

ORF522 – Linear and Nonlinear Optimization

13. Subgradient method and proximal operators

Today's lecture

[Chapter 3 and 6, FMO] [PA]

Subgradient method

Proximal methods

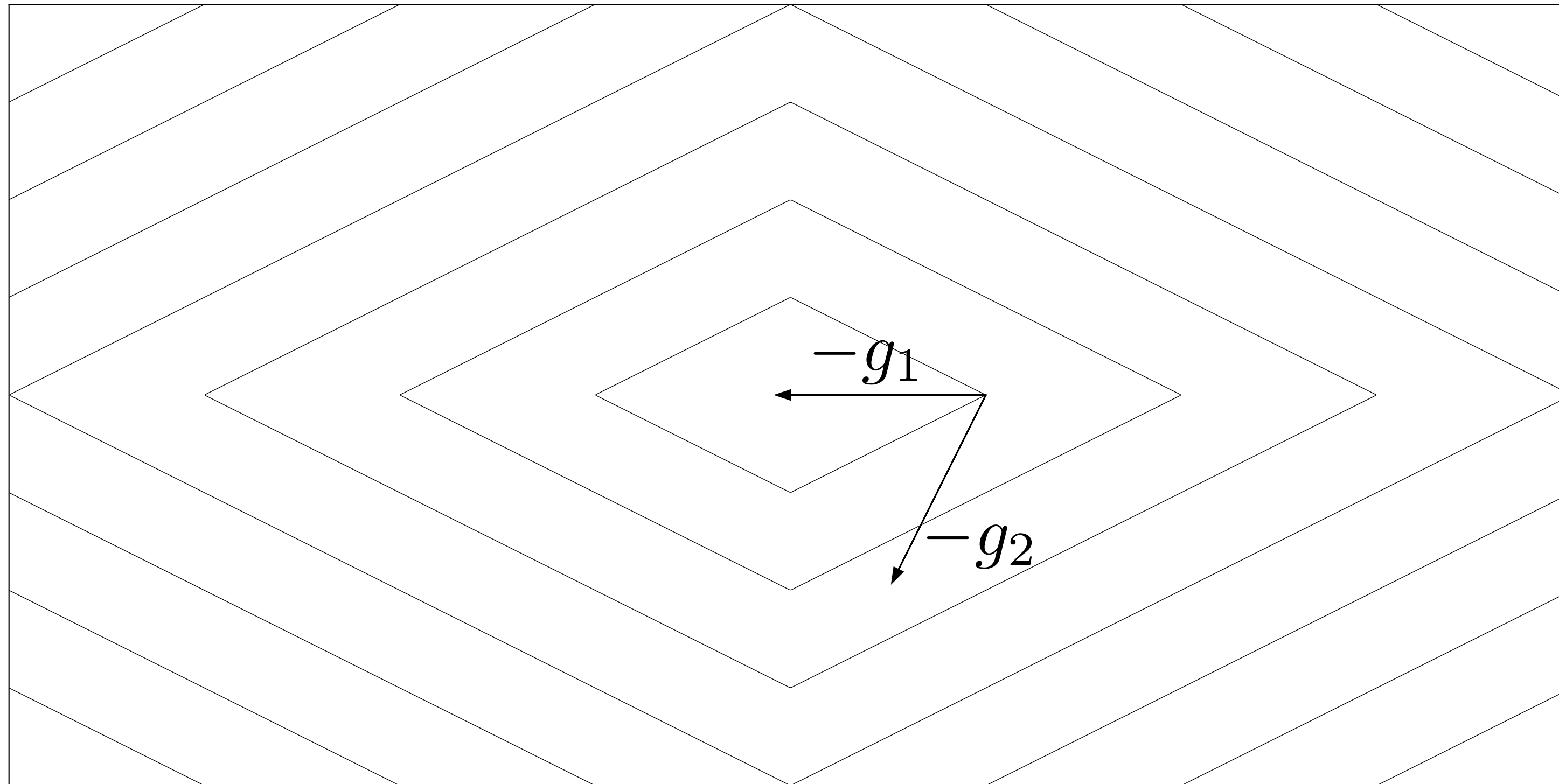
- Proximal operators
- Proximal gradient method

Subgradient method

Negative subgradients are not necessarily descent directions

$$f(x) = |x_1| + 2|x_2|$$

$$x = (1, 0)$$



$g_1 = (1, 0) \in \partial f(x)$ and
 $-g_1$ is a descent direction

$g_2 = (1, 2) \in \partial f(x)$ and
 $-g_2$ is not a descent direction

Subgradient method

Convex optimization problem

$$\text{minimize } f(x) \quad (\text{optimal cost } f^*)$$

Iterations

$$x^{k+1} = x^k - t_k g^k, \quad g^k \in \partial f(x^k)$$

g^k is **any subgradient** of f at x^k

Not a descent method, keep track of the best point

$$f_{\text{best}}^k = \min_{i=1, \dots, k} f(x^i)$$

Step sizes

Line search can lead to **suboptimal points**

Step sizes *pre-specified*, not adaptively computed
(different than gradient descent)

Fixed: $t_k = t$ for $k = 0, \dots$

Diminishing: $t_k \rightarrow 0$, $\sum_{k=0}^{\infty} t_k = \infty$

goes to 0 but not too fast
e.g., $t_k = O(1/k)$

*(include sequences that are
square summable but not summable)*

Convergence

Assumptions

- f is convex with $\text{dom } f = \mathbf{R}^n$
- $f(x^*) > -\infty$ (finite optimal value)
- f is Lipschitz continuous with constant $G > 0$, i.e.

$$|f(x) - f(y)| \leq G\|x - y\|_2, \quad \forall x, y$$

which is equivalent to $\|g\|_2 \leq G, \quad \forall g \in \partial f(x), \forall x$

Convergence

Lipschitz continuity equivalence

f is Lipschitz continuous with constant $G > 0$, i.e.

$$|f(x) - f(y)| \leq G\|x - y\|_2, \quad \forall x, y$$

which is equivalent to $\|g\|_2 \leq G, \quad \forall g \in \partial f(x), \forall x$

Proof

If $\|g\| \leq G$ for all subgradients, pick $x, g_x \in \partial f(x)$ and $y, g_y \in \partial f(y)$. Then,

$$\begin{aligned} g_x^T(x - y) &\geq f(x) - f(y) \geq g_y^T(x - y) \\ \implies G\|x - y\|_2 &\geq f(x) - f(y) \geq -G\|x - y\|_2 \end{aligned}$$

If $\|g\|_2 > G$ for some $g \in \partial f(x)$. Take $y = x + g/\|g\|_2$ such that $\|x - y\|_2 = 1$:

$$f(y) \geq f(x) + g^T(y - x) = f(x) + \|g\|_2 > f(x) + G \quad \blacksquare$$

Convergence

Theorem

Given a convex, G -Lipschitz continuous f with finite optimal value, the subgradient method obeys

$$f_{\text{best}}^k - f^* \leq \frac{R^2 + G^2 \sum_{i=0}^k t_i^2}{2 \sum_{i=0}^k t_i}$$

where $\|x^0 - x^*\|_2 \leq R$

Convergence

Proof

Key quantity: euclidean distance to optimal set
(not function value since it can go up and down)

$$\begin{aligned}\|x^{k+1} - x^*\|_2^2 &= \|x^k - t_k g^k - x^*\|_2^2 \\ &= \|x^k - x^*\|_2^2 - 2t_k (g^k)^T (x^k - x^*) + t_k^2 \|g^k\|_2^2 \\ &\leq \|x^k - x^*\|_2^2 - 2t_k (f(x^k) - f^*) + t_k^2 \|g^k\|_2^2\end{aligned}$$

using subgradient definition $f^* = f(x^*) \geq f(x^k) + (g^k)^T (x^* - x^k)$

Convergence

Proof (continued)

Combine inequalities for $i = 0, \dots, k$

$$\begin{aligned}\|x^{k+1} - x^*\|_2^2 &\leq \|x^0 - x^*\|_2^2 - 2 \sum_{i=0}^k t_i (f(x^i) - f^*) + \sum_{i=0}^k t_i^2 \|g^i\|_2^2 \\ &\leq R^2 - 2 \sum_{i=0}^k t_i (f(x^i) - f^*) + G^2 \sum_{i=0}^k t_i^2\end{aligned}$$

Using $\|x^{k+1} - x^*\|_2^2 \geq 0$ we get

$$2 \sum_{i=0}^k t_i (f(x^i) - f^*) \leq R^2 + G^2 \sum_{i=0}^k t_i^2$$

Convergence

Proof (continued)

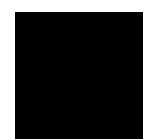
$$2 \sum_{i=0}^k t_i (f(x^i) - f^*) \leq R^2 + G^2 \sum_{i=0}^k t_i^2$$

Combine it with

$$\sum_{i=0}^k t_i (f(x^i) - f(x^*)) \geq \left(\sum_{i=0}^k t_i \right) \min_{i=0, \dots, k} (f(x^i) - f^*) = \left(\sum_{i=0}^k t_i \right) (f_{\text{best}}^k - f^*)$$

to get

$$f_{\text{best}}^k - f^* \leq \frac{R^2 + G^2 \sum_{i=0}^k t_i^2}{2 \sum_{i=0}^k t_i}$$



Implications for step size rules

$$f_{\text{best}}^k - f^* \leq \frac{R^2 + G^2 \sum_{i=0}^k t_i^2}{2 \sum_{i=0}^k t_i}$$

Fixed: $t_k = t$ for $k = 0, \dots$

$$f_{\text{best}}^k - f^* \leq \frac{R^2 + G^2(k+1)t^2}{2(k+1)t}$$

May be suboptimal

$$\lim_{k \rightarrow \infty} f_{\text{best}}^k \leq f^* + \frac{G^2 t}{2}$$

Diminishing:

$$t_k \rightarrow 0, \quad \sum_{k=0}^{\infty} t_k = \infty$$

Optimal

$$\lim_{k \rightarrow \infty} f_{\text{best}}^k = f^*$$

e.g., $t_k = \tau/(k+1)$ or $t_k = \tau/\sqrt{k+1}$

Optimal step size and convergence rate

For a tolerance $\epsilon > 0$, let's find the optimal t_k for a fixed k :

$$\frac{R^2 + G^2 \sum_{i=0}^k t_i^2}{2 \sum_{i=0}^k t_i} \leq \epsilon$$

Convex and symmetric in (t_0, \dots, t_k)

Hence, minimum when $t_i = t$

$$\longrightarrow \frac{R^2 + G^2(k+1)t^2}{2(k+1)t}$$

Optimal choice $t = \frac{R}{G\sqrt{k+1}}$

Convergence rate

$$f_{\text{best}}^k - f^* \leq \frac{RG}{\sqrt{k+1}}$$

Iterations required

$$k = O(1/\epsilon^2)$$

(gradient descent $k = O(1/\epsilon)$)

Stopping criterion

Terminating when

$$\frac{R^2 + G^2 \sum_{i=0}^k t_i^2}{2 \sum_{i=0}^k t_i} \leq \epsilon$$

is really, really slow.

Bad news

There is not really a good stopping criterion for the subgradient method

Optimal step size when f^* is known

Polyak step size

$$t_k = \frac{f(x^k) - f^*}{\|g^k\|_2^2}$$

Motivation: minimize righthand side of

$$\|x^{k+1} - x^*\|_2^2 \leq \|x^k - x^*\|_2^2 - 2t_k(f(x^k) - f^*) + t_k^2 \|g^k\|_2^2$$

Obtaining $(f(x^k) - f^*)^2 \leq (\|x^{k+1} - x^*\|_2^2 - \|x^k - x^*\|_2^2) G^2$

Applying recursively, $f_{\text{best}}^k - f^* \leq \frac{GR}{\sqrt{k+1}}$

Iterations required

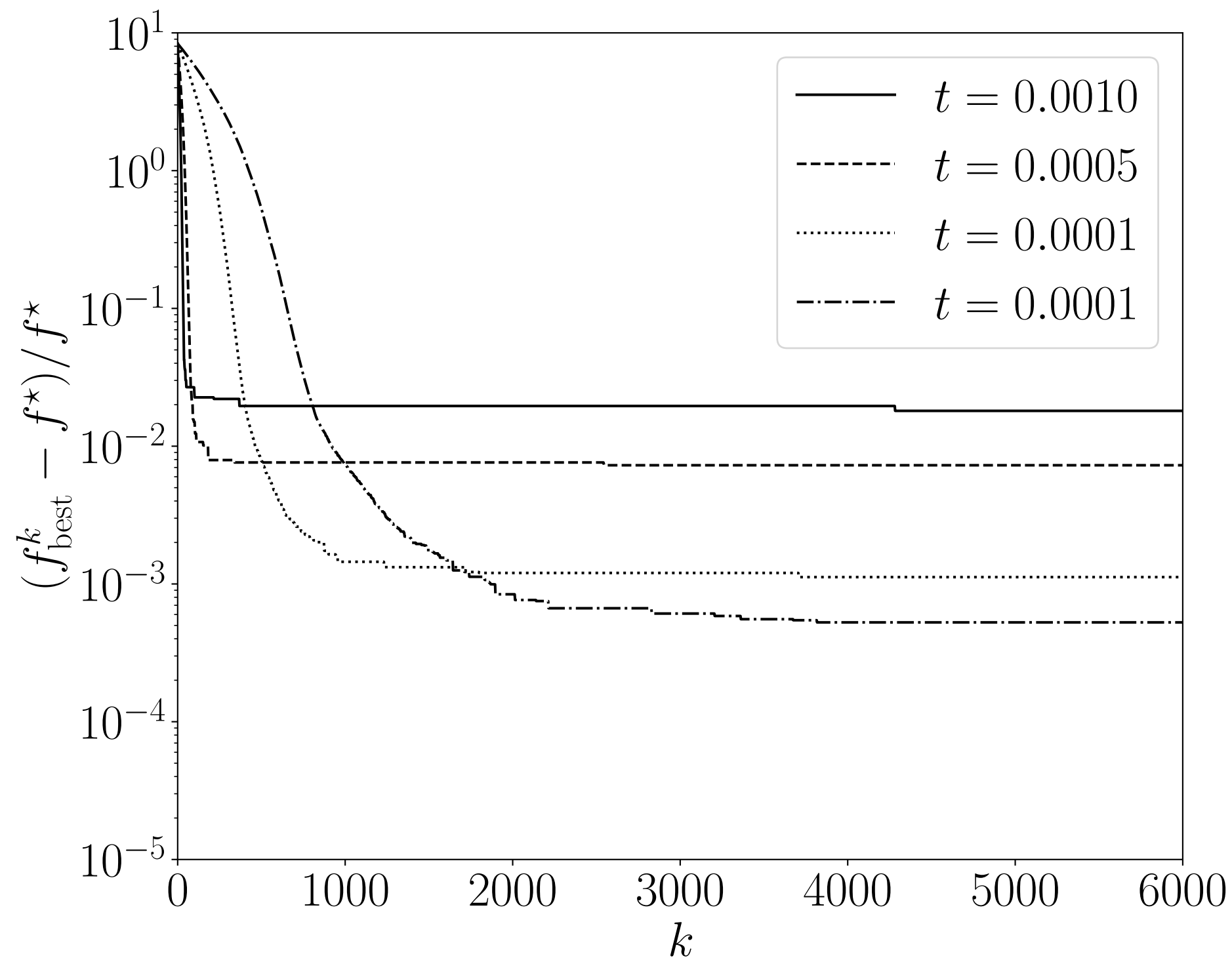
$$k = O(1/\epsilon^2)$$

still slow

Example: 1-norm minimization

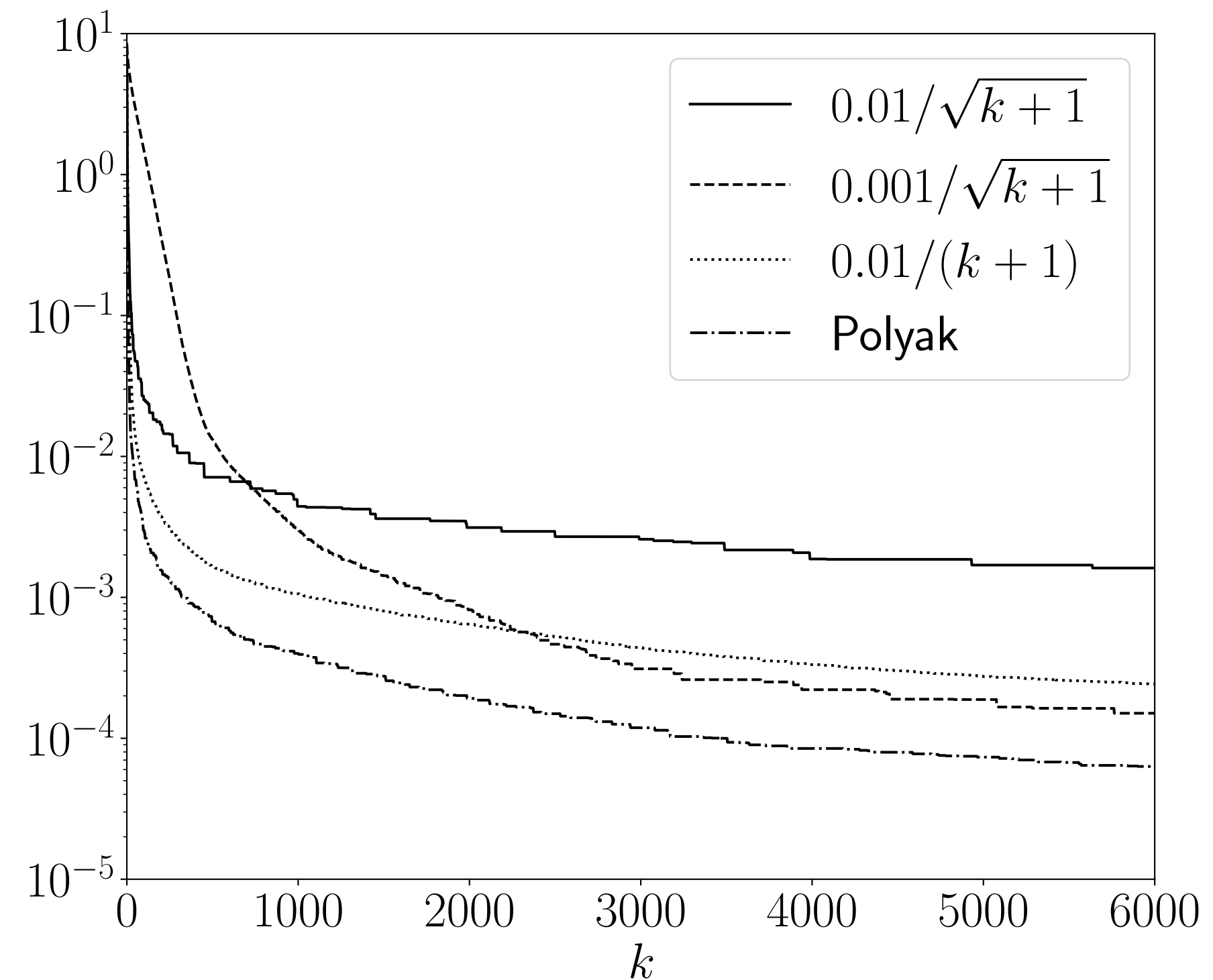
minimize $f(x) = \|Ax - b\|_1$

Fixed step size



$g = A^T \text{sign}(Ax - b) \in \partial f(x)$

Diminishing step size



Efficient packages to automatically compute (sub)gradients:

Python: JAX, PyTorch

Julia: Zygote.jl, ForwardDiff.jl, ReverseDiff.jl

Remarks on subgradient method

- Simple
- Handles general nondifferentiable convex functions
- Very slow convergence $O(1/\epsilon^2)$
- No good stopping criterion

Can we do better?

Can we incorporate constraints?

Proximal operators

Composite models

$$\text{minimize } f(x) + g(x)$$

$f(x)$ convex and smooth

$g(x)$ convex (may be not differentiable)

Examples

- Regularized regression: $g(x) = \|x\|_1$
- Constrained optimization: $g(x) = \mathcal{I}_C(x)$

Proximal operator

Definition

The **proximal operator** of the function $g : \mathbf{R}^n \rightarrow \mathbf{R}$ is

$$\text{prox}_g(x) = \underset{z}{\operatorname{argmin}} \left(g(z) + \frac{1}{2} \|z - x\|_2^2 \right)$$

Optimality conditions of prox

$$0 \in \partial g(z) + z - x \quad \Longrightarrow \quad x - z \in \partial g(z)$$

Properties

- It involves solving an optimization problem (not always easy!)
- Well-defined for CCP functions
- Easy to evaluate for many standard functions, i.e. *proxable functions*
- Generalizes many well-known algorithms

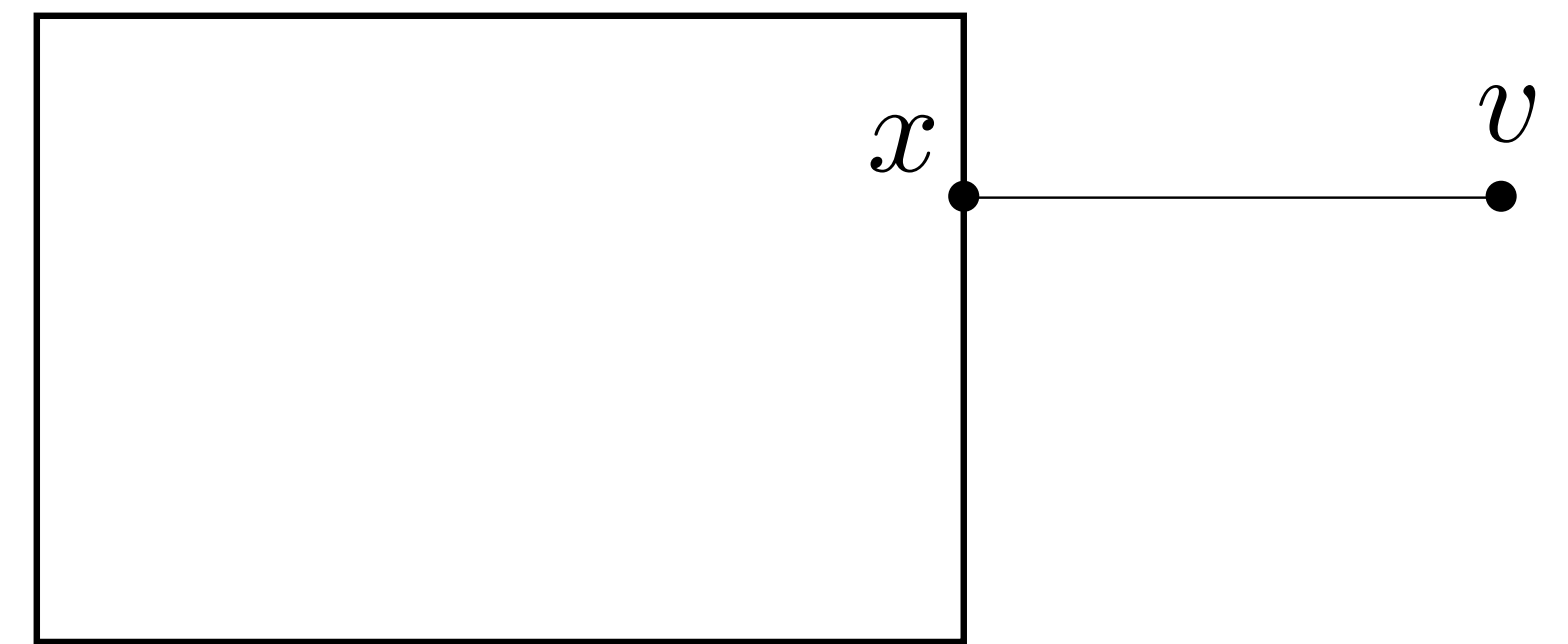
Generalized projection

The prox operator of the indicator function \mathcal{I}_C is the projection onto C

$$\mathbf{prox}_{\mathcal{I}_C}(v) = \operatorname{argmin}_{x \in C} \|x - v\|_2 = \Pi_C(v)$$

Example projection onto a box $C = \{x \mid l \leq x \leq u\}$

$$\Pi_C(v)_i = \begin{cases} l_i & v_i \leq l_i \\ v_i & l_i \leq v_i \leq u_i \\ u_i & v_i \geq u_i \end{cases}$$



Remarks

- Easy for many common sets (e.g., closed form)
- Can be “hard” for surprisingly simple sets, e.g., $C = \{Ax \leq b\}$

Quadratic functions

If $g(x) = (1/2)x^T P x + q^T x + r$ with $P \succeq 0$, then

$$\mathbf{prox}_g(v) = (I + P)^{-1}(v - q)$$

Remarks

- Closed-form always solvable (even with P not full rank)
- Symmetric, positive definite and usually sparse linear system
- Can prefactor $I + P$ and solve for different v

Separable sum

If $g(x)$ is block separable, i.e., $g(x) = \sum_{i=1}^N g_i(x_i)$

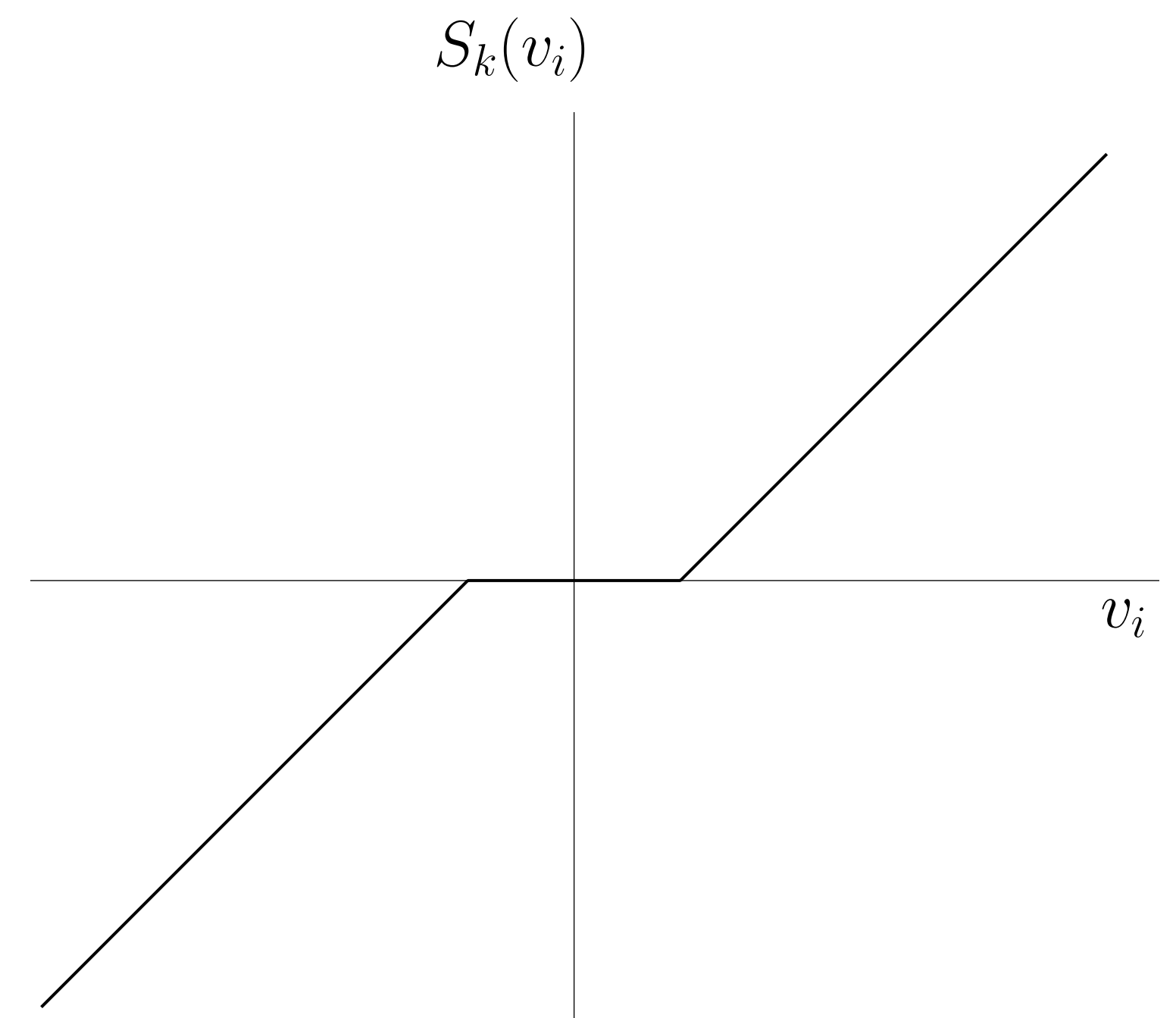
then, $(\mathbf{prox}_g(v))_i = \mathbf{prox}_{g_i}(v_i), \quad i = 1, \dots, N$

(key to parallel/distributed proximal algorithms)

Example: $g(x) = \lambda \|x\|_1 = \sum_{i=1}^n \lambda |x_i|$

soft-thresholding

$$(\mathbf{prox}_g(v))_i = \mathbf{prox}_{\lambda|\cdot|}(v_i) = S_\lambda(v_i) = \begin{cases} v_i - \lambda & v_i > \lambda \\ 0 & |v_i| \leq \lambda \\ v_i + \lambda & v_i < -\lambda \end{cases}$$



Basic rules

- **Scaling and translation:** $g(x) = ah(x) + b$ with $a > 0$, then

$$\mathbf{prox}_g(x) = \mathbf{prox}_{ah}(x)$$

Examples

- **Affine addition:** $g(x) = h(x) + a^T x + b$, then

$$\mathbf{prox}_g(x) = \mathbf{prox}_h(x - a)$$

- **Affine transformation:** $g(x) = h(ax + b)$, with $a \neq 0, a \in \mathbf{R}$,

$$\mathbf{prox}_g(x) = \frac{1}{a} (\mathbf{prox}_{a^2 h}(ax + b) - b)$$

Proofs (exercise):

- Rearrange proximal term: $(1/2)\|z - x\|_2^2$
- Apply prox optimality conditions

Proximal gradient method

Remember: gradient descent interpretation

Problem

$$\text{minimize } f(x)$$

Iterations

$$x^{k+1} = x^k - t \nabla f(x^k)$$

Quadratic approximation, replacing Hessian $\nabla^2 f(x^k)$ with $\frac{1}{t} I$

$$x^{k+1} = \underset{z}{\operatorname{argmin}} f(x^k) + \nabla f(x^k)^T (z - x^k) + \frac{1}{2t} \|z - x^k\|_2^2$$

Let's exploit the smooth part

$$\text{minimize } f(x) + g(x)$$

$f(x)$ convex and smooth

$g(x)$ convex (may be not differentiable)

Quadratic approximation of f while keeping g

$$x^{k+1} = \underset{z}{\operatorname{argmin}} g(z) + f(x^k) + \nabla f(x^k)^T (z - x^k) + \frac{1}{2t} \|z - x^k\|_2^2 \leftarrow \text{same as gradient descent}$$

Equivalent to

Proximal operator

$$x^{k+1} = \underset{z}{\operatorname{argmin}} tg(z) + \frac{1}{2} \|z - (x^k - t\nabla f(x^k))\|_2^2 = \mathbf{prox}_{tg} (x^k - t\nabla f(x^k))$$

↑
make g
small

↑
stay close to
gradient update

Proximal gradient method

minimize $f(x) + g(x)$

$f(x)$ convex and smooth

$g(x)$ convex (may be not differentiable)

Iterations

$$x^{k+1} = \text{prox}_{tg} (x^k - t \nabla f(x^k))$$

Properties

- Alternates between gradient updates of f and proximal updates on g
- Useful if prox_{tg} is inexpensive
- Can handle nonsmooth and constrained problems

Special cases

Generalized gradient descent

Smooth

$$g(x) = 0 \implies \mathbf{prox}_{tg}(x) = x$$

Constraints

$$g(x) = \mathcal{I}_C(x) \implies \mathbf{prox}_{tg}(x) = \Pi_C(x)$$

Non smooth

$$f(x) = 0$$

Problem

$$\text{minimize } f(x) + g(x)$$

Iterations

$$x^{k+1} = \mathbf{prox}_{tg}(x^k - t\nabla f(x^k))$$

Gradient descent

$$\implies x^{k+1} = x^k - t\nabla f(x^k)$$

Projected gradient descent

$$\implies x^{k+1} = \Pi_C(x^k - t\nabla f(x^k))$$

Proximal minimization

$$\implies x^{k+1} = \mathbf{prox}_{tg}(x^k)$$

Note: useful if \mathbf{prox}_{tg} is cheap ³⁰

What happens if we cannot evaluate the prox?

At every iteration, it can be very expensive to evaluate

$$\mathbf{prox}_g(x) = \operatorname{argmin}_z \left(g(z) + \frac{1}{2} \|z - x\|_2^2 \right)$$

Idea: solve it approximately!

If you precisely control the $\mathbf{prox}_g(x)$ evaluation errors you can obtain the same convergence guarantees (and rates) as the exact evaluations.

Example: Lasso

Iterative Soft Thresholding Algorithm (ISTA)

$$\text{minimize } \underbrace{(1/2)\|Ax - b\|_2^2}_{f(x)} + \underbrace{\lambda\|x\|_1}_{g(x)}$$

Proximal gradient descent

$$x^{k+1} = \text{prox}_{tg} \left(x^k - t \nabla f(x^k) \right)$$

$$\nabla f(x) = A^T (Ax - b)$$

$$\text{prox}_{tg}(x) = S_{\lambda t}(x) \quad (\text{component wise soft-thresholding})$$

Closed-form iterations

$$x^{k+1} = S_{\lambda t} \left(x^k - t A^T (Ax^k - b) \right)$$

Example: Lasso

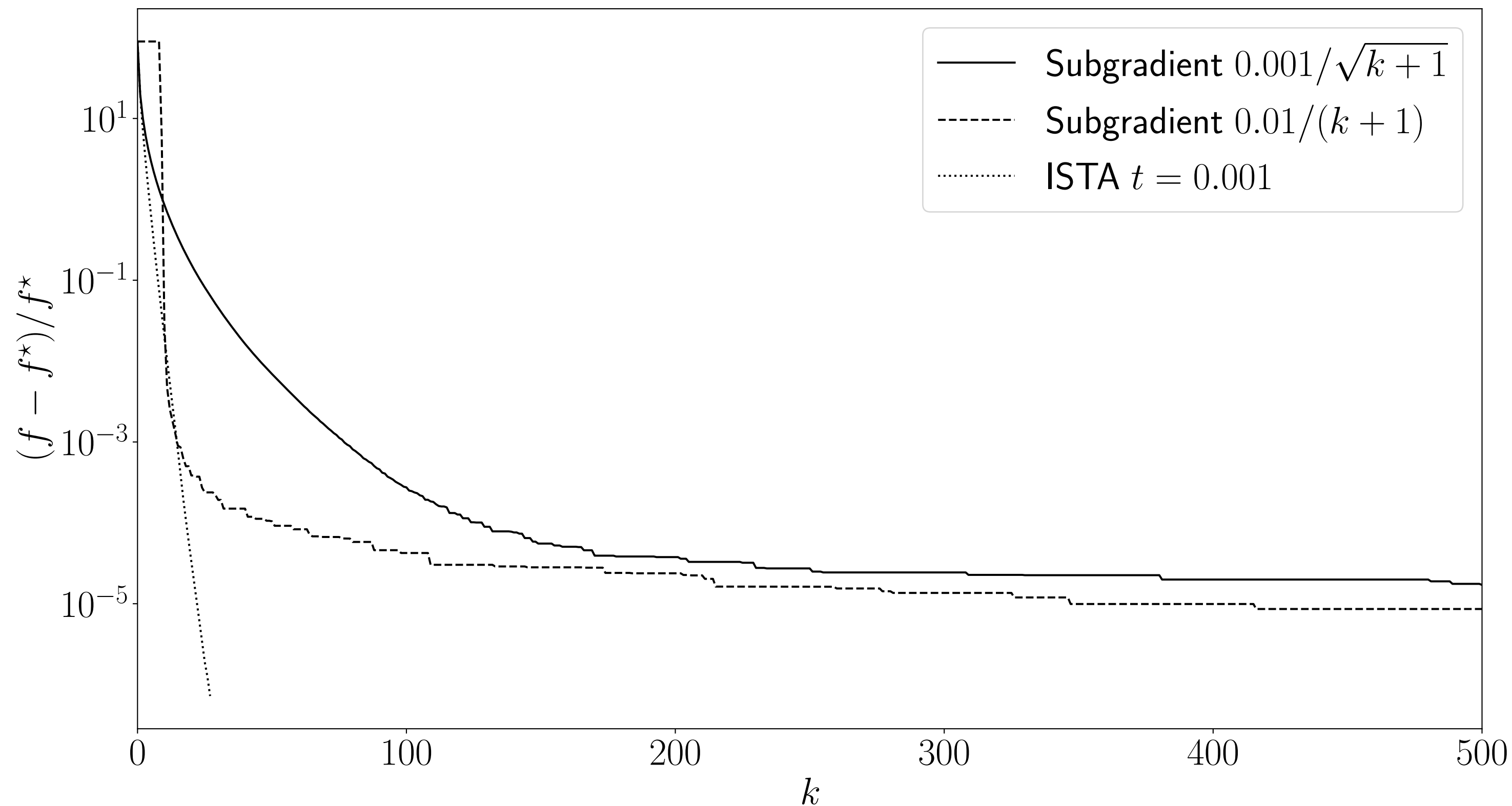
Iterative Soft Thresholding Algorithm (ISTA)

$$A \in \mathbf{R}^{500 \times 100}$$

$$\text{minimize} \quad (1/2)\|Ax - b\|_2^2 + \lambda\|x\|_1$$

Closed-form iterations

$$x^{k+1} = S_{\lambda t} \left(x^k - tA^T (Ax^k - b) \right)$$



Better convergence

Can we prove convergence generally?

Can we combine different operators?

Subgradient method and proximal operators

Today, we learned to:

- **Develop** subgradient method and **analyze** its convergence
- **Define** and **evaluate** proximal operators for various common functions
- **Apply** proximal operators to generalize gradient descent (vanilla, projected, proximal)

Next lecture

- Operator theory