

# **ORF307 – Optimization**

## **2. Solving linear systems in practice**

# Ed Forum

- Please select lecture when adding notes to the forum

The screenshot shows the Ed Forum post creation interface. At the top, there are three tabs: 'Question', 'Post', and 'Announcement'. The 'Post' tab is selected. Below the tabs, there is a 'Title' field containing 'Lecture 2 Note'. Underneath the title, there are three category buttons: 'General', 'Questions', and 'Notes'. The 'Notes' button is selected. Below the categories, there is a 'Subcategory' dropdown menu. The dropdown menu is open, showing a list of options: '(None)', 'Lecture 1', 'Lecture 2', 'Lecture 3', 'Lecture 4', 'Lecture 5', 'Lecture 6', 'Lecture 7', 'Lecture 8', 'Lecture 9', and 'Lecture 10'. The 'Lecture 2' option is selected and highlighted in blue. Below the dropdown menu, there is a warning icon and the text 'Select a subcategory'. To the right of the dropdown menu, there are four checkboxes: 'Pinned', 'Anonymous', 'Anonymous Comments', and 'Megathread'. The 'Pinned' checkbox is checked, and the 'Anonymous' checkbox is checked. The 'Anonymous Comments' and 'Megathread' checkboxes are unchecked. Below the checkboxes, there is a 'Post' button with an upward arrow.

- Participation questions from today!

# Today's lecture

## Solving linear systems in practice

- Matrices: definition, operations, special cases
- Linear systems solutions
- Solving linear systems

# Matrices

# Matrices

matrix of size  $m \times n$

$$A = \begin{bmatrix} A_{11} & A_{12} & \dots & A_{1n} \\ A_{21} & A_{22} & \dots & A_{2n} \\ \vdots & \vdots & & \vdots \\ A_{m1} & A_{m2} & \dots & A_{mn} \end{bmatrix}$$

- $A_{ij}$  is the  $i, j$  element (also called *entry* or *coefficient*)
- $i$  is the row index,  $j$  the column index
- indices start at 1 (when you code, at 0)
- vectors are like matrices with 1 column

# Special matrices

$$I = \begin{bmatrix} 1 & & & \\ & 1 & & \\ & & \ddots & \\ & & & 1 \end{bmatrix}$$

## Special matrices

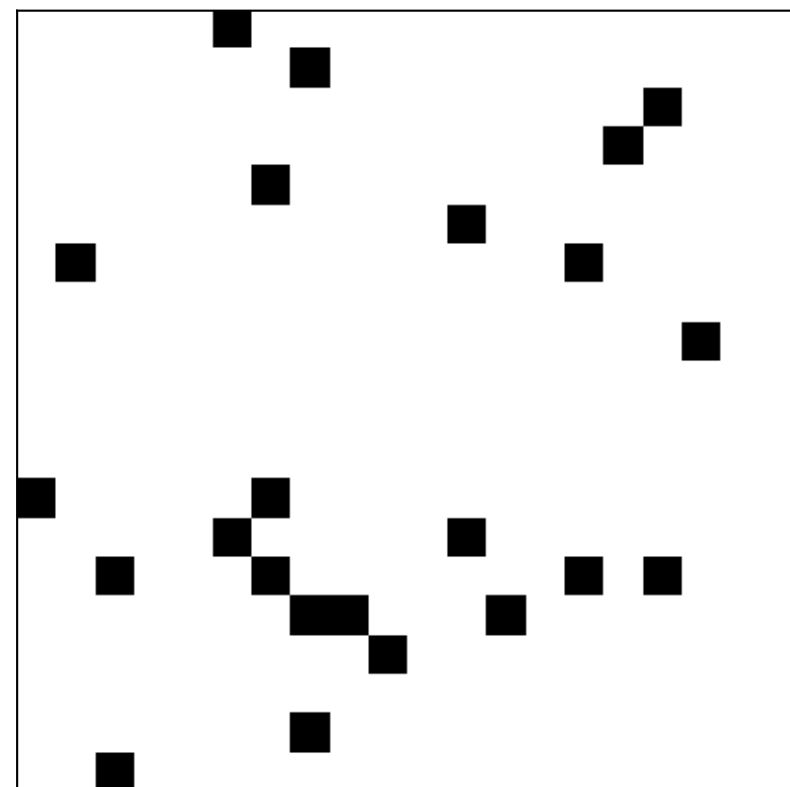
- $A = 0$  (zero matrix):  $A_{ij} = 0$ ,  $i = 1, \dots, m$ ,  $j = 1, \dots, n$
- $A = I$  (identity matrix):  $m = n$  with  $A_{ii} = 1$  and  $A_{ij} = 0$  for  $i \neq j$

# Special matrices

## Special matrices

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- $A = I$  (identity matrix):  $m = n$  with  $A_{ii} = 1$  and  $A_{ij} = 0$  for  $i \neq j$

## Sparse matrices (most entries are 0)



- Examples:  $0$  and  $I$
- Can be stored and manipulated efficiently
- $\text{nnz}(A)$  is the number of nonzero entries

# Diagonal and triangular matrices

## diagonal matrix

- A square matrix  $n \times n$  with  $A_{ij} = 0$  when  $i \neq j$
- $\text{diag}(a_1, \dots, a_n)$  is the diagonal matrix with  $A_{ii} = a_i$  for  $i = 1, \dots, n$

$$\text{diag}(0.2, -3) = \begin{bmatrix} 0.2 & 0 \\ 0 & -3 \end{bmatrix}$$



# Diagonal and triangular matrices

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## lower triangular matrix

$$A_{ij} = 0 \text{ for } i < j$$

$$\begin{bmatrix} -0.6 & 0 \\ 1.6 & -2 \end{bmatrix}$$

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## lower triangular matrix

$$A_{ij} = 0 \text{ for } i < j$$

$$\begin{bmatrix} -0.6 & 0 \\ 1.6 & -2 \end{bmatrix}$$

## upper triangular matrix

$$A_{ij} = 0 \text{ for } i > j$$

$$\begin{bmatrix} -0.2 & 0.3 \\ 0 & -1 \end{bmatrix}$$

# Block matrices

**Matrices whose entries are matrices**

$$n \times m \text{ matrix } A = \begin{bmatrix} B & C \\ D & E \end{bmatrix}$$

where  $B, C, D, E$  are  
*submatrices of blocks of  $A$*

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**column representation**

$$A = \begin{bmatrix} \downarrow & \downarrow & \dots & \downarrow \\ a_1 & a_2 & \dots & a_n \\ \uparrow & \uparrow & & \uparrow \end{bmatrix}$$

$(a_i \text{ are } m\text{-vectors})$

# Block matrices

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**column representation**

$$A = \begin{bmatrix} a_1 & a_2 & \dots & a_n \end{bmatrix}$$

( $a_i$  are  $m$ -vectors)

**row representation**

$$A = \begin{bmatrix} \overline{b_1^T} \\ \overline{b_2^T} \\ \vdots \\ \overline{b_n^T} \end{bmatrix}$$

( $b_i$  are  $n$ -row-vectors)

# Matrix operations

## transpose

A *transpose* of a matrix  $A$  is denoted as  $A^T$  where

$$(A^T)_{ij} = A_{ji}, \quad i = 1, \dots, m, \quad j = 1, \dots, n$$

$$A = \begin{bmatrix} 0.2 & \Theta \\ 1 & -0.3 \end{bmatrix} \rightarrow A^T = \begin{bmatrix} 0.2 & 1 \\ \Theta & -0.3 \end{bmatrix}$$

$$e = \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix} \quad e^T = [1 \ 2 \ 3]$$

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$$(A^T)_{ij} = A_{ji}, \quad i = 1, \dots, m, \quad j = 1, \dots, n$$

## addition (just like vectors)

$$(A + B)_{ij} = A_{ij} + B_{ij}, \quad i = 1, \dots, m, \quad j = 1, \dots, n$$

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## scalar multiplication (just like vectors)

$$(\alpha A)_{ij} = \alpha A_{ij}, \quad i = 1, \dots, m, \quad j = 1, \dots, n$$



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$$(\alpha A)_{ij} = \alpha A_{ij}, \quad i = 1, \dots, m, \quad j = 1, \dots, n$$

## Many properties

$$(A^T)^T = A, \quad A + B = B + A, \quad \alpha(A + B) = \alpha A + \alpha B$$

# Matrix-vector multiplication

## dot product

A matrix-vector product of an  $m \times n$  matrix  $A$  and a  $n$ -vector  $x$  is denoted as

$$y = Ax, \quad \text{where } y_i = A_{i1}x_1 + \cdots + A_{in}x_n, \quad i = 1, \dots, m$$

$$\begin{bmatrix} 0 & 2 & -1 \\ -2 & 1 & 1 \end{bmatrix} \begin{bmatrix} 2 \\ 1 \\ -1 \end{bmatrix} = \begin{bmatrix} 3 \\ -4 \end{bmatrix}$$

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## row interpretation

$$y_i = b_i^T x$$

where  $b_1^T, \dots, b_m^T$  are rows of  $A$

**example A1**

# Matrix-vector multiplication

$$e_j = \begin{bmatrix} 0 \\ \vdots \\ 1 \\ \vdots \\ 0 \end{bmatrix} \quad \text{jTH ELEMENT}$$

## dot product

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## row interpretation

$$y_i = b_i^T x$$

where  $b_1^T, \dots, b_m^T$  are rows of  $A$

**example**  $A1$

$$A = \begin{bmatrix} b_1^T \\ b_2^T \\ \vdots \\ b_m^T \end{bmatrix}$$

## column interpretation

$$y = x_1 a_1 + \dots + x_n a_n$$

where  $a_1, \dots, a_n$  are the columns of  $A$

**example**  $Ae_j = a_j$

$$A = [a_1 \ a_2 \ \dots \ a_n]$$

# Return matrix — portfolio vector

$R$  is the  $T \times n$  matrix of **asset returns**

	AAPL	GOOG	MMM	AMZN	
$R =$	0.00219	0.0006	-0.00113	0.00202	Mar 1, 2016
	0.00744	-0.00894	-0.00019	-0.00468	Mar 2, 2016
	0.01488	-0.00215	0.00433	-0.00407	Mar 3, 2016

$$R_{ti} = \frac{p_{ti}^{\text{final}} - p_{ti}^{\text{initial}}}{p_{ti}^{\text{initial}}}$$

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constant investment  $w$



$Rw$  is the vector of portfolio returns over periods  $1, \dots, T$

$$R_{ti} = \frac{p_{ti}^{\text{final}} - p_{ti}^{\text{initial}}}{p_{ti}^{\text{initial}}}$$

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$$R_{ti} = \frac{p_{ti}^{\text{final}} - p_{ti}^{\text{initial}}}{p_{ti}^{\text{initial}}}$$

**example**  $w = (0.4, 0.3, -0.2, 0.5)$

$Rw = (0.00213, -0.00201, 0.00241)$  11

# Symmetric positive semidefinite matrices

**symmetric matrix**

$$A^T = A$$

**positive semidefinite matrix**

$$x^T Ax \geq 0 \text{ for any } x \in \mathbf{R}^n$$

**positive definite matrix**

$$x^T Ax > 0 \text{ for any } x \neq 0$$



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**example**

$$A = \begin{bmatrix} 2 & 6 \\ 6 & 20 \end{bmatrix}$$

# Symmetric positive semidefinite matrices

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$$x^T Ax > 0 \text{ for any } x \neq 0$$

**example**

$$A = \begin{bmatrix} 2 & 6 \\ 6 & 20 \end{bmatrix}$$

$$\begin{bmatrix} x_1 & x_2 \end{bmatrix} \begin{bmatrix} 2 & 6 \\ 6 & 20 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} x_1 & x_2 \end{bmatrix} \begin{bmatrix} 2x_1 + 6x_2 \\ 6x_1 + 20x_2 \end{bmatrix}$$

$$= 2x_1^2 + 12x_1x_2 + 20x_2^2$$

$$= 2(x_1 + 3x_2)^2 + 2x_2^2$$

**sum of squares**

# Matrix multiplication

## matrix product

A matrix product of an  $m \times p$  matrix  $A$  and a  $p \times n$  matrix  $B$  is

$$C = AB, \quad \text{where } C_{ij} = A_{i1}B_{1j} + \cdots + A_{ip}B_{pj}, \quad i = 1, \dots, m, j = 1, \dots, n$$

# Matrix multiplication

## matrix product

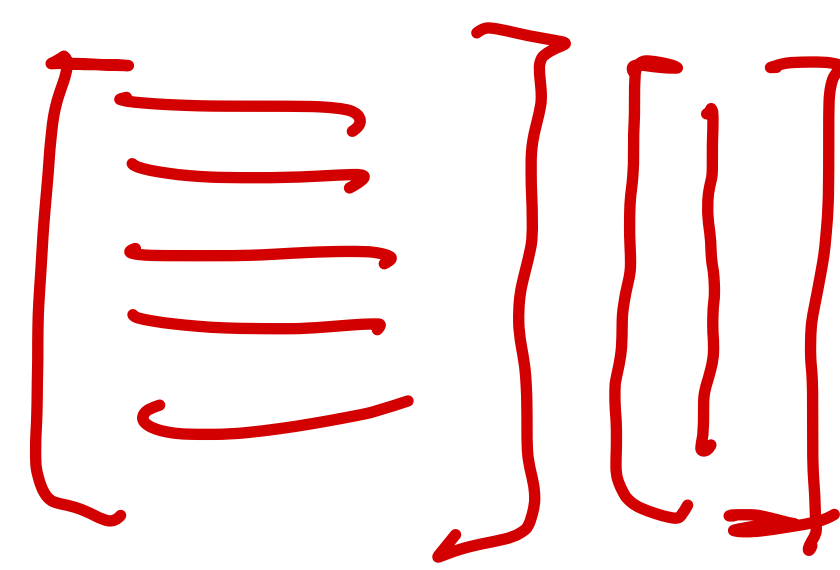
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(move along  $i$ th row of  $A$  and  $j$ th column of  $B$ )

$$\begin{bmatrix} -1.5 & 3 & 2 \\ 1 & -1 & 0 \end{bmatrix} \begin{bmatrix} -1 & -1 \\ 0 & -2 \\ 1 & 0 \end{bmatrix} = \begin{bmatrix} 3.5 & -4.5 \\ -1 & 1 \end{bmatrix}$$

# Complexity



VECTOR-VECTOR  
MULTIPLY

Given  $m \times n$  matrix  $A$

- Matrix addition, scalar-matrix multiplication:  $mn$  flops
- Matrix-vector multiplication:  $m(2n - 1) \approx 2mn$  flops
- Matrix-matrix multiplication ( $A \in \mathbf{R}^{m \times p}$ ,  $B \in \mathbf{R}^{p \times n}$ ):  $(mn)(2p - 1) \approx 2mnp$  flops  
(inner product of  $p$  vectors)

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## Questions

- How many flops does it take to multiply two  $1000 \times 1000$  matrices?
- How long does it take on a computer?

# Linear systems solutions

# Linear systems of equations

Given an  $m \times n$  matrix  $A$  and a  $m$ -vector  $b$ , find a  $n$ -vectors  $x$  such that

$$Ax = b$$



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## typical scenarios

underdetermined  
(wide)

$$m < n$$

$$\begin{bmatrix} * & * & * \\ * & * & * \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} * \\ * \end{bmatrix}$$

infinite  
solutions

# Linear systems of equations

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infinite  
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square

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unique  
solution

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unique  
solution

overdetermined  
(tall)

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no  
solution

# Linear systems of equations

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$$Ax = b$$

## typical scenarios

underdetermined  
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$$m < n$$

$$\begin{bmatrix} * & * & * \\ * & * & * \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} * \\ * \end{bmatrix}$$

infinite  
solutions

square

$$m = n$$

$$\begin{bmatrix} * & * \\ * & * \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} * \\ * \end{bmatrix}$$

unique  
solution

**most common**

overdetermined  
(tall)

$$m > n$$

$$\begin{bmatrix} * & * \\ * & * \\ * & * \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} * \\ * \\ * \end{bmatrix}$$

no  
solution

# Solving square linear systems

Given an  $n \times n$  matrix  $A$  and a  $n$ -vector  $b$ , find a  $n$ -vector  $x$  such that

$$Ax = b$$

# Solving square linear systems

Given an  $n \times n$  matrix  $A$  and a  $n$ -vector  $b$ , find a  $n$ -vector  $x$  such that

$$Ax = b \longrightarrow A^{-1}Ax = A^{-1}b$$

# Solving square linear systems

Given an  $n \times n$  matrix  $A$  and a  $n$ -vector  $b$ , find a  $n$ -vector  $x$  such that

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↑  
inverse

# Solving square linear systems

Given an  $n \times n$  matrix  $A$  and a  $n$ -vector  $b$ , find a  $n$ -vector  $x$  such that

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**When does it work?**

↑  
inverse



# Solving square linear systems

Given an  $n \times n$  matrix  $A$  and a  $n$ -vector  $b$ , find a  $n$ -vector  $x$  such that

$$Ax = b \longrightarrow A^{-1}Ax = A^{-1}b \longrightarrow x = A^{-1}b$$

**When does it work?**

↓

$A$  must be invertible

↑  
inverse

- Columns of  $A$  are linearly independent
- Rows of  $A$  are linearly independent
- Columns/rows form a **basis** of  $\mathbf{R}^n$

# Solving linear systems

# How do we solve linear systems in practice?

$$Ax = b$$

## Idea

- compute  $A^{-1}$
- multiply  $A^{-1}b$

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## Idea

- compute  $A^{-1}$
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## Example

5000  $\times$  5000 matrix  $A$  and a 5000-vector  $b$

- Solve by computing  $A^{-1}$
- Solve with `numpy.linalg.solve`

# How do we solve linear systems in practice?

$$Ax = b$$

## Idea

- compute  $A^{-1}$
- multiply  $A^{-1}b$

## Example

5000  $\times$  5000 matrix  $A$  and a 5000-vector  $b$

- Solve by computing  $A^{-1}$
- Solve with `numpy.linalg.solve`

**What's happening inside?**

# Easy linear systems

## Diagonal matrix

$$\begin{bmatrix} A_{11} & & & \\ & \ddots & & \\ & & A_{nn} & \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix}$$



$$\begin{aligned} A_{11}x_1 &= b_1 \\ A_{22}x_2 &= b_2 \\ &\vdots \\ A_{nn}x_n &= b_n \end{aligned}$$

# Easy linear systems

## Diagonal matrix

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## Solution

$$x = A^{-1}b = (b_1/A_{11}, \dots, b_n/A_{nn})$$

# Easy linear systems

## Diagonal matrix

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## Solution

$$x = A^{-1}b = (b_1/A_{11}, \dots, b_n/A_{nn})$$

## Complexity

$n$  flops



# Easy linear systems

## Lower triangular matrix

$$\begin{bmatrix} A_{11} & & & \\ A_{21} & A_{22} & & \\ \vdots & & \ddots & \\ A_{n1} & A_{n2} & \dots & A_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix}$$



$$\begin{aligned} A_{11}x_1 &= b_1 \\ A_{21}x_1 + A_{22}x_2 &= b_2 \\ &\vdots \\ A_{n1}x_1 + A_{n2}x_2 + \dots + A_{nn}x_n &= b_n \end{aligned}$$

# Easy linear systems

## Lower triangular matrix

$$\begin{bmatrix} A_{11} & & & \\ A_{21} & A_{22} & & \\ \vdots & & \ddots & \\ A_{n1} & A_{n2} & \dots & A_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix}$$



$$\begin{aligned} & A_{11}x_1 = b_1 \\ & \underbrace{A_{21}x_1} + \underbrace{A_{22}x_2} = b_2 \\ & \vdots \\ & A_{n1}x_1 + A_{n2}x_2 + \dots + A_{nn}x_n = b_n \end{aligned}$$

### Solution: “forward substitution”

- First equation:  $x_1 = b_1/A_{11}$
- Second equation:  $x_2 = (b_2 - A_{21}x_1)/A_{22}$
- Repeat to get  $x_3, \dots, x_n$

# Easy linear systems

## Lower triangular matrix

$$\begin{bmatrix} A_{11} & & & \\ A_{21} & A_{22} & & \\ \vdots & & \ddots & \\ A_{n1} & A_{n2} & \dots & A_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix}$$



$$\sum_{i=1}^n 2i - 1 = 2 \left( \sum_{i=1}^n i \right) - n = 2 \left( \frac{n(n+1)}{2} \right) - n = n^2 + n - n$$

$$\begin{aligned} A_{11}x_1 &= b_1 \\ A_{21}x_1 + A_{22}x_2 &= b_2 \\ &\vdots \\ A_{n1}x_1 + A_{n2}x_2 + \dots + A_{nn}x_n &= b_n \end{aligned}$$

### Solution: “forward substitution”

- First equation:  $x_1 = b_1/A_{11}$
- Second equation:  $x_2 = (b_2 - A_{21}x_1)/A_{22}$
- Repeat to get  $x_3, \dots, x_n$

### Complexity

- First equation: 1 flop (division)
- Second equation: 3 flops
- $i$ th step needs  $2i - 1$  flops

$$1 + 3 + \dots + (2n - 1) = n^2 \text{ flops}$$

# Easy linear systems

## Upper triangular matrix

$$\begin{bmatrix} A_{11} & \dots & A_{n-1,n} & A_{1n} \\ & \ddots & & \vdots \\ & & A_{n-1,n-1} & A_{n-1,n} \\ & & & A_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix} \longrightarrow \begin{array}{l} A_{11}x_1 + \dots + A_{1,n-1}x_{n-1} + A_{1n}x_n = b_1 \\ \vdots \\ A_{n-1,n-1}x_{n-1} + A_{n-1,n}x_n = b_{n-1} \\ A_{nn}x_n = b_n \end{array}$$

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**Solution:** “backward substitution”

- Last equation:  $x_n = b_n/A_{nn}$
- Second to last equation:  
 $x_{n-1} = (b_{n-1} - A_{n-1,n}x_n)/A_{n-1,n-1}$
- Repeat to get  $x_{n-2}, \dots, x_1$

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## Permutation matrices

$\pi = (\pi_1, \dots, \pi_n)$  is a permutation of  $(1, 2, \dots, n)$

A  $n \times n$  permutation matrix  $P$ ,  
permutes the vector  $x$

$$Px = (x_{\pi_1}, \dots, x_{\pi_n})$$

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**example**

$$\pi = (2, 3, 1)$$



$$P \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} x_2 \\ x_3 \\ x_1 \end{bmatrix}$$



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- $P_{ij} = \begin{cases} 1 & j = \pi_i \\ 0 & \text{otherwise} \end{cases}$
- $P^{-1} = P^T$  (inverse permutation)

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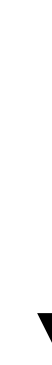
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- $P^{-1} = P^T$  (inverse permutation)

## Complexity

Solve  $Px = b$ : 0 flops (no operations)

## example

$$\pi = (2, 3, 1)$$

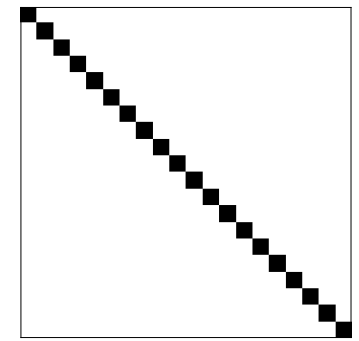


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# Summary of easy linear systems



**diagonal**

$$A = \text{diag}(a_1, \dots, a_n)$$

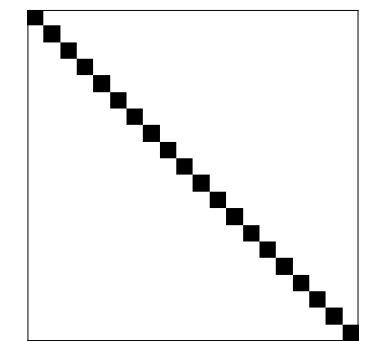
**method**

$$x_i = b_i / a_i$$

**flops**

$n$

# Summary of easy linear systems



**diagonal**

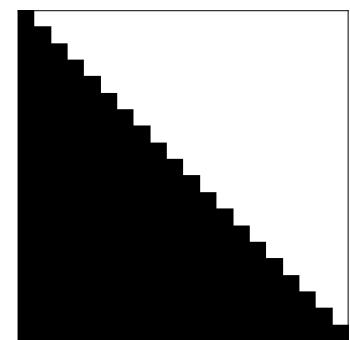
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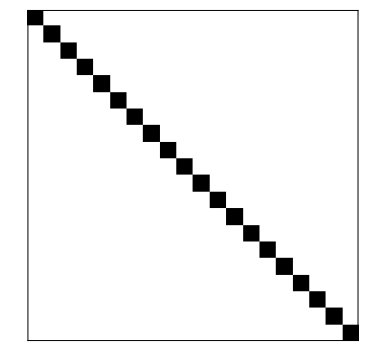
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$$A_{ij} = 0 \text{ for } i < j$$

**forward  
substitution**

$$n^2$$

# Summary of easy linear systems



**diagonal**

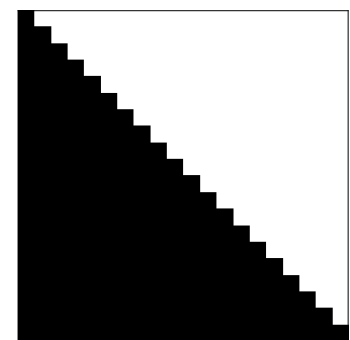
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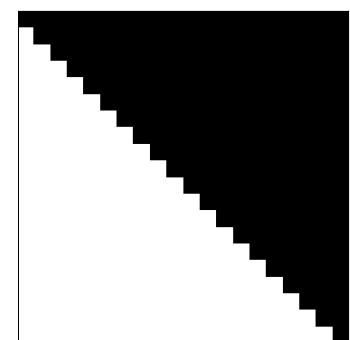


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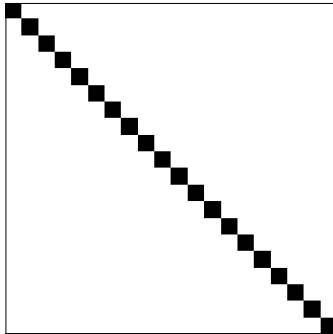
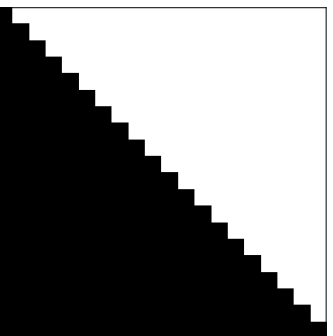
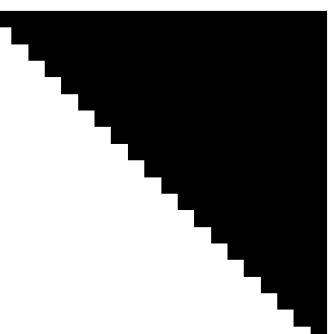
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# Summary of easy linear systems

		method	flops
	<b>diagonal</b> $A = \text{diag}(a_1, \dots, a_n)$	$x_i = b_i/a_i$	$n$
	<b>lower triangular</b> $A_{ij} = 0$ for $i < j$	forward substitution	$n^2$
	<b>upper triangular</b> $A_{ij} = 0$ for $i > j$	backward substitution	$n^2$
	<b>permutation</b> $P_{ij} = 1$ if $j = \pi_i$ else 0	inverse permutation	0

# How do we solve linear systems in practice?

$$Ax = b$$



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**Any idea?**

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**Any idea?**

We know how to solve special ones

Let's use that!

# The factor-solve method for solving $Ax = b$

1. **Factor**  $A$  as a product of simple matrices:

$$A = A_1 A_2 \cdots A_k, \quad \longrightarrow \quad A_1 A_2, \dots, A_k x = b$$

( $A_i$  diagonal, upper/lower triangular, permutation, etc)

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$$A_2 A_3 \dots A_k x \rightarrow x_1$$

$x_2$

2. **Compute**  $x = A^{-1}b = A_k^{-1} \cdots A_1^{-1}b$   
by solving  $k$  “easy” systems

$$\rightarrow A_1 x_1 = b$$

$$\rightarrow A_2 x_2 = x_1$$

$\vdots$

$$A_k x = x_{k-1}$$

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by solving  $k$  “easy” systems  $\longrightarrow$

$$A_1 x_1 = b$$

$$A_2 x_2 = x_1$$

$$\vdots$$

$$A_k x = x_{k-1}$$

**Note:** step 2 is much cheaper than step 1

# Multiple right-hand sides

You now have factored  $A$  and you want to solve  $d$  linear systems with different right-hand side  $m$ -vectors  $b_i$

$$Ax = b_1 \quad Ax = b_2 \quad \dots \quad Ax = b_d$$

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1. Factor  $A = A_1, \dots, A_k$  **only once** (expensive)
2. Solve all linear systems using **the same factorization** (cheap)

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**Solve many “at the price of one”**



# *LU* Factorization

Every invertible matrix  $A$  can be factored as

$$A = PLU \longrightarrow P^T A = LU$$

$P$  permutation,  $L$  lower triangular,  $U$  upper triangular

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## Procedure

- Similar to Gaussian elimination (but we can reuse  $P$ ,  $L$ , and  $U$ !)
- Permutation  $P$  avoids divisions by 0
- One of infinite possible combinations of  $P, L, U$

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## Complexity

- $(2/3)n^3$  flops
- Less if  $A$  has special structure (sparse, diagonal, etc)

# *LU* Solution

$$Ax = b, \quad \Rightarrow \quad PLUx = b$$

## Iterations

1. *Permutation*: Solve  $Px_1 = b$  (0 flops)
2. *Forward substitution*: Solve  $Lx_2 = x_1$  ( $n^2$  flops)
3. *Backward substitution*: Solve  $Ux = x_2$  ( $n^2$  flops)

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## Complexity

- Factor + solve:  $(2/3)n^3 + 2n^2 \approx (2/3)n^3$  (for large  $n$ )
- Just solve (prefactored):  $2n^2$

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Every positive definite matrix  $A$  can be factored as

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- Works only on **symmetric with positive definite** matrices
- No need to permute as in  $LU$
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- $(1/3)n^3$  flops (half of  $LU$  decomposition)
- Less if  $A$  has special structure (sparse, diagonal, etc)



# $LL^T$ (Cholesky) Solution

$$Ax = b, \quad \Rightarrow \quad LL^T x = b$$

## Iterations

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2. *Backward substitution*: Solve  $L^T x = x_1$  ( $n^2$  flops)

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# What complexity really means?

## Example

Large matrix  $n \times n$  matrix  $A$  with  $n = 10,000$

$$\begin{array}{l} \text{Factor + solve: } (2/3)n^3 \\ \text{Just solve: } 2n^2 \end{array} \longrightarrow \text{Gains: } \frac{(2/3)n^3}{2n^2} = (1/3)n \approx 3,333 \text{ times}$$

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**3 thousand times!**

Something that takes 1 second  $\longrightarrow$   $\approx$  1 hour

# Linear system example

## Polynomial interpolation

Given a cubic polynomial

$$p(x) = c_1 + c_2x + c_3x^2 + c_4x^3$$

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$$p(0.1) = b_3$$

$$p(0.8) = b_4$$

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Equivalent linear system

$$Ac = b$$

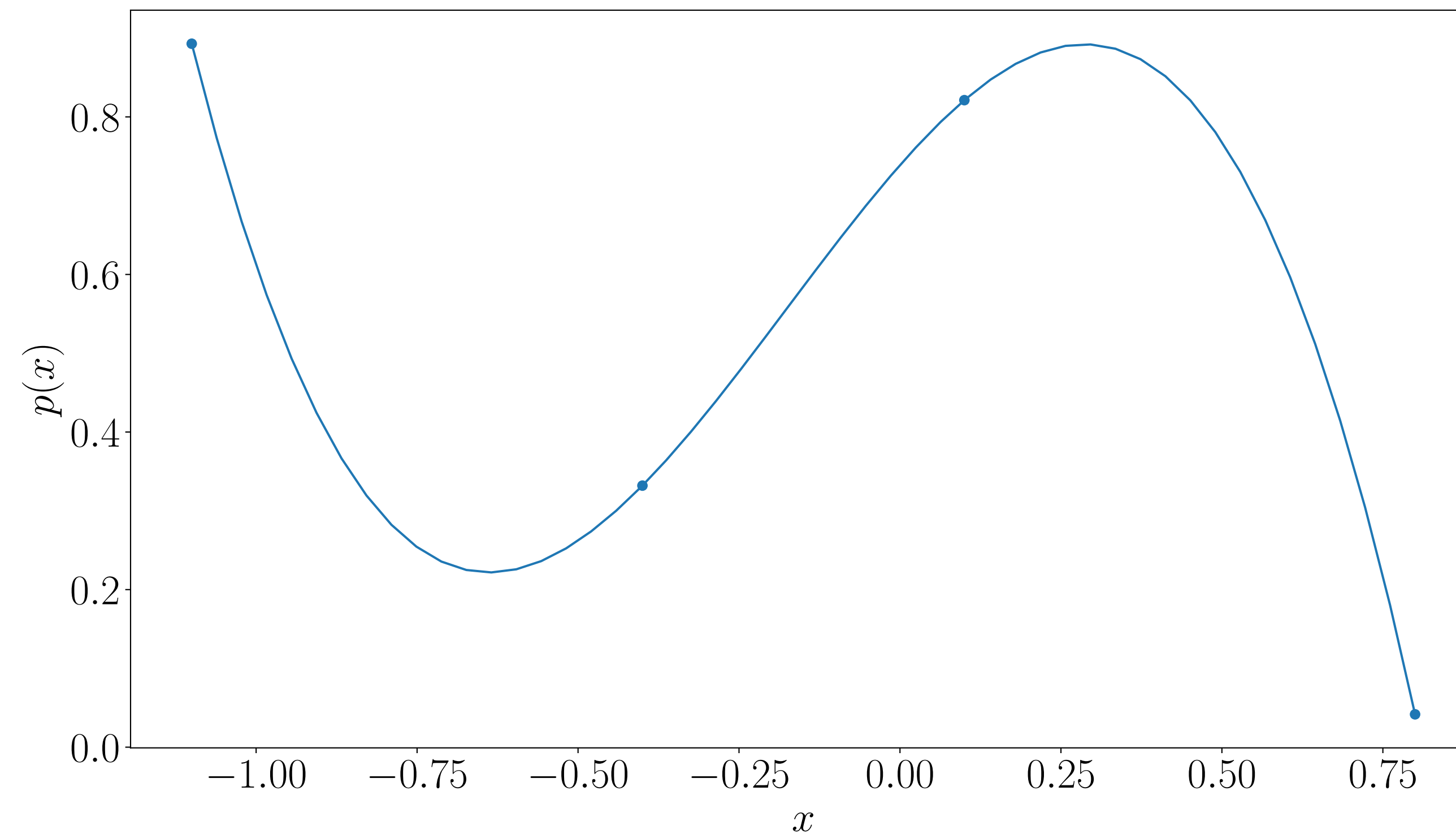
$$\begin{bmatrix} 1 & -1.1 & (-1.1)^2 & (-1.1)^3 \\ 1 & -0.4 & (-0.4)^2 & (-0.4)^3 \\ 1 & 0.1 & (0.1)^2 & (0.1)^3 \\ 1 & 0.8 & (0.8)^2 & (0.8)^3 \end{bmatrix} \begin{bmatrix} c_1 \\ c_2 \\ c_3 \\ c_4 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ b_4 \end{bmatrix}$$

# Polynomial interpolation

## Plot

$$p(x) = c_1 + c_2x + c_3x^2 + c_4x^3$$

$$c = (0.74, 0.93, -0.89, -1.70)$$





# Solving linear systems in practice

Today, we learned to:

- **Avoid** computing inverses
- **Solve** linear systems using the factor-solve method
- **Understand** the complexity of solving linear systems (useful to build optimization algorithms!)

# References

- S. Boyd, L. Vandenberghe: Introduction to Applied Linear Algebra – Vectors, Matrices, and Least Squares
  - Chapter 6: matrices
  - Chapter 10: matrix multiplication
  - Chapter 11: matrix inverses/solving linear systems

# Next lecture

- Solve optimization problems: least squares