}$$

Optimization layers

$$x^*(\theta) = \underset{x}{\operatorname{argmin}} \quad c^T x$$

subject to $Ax \leq b$

Parameters: $\theta = \{c, A, b\}$

Solution $x^*(\theta)$

Features

- Add **domain knowledge** and **hard constraints**
- **End-to-end** training and optimization
- Nice theory and algorithms for general **convex optimization**
- **Applications** in RL, control, meta-learning, game theory, etc.

Goal

Compute $\frac{\partial x^*(\theta)}{\partial \theta}$

Optimality conditions

$$\begin{aligned} &\text{minimize} && c^T x \\ &\text{subject to} && Ax \leq b \end{aligned}$$

Parameters: $\theta = \{c, A, b\}$
Solution $x^*(\theta)$

Solve and obtain primal-dual pair x^*, y^* (forward-pass)

Optimality conditions

$$\begin{aligned} &A^T y + c = 0 \\ &\text{diag}(y)(Ax - b) = 0 \\ &y \geq 0, \quad b - Ax \geq 0 \end{aligned}$$

Mapping $r(\theta, x(\theta)) = 0$

Computing derivatives

Take differentials

$$\begin{array}{l} A^T y^* + c = 0 \\ \mathbf{diag}(y^*)(Ax - b) = 0 \end{array} \longrightarrow \begin{array}{l} dA^T y^* + A^T dy = 0 \\ \mathbf{diag}(Ax - b)dy + \mathbf{diag}(y^*)(dAx^* + Adx - db) + dc = 0 \end{array}$$

Linear system

$$\begin{bmatrix} 0 & A^T \\ \mathbf{diag}(y^*)A & \mathbf{diag}(Ax^* - b) \end{bmatrix} \begin{bmatrix} dx \\ dy \end{bmatrix} = - \begin{bmatrix} dA^T y^* + dc \\ \mathbf{diag}(y^*)(dAx^* - db) \end{bmatrix}$$

Example: How does x^* change with b_1 ?

Set $db = e_1$, $dA = 0$, $dc = 0$ and solve the linear system.

The solution dx will correspond to $\frac{\partial x}{\partial b_1}$

Is it always differentiable?

The linear system matrix must be invertible
(the problem must have unique solution)

$$\begin{bmatrix} 0 & A^T \\ \text{diag}(y^*)A & \text{diag}(Ax^* - b) \end{bmatrix} \begin{bmatrix} dx \\ dy \end{bmatrix} = - \begin{bmatrix} dA^T y^* + dc \\ \text{diag}(y^*)(dAx^* - db) \end{bmatrix}$$

M q

Remember. implicit function theorem

$$\frac{\partial x(\theta)}{\partial \theta} = - \left(\frac{\partial r(\theta, x(\theta))}{\partial x} \right)^{-1} \frac{\partial r(\theta, x(\theta))}{\partial \theta}$$

If not, **least squares** “subdifferential”

$$\text{minimize} \left\| M \begin{bmatrix} dx \\ dy \end{bmatrix} + q \right\|_2^2$$

Example

Learning to play Sudoku

			3
1			
		4	
4			1

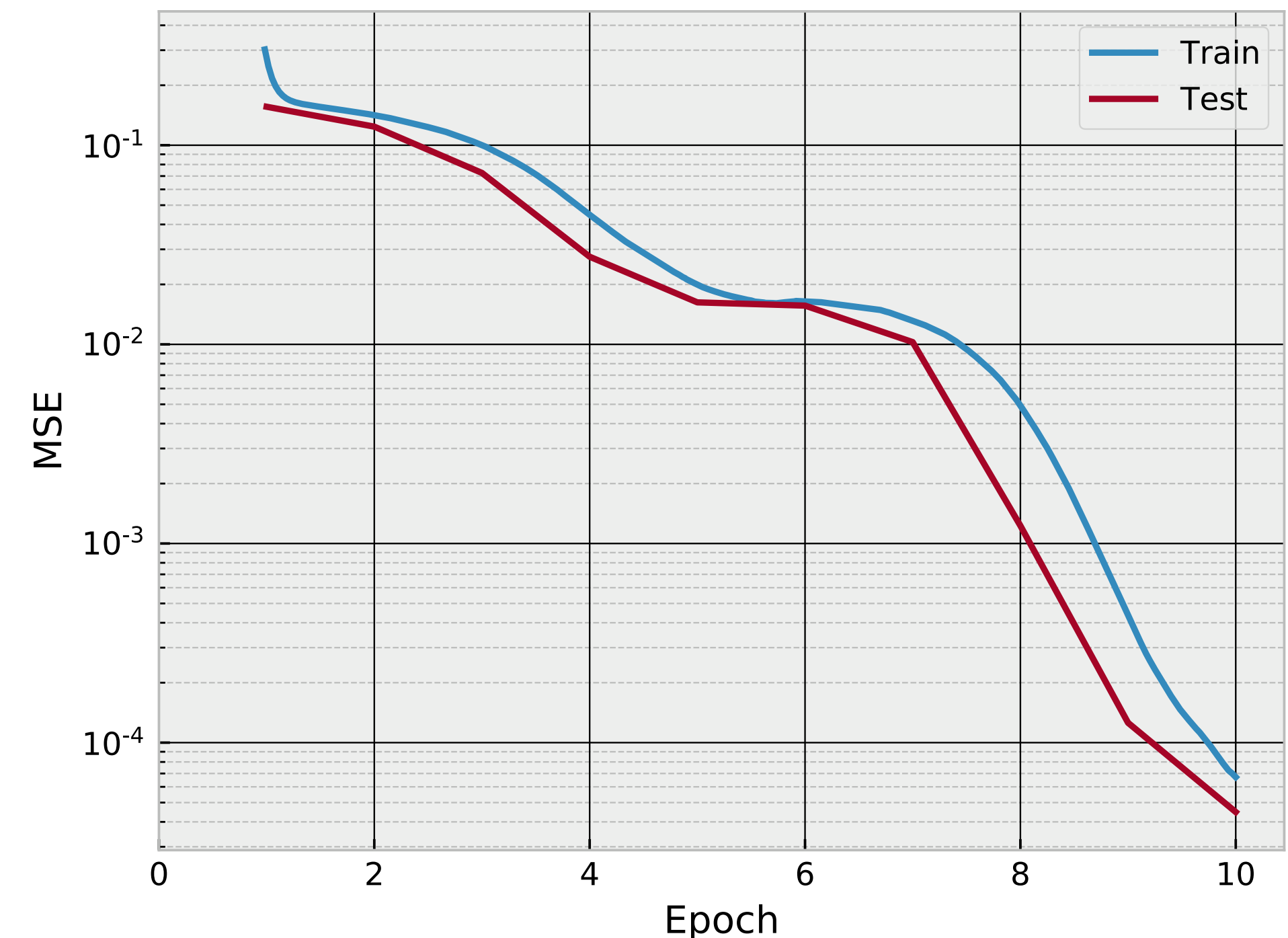
2	4	1	3
1	3	2	4
3	1	4	2
4	2	3	1

Sudoku constraint satisfaction problem

$$\begin{aligned} & \text{minimize} && 0 \\ & \text{subject to} && Ax = b \\ & && x \geq 0, x \in \mathbf{Z}^d \end{aligned}$$

Linear optimization layer (parameters $\theta = \{A, b\}$)

$$\begin{aligned} x^* = & \underset{x}{\operatorname{argmin}} && 0 \\ & \text{subject to} && Ax = b \\ & && x \geq 0 \end{aligned}$$



Sensitivity analysis in linear optimization

Today, we learned to:

- **Use** the most appropriate primal/dual simplex algorithm when variables and/or constraints are added
- **Analyze** sensitivity of the cost with respect to change in the data
- **Apply** sensitivity analysis to differentiable linear optimization layers

Next lecture

- Barrier methods for linear optimization