

The Voice of Optimization

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joint work with Dimitris Bertsimas

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Inventory management

minimize $\sum_{t=0}^{T-1} h(x_t) + o(u_t)$

subject to $x_{t+1} = x_t + u_t - d_t$

$x_0 = x_{\text{init}}$

$0 \leq u_t \leq M$

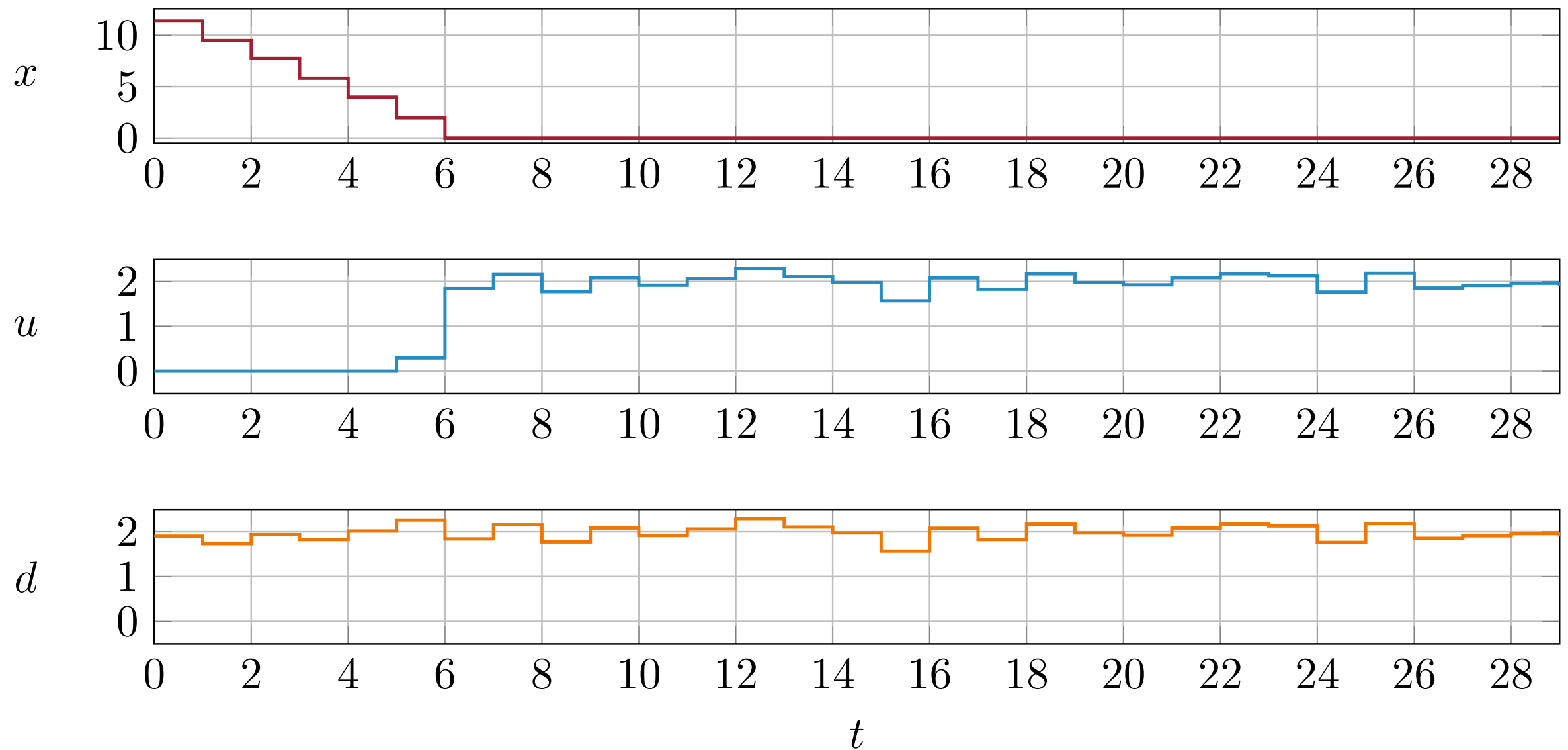
Inventory

Order

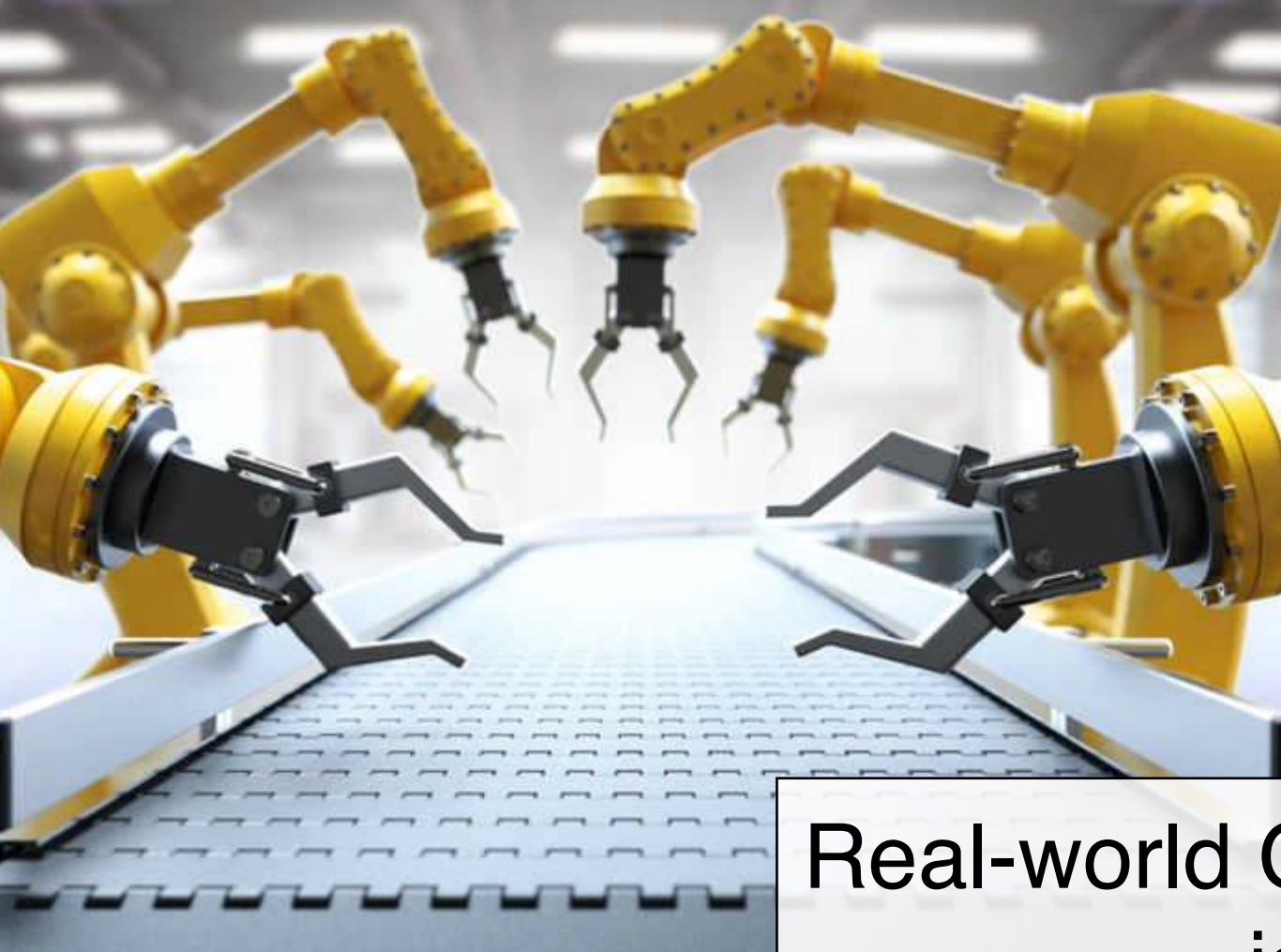
Demand

Parameters

What does it mean?

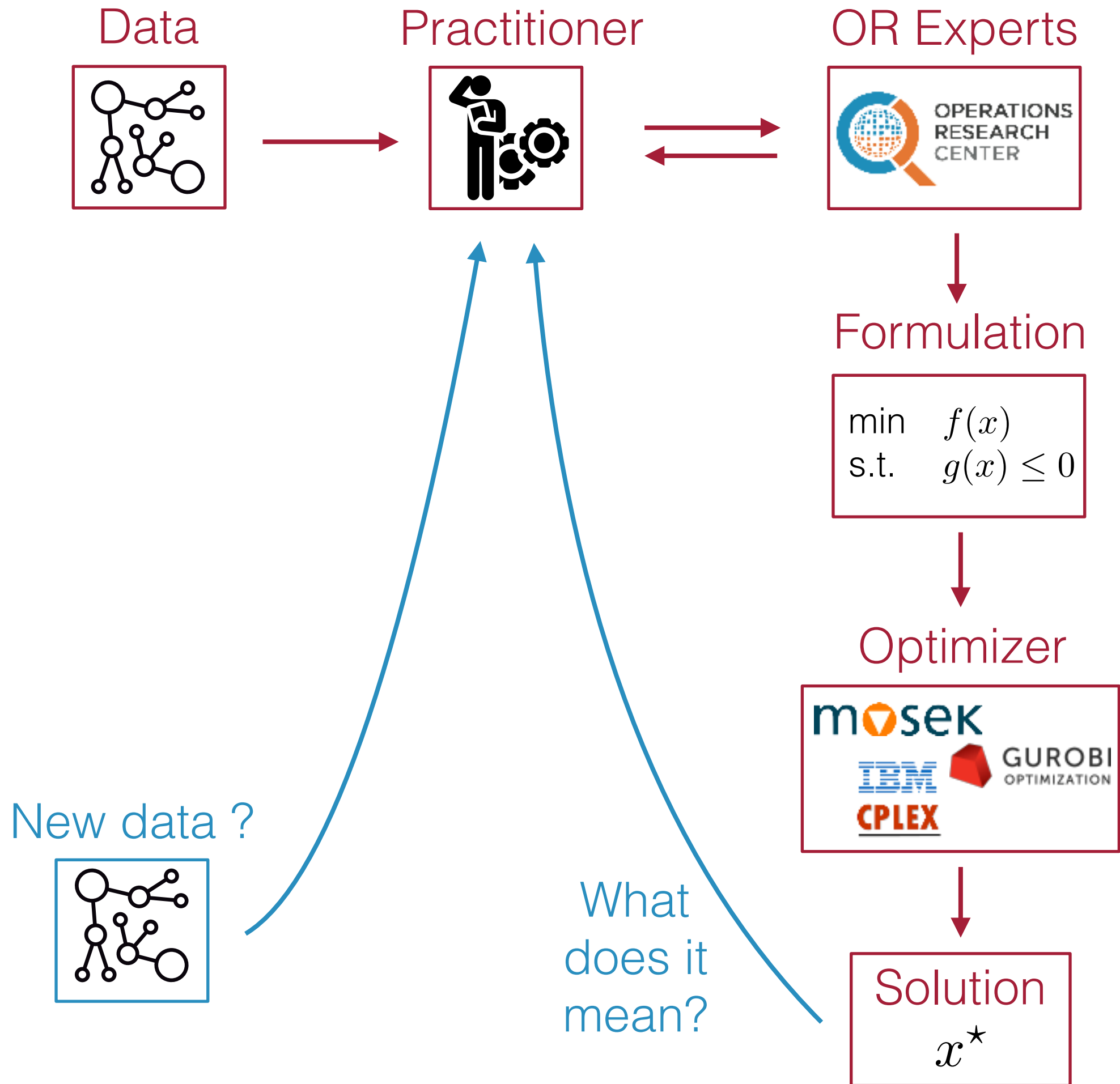


x_{init} ? d_t ?



Real-world Optimization
is
Parametric

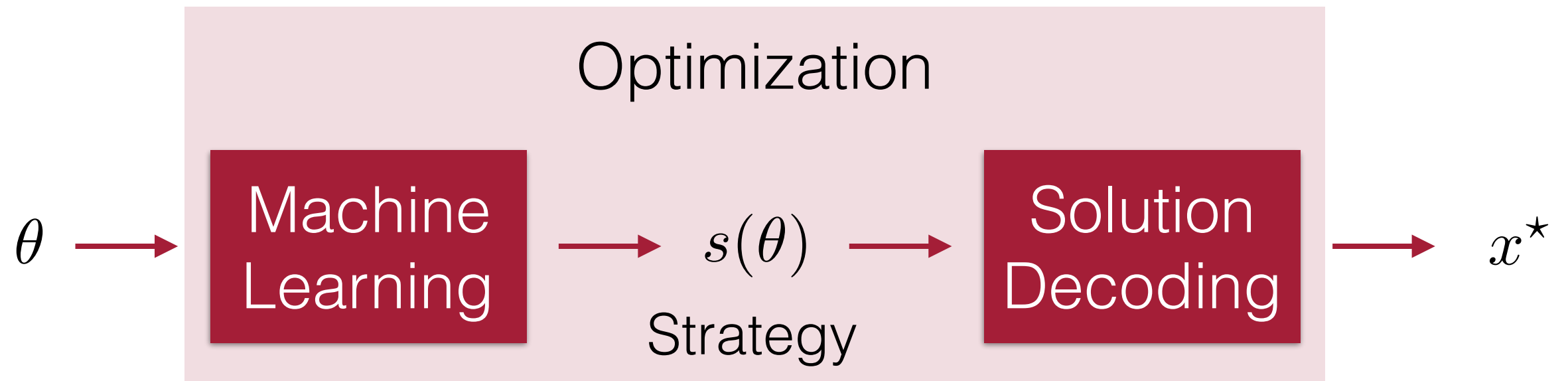






Black box $\longrightarrow$ Open the box

Too slow $\longrightarrow$ Solve in milliseconds



Black box $\longrightarrow$ Open the box

Too slow $\longrightarrow$ Solve in milliseconds

Optimal strategies

Strategy

*The complete information to efficiently get
the optimal solution*

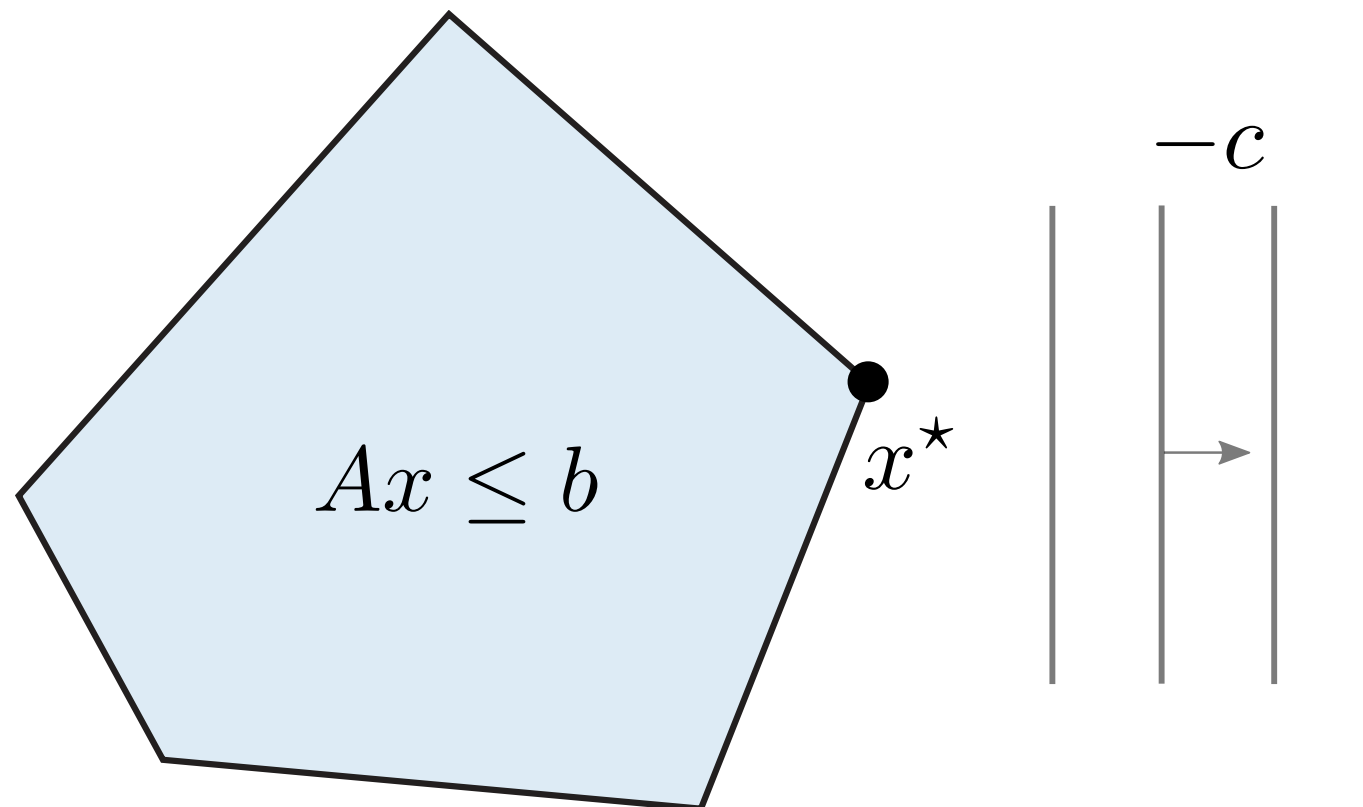
$$s(\theta)$$

Linear optimization

$$\begin{array}{ll} \text{minimize} & c(\theta)^T x \\ \text{subject to} & A(\theta)x \leq b(\theta) \end{array}$$

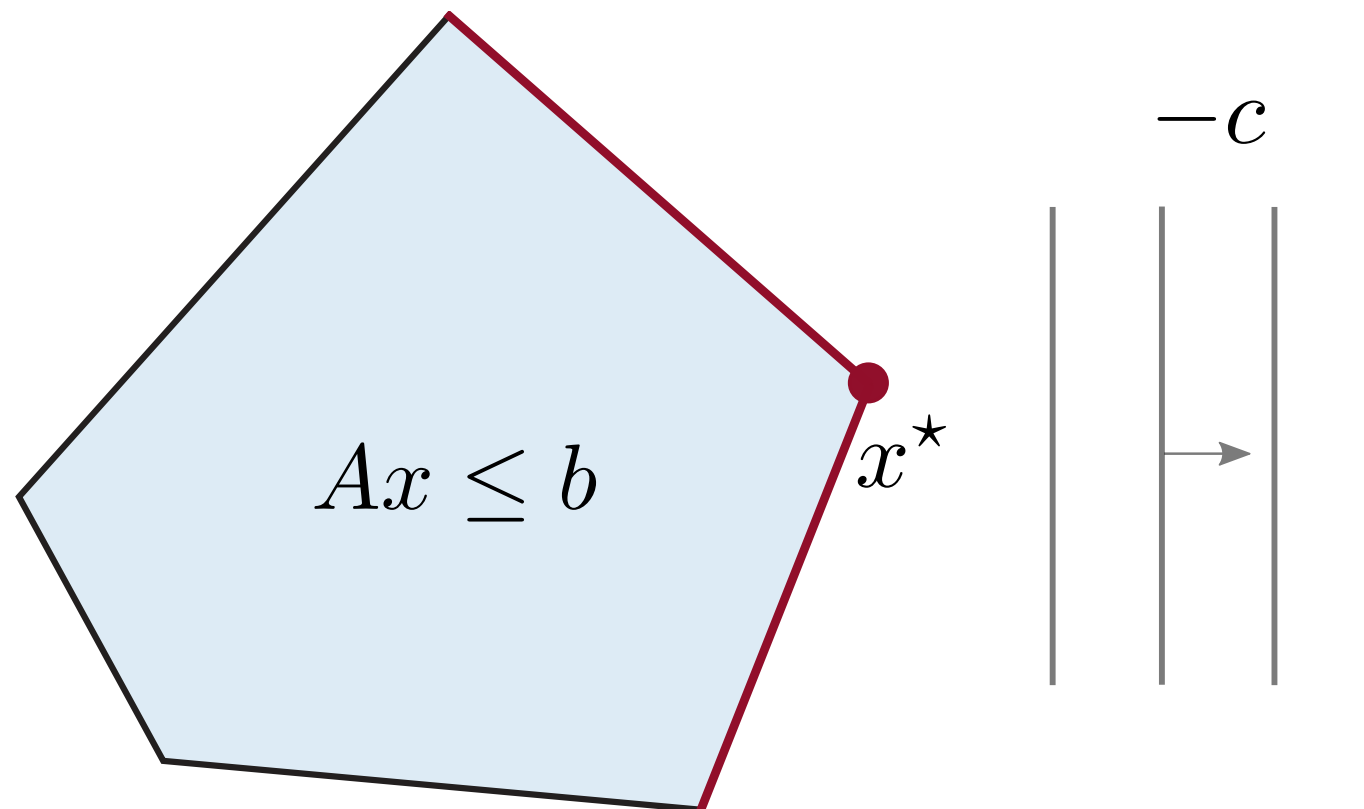
Parameters

Variables



Tight constraints

$$\mathcal{T}(\theta) = \{i \mid A_i(\theta)x^* = b_i(\theta)\}$$



$|\mathcal{T}(\theta)| = \# \text{ variables}$

$|\mathcal{T}(\theta)| \ll \# \text{ constraints}$

if non-degenerate

in general

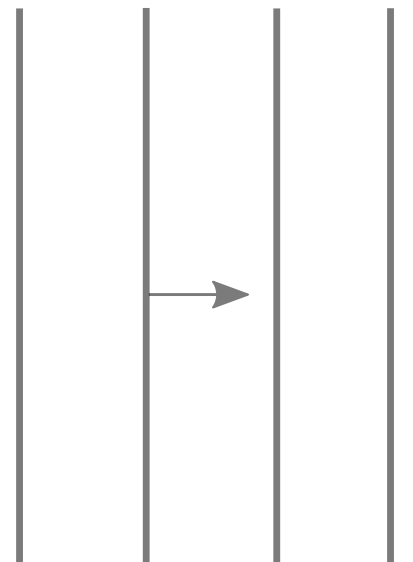
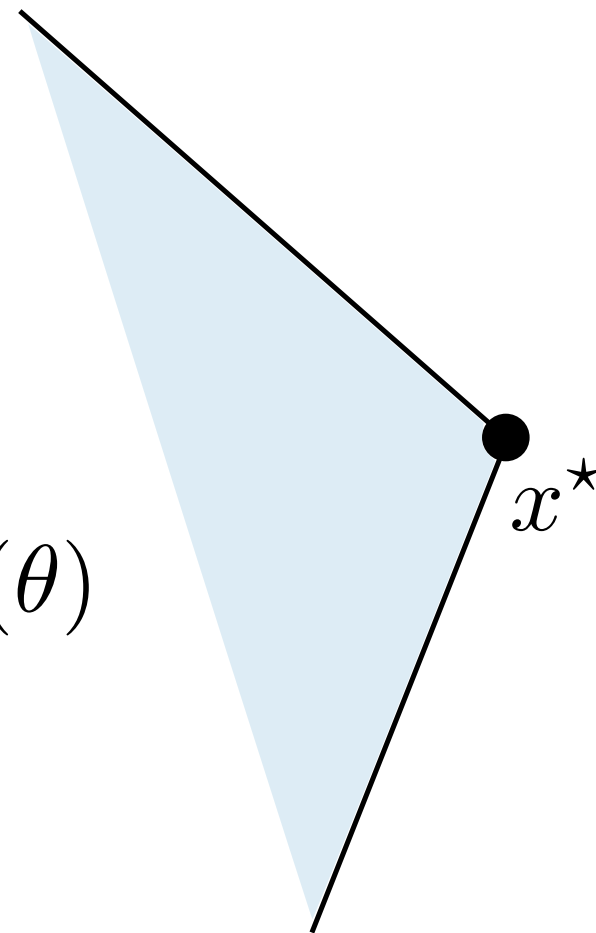
Strategy

*The complete information to efficiently get
the optimal solution*

$$s(\theta) = \mathcal{T}(\theta)$$



minimize $c(\theta)^T x$
 subject to $A_i(\theta)x = b_i(\theta) \quad \forall i \in \mathcal{T}(\theta)$



Inventory management

minimize $\sum_{t=0}^{T-1} h(x_t) + o(u_t)$

subject to $x_{t+1} = x_t + u_t - d_t$

$x_0 = x_{\text{init}}$

$0 \leq u_t \leq M$

Inventory

Inequalities

Parameters

Order

Demand

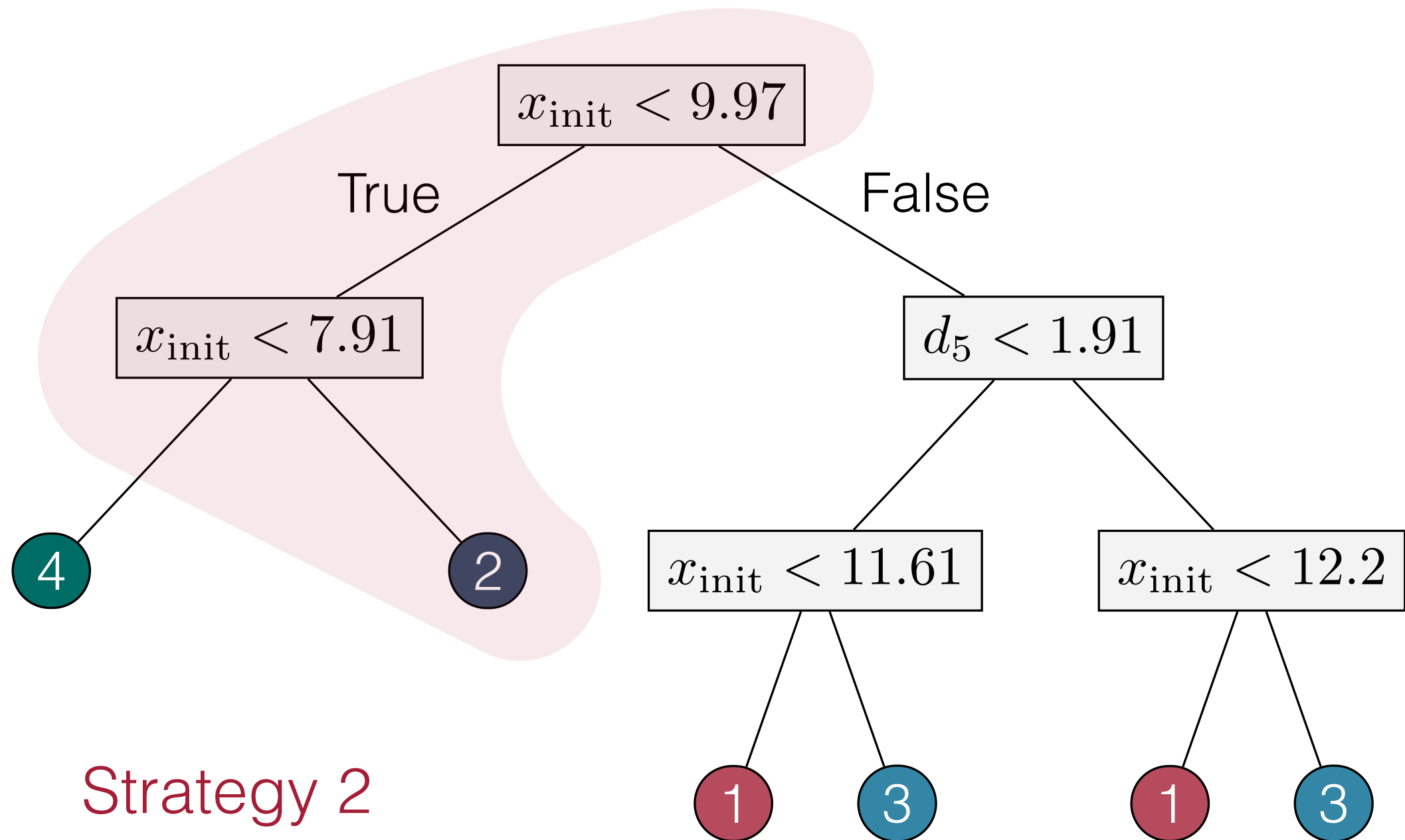
The diagram illustrates the inventory management problem with the following components:

- Equation 1:** $x_{t+1} = x_t + u_t - d_t$. The terms are color-coded: x_{t+1} is in a pink box, x_t is in a light blue box, u_t is in a light blue box, and d_t is in a grey box.
- Equation 2:** $x_0 = x_{\text{init}}$. The term x_{init} is in a light blue box.
- Equation 3:** $0 \leq u_t \leq M$. The entire expression is in a pink box.

Arrows indicate the relationships between the variables and the labels:

- A red arrow points from the word "Inventory" to the x_t term in Equation 1.
- A red arrow points from the word "Inequalities" to the pink box containing Equation 3.
- A teal arrow points from the word "Parameters" to the u_t term in Equation 1.
- A blue arrow points from the word "Order" to the u_t term in Equation 1.
- An orange arrow points from the word "Demand" to the d_t term in Equation 1.

Strategy selection



Strategy 2

$$u_t = 0 \quad t \leq 4$$

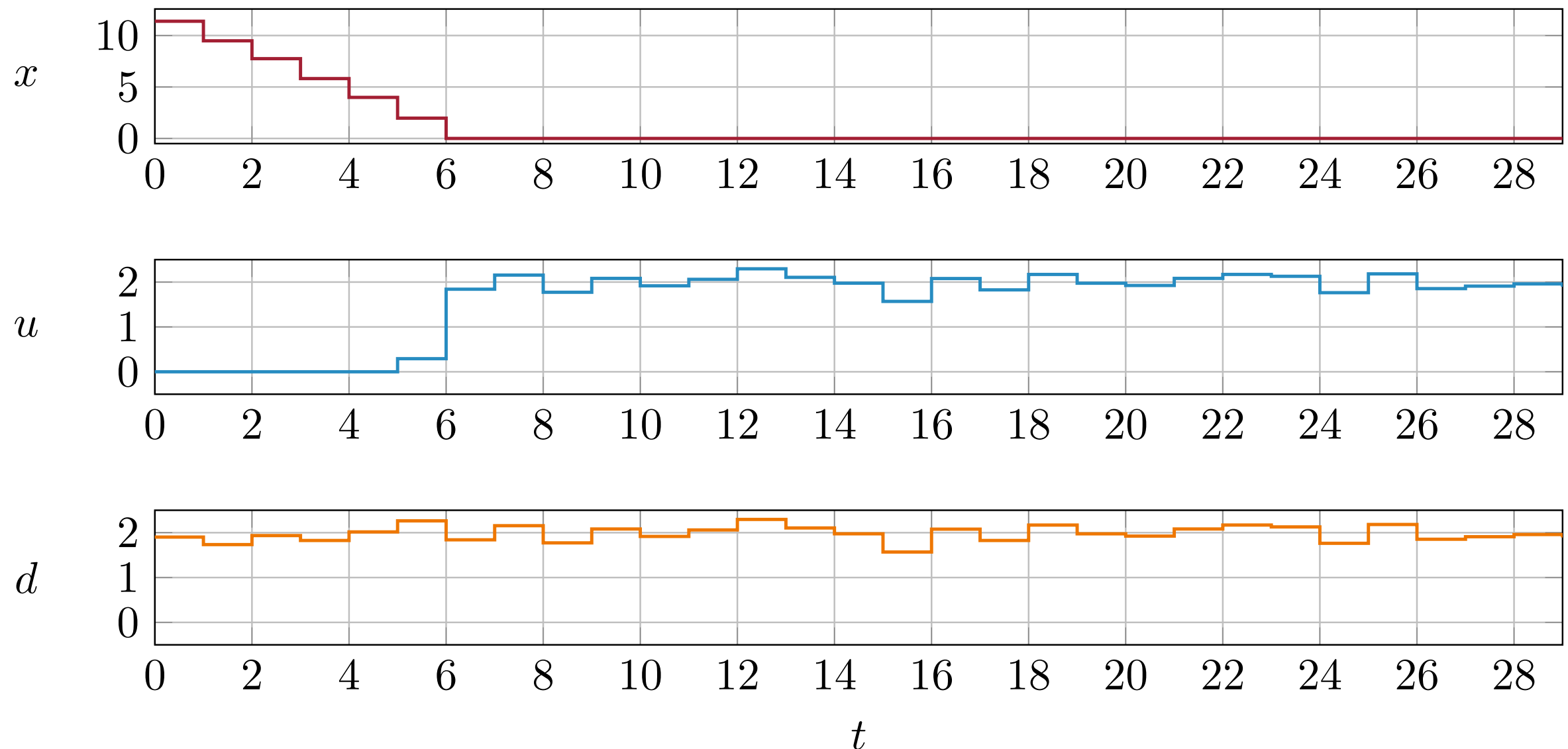
$$0 \leq u_t \leq M \quad t > 4$$

Strategies for inventory

Strategy 2

$$u_t = 0 \quad t \leq 4$$

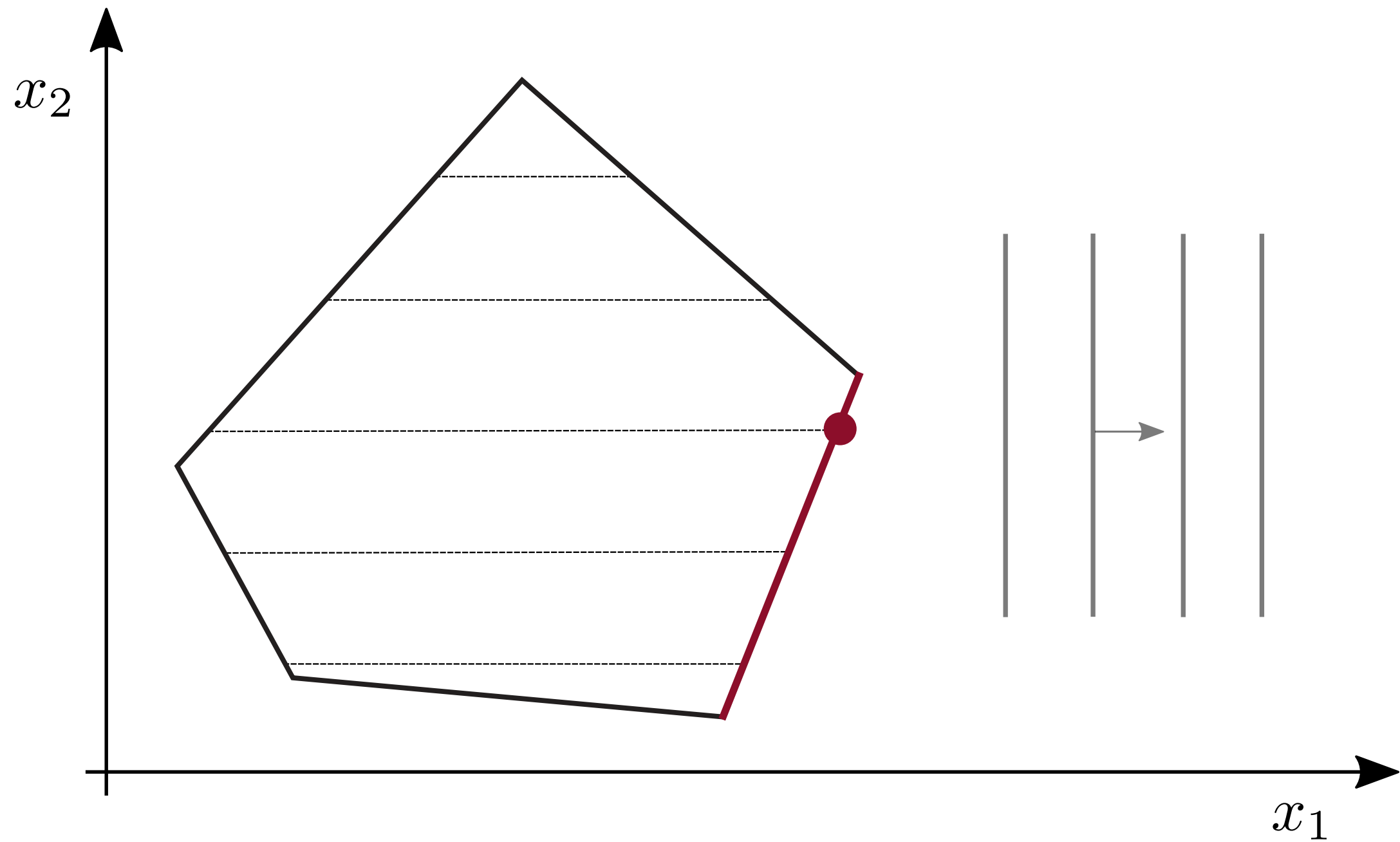
$$0 \leq u_t \leq M \quad t > 4$$



Mixed-integer optimization

$$\begin{array}{ll} \text{minimize} & c(\theta)^T x \\ \text{subject to} & A(\theta)x \leq b(\theta) \\ & x_{\mathcal{I}} \in \mathbf{Z}^d \quad \text{Integers} \end{array}$$

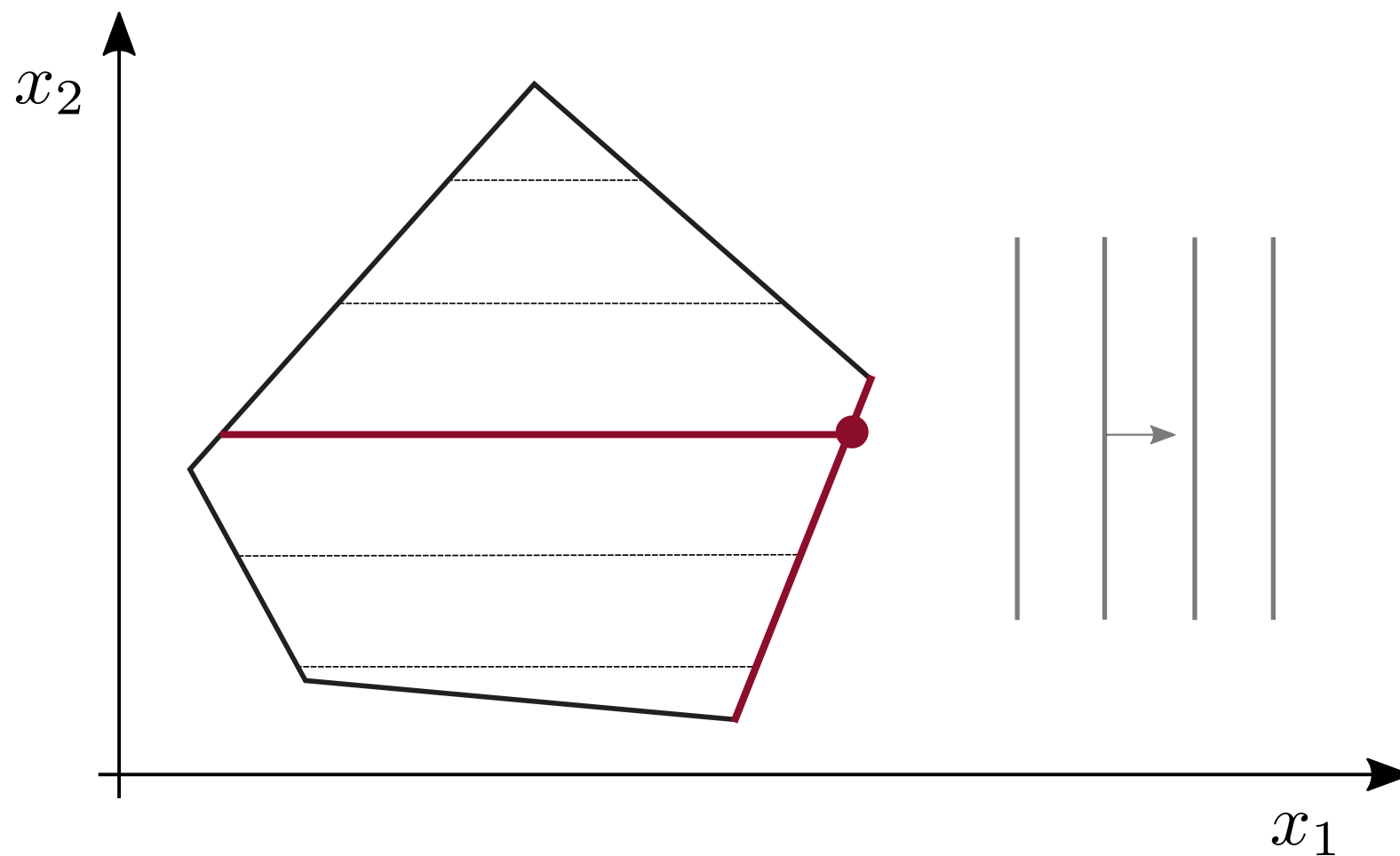
Tight constraints are not enough



Strategy

$$s(\theta) = (\mathcal{T}(\theta), x_{\mathcal{I}}^*(\theta))$$

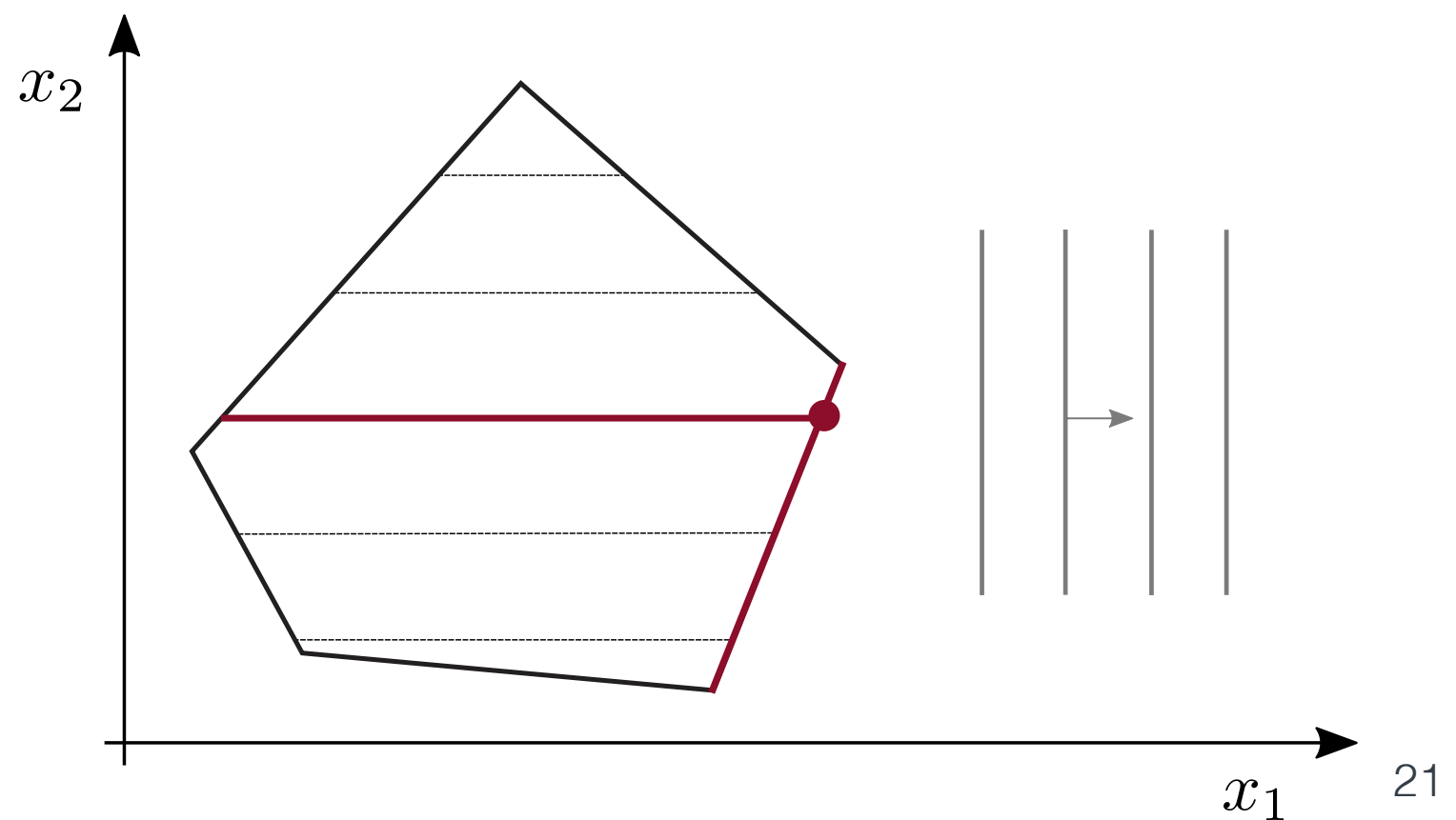
Integer variables





minimize $c(\theta)^T x$
 subject to $A_i(\theta)x = b_i(\theta) \quad \forall i \in \mathcal{T}(\theta)$
 $x_{\mathcal{I}} = x_{\mathcal{I}}^\star(\theta)$

↑
Integers



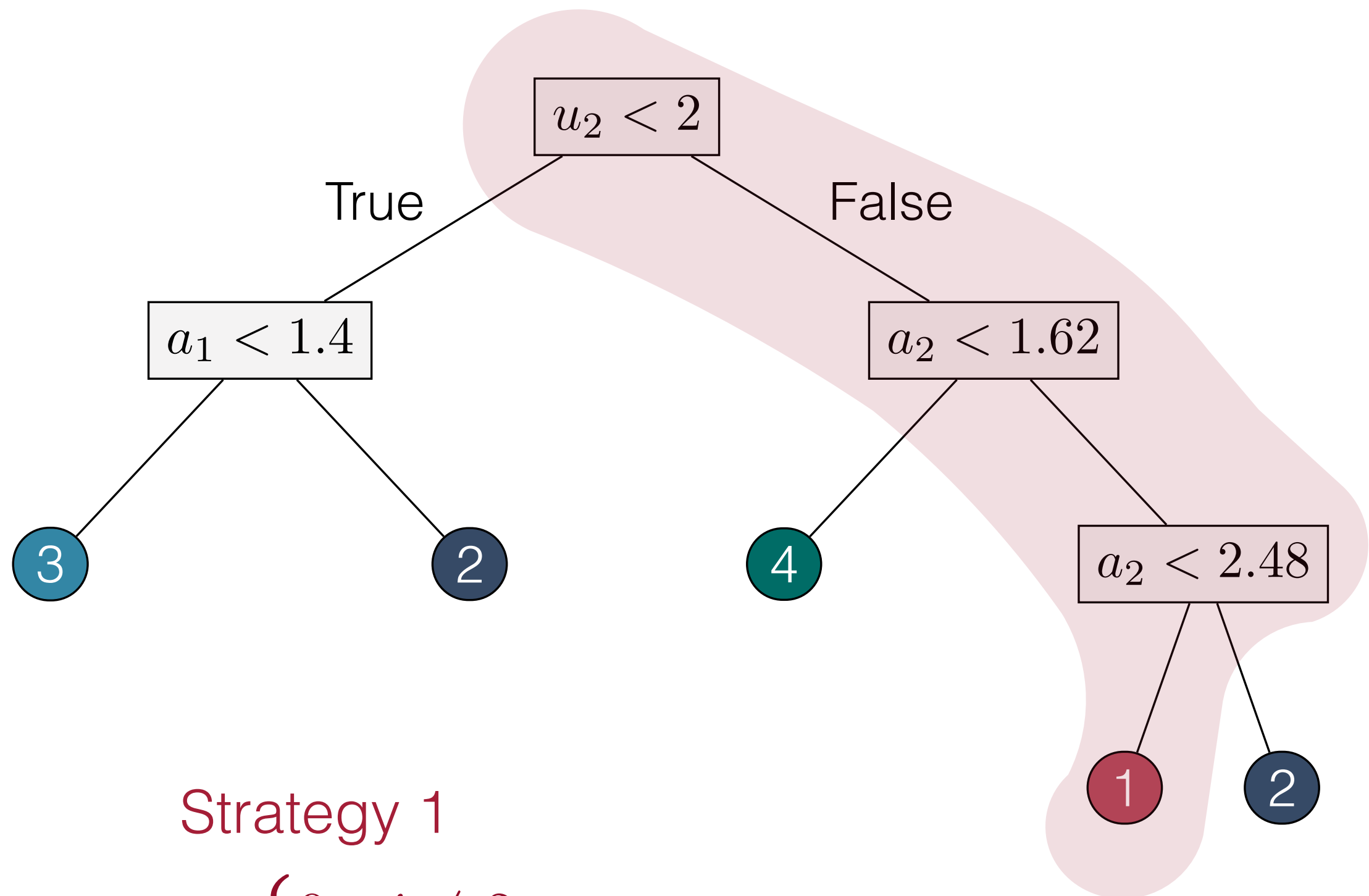
Knapsack

$$\begin{array}{ll}\text{maximize} & c^T x \\ \text{subject to} & a^T x \leq b \\ & 0 \leq x \leq u \\ & x \in \mathbf{Z}^n\end{array}$$

Parameters



Strategy selection

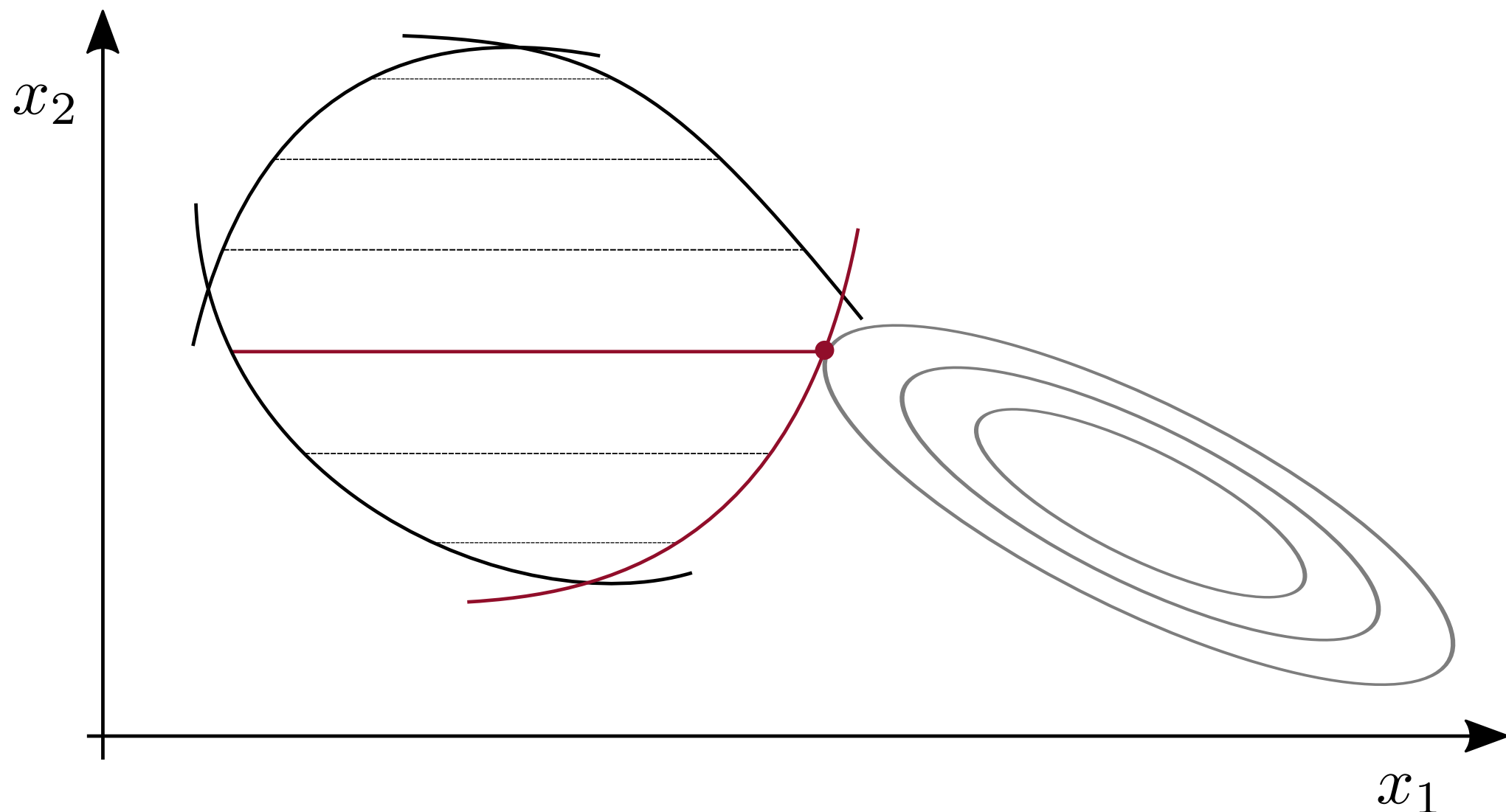


Strategy 1

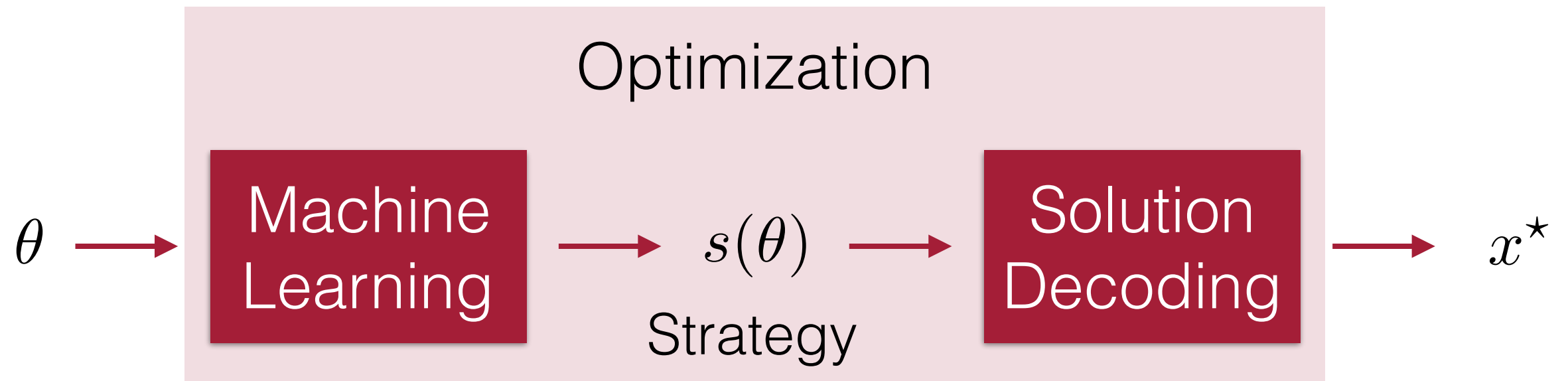
$$x_i = \begin{cases} 0 & i \neq 2 \\ 2 & i = 2 \end{cases}$$

Mixed-integer convex optimization

$$\begin{array}{ll}\text{minimize} & f(\theta, x) \\ \text{subject to} & g(\theta, x) \leq 0 \\ & x_{\mathcal{I}} \in \mathbf{Z}^d\end{array}$$



Learning the strategies



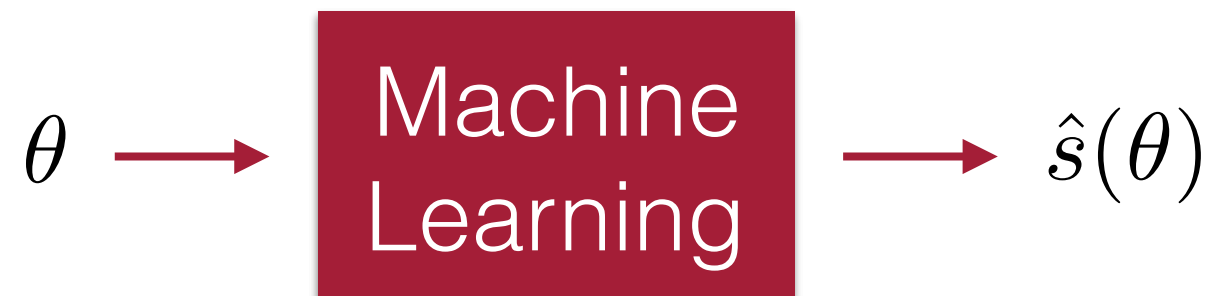
Black box	→	Open the box
Too slow	→	Solve in milliseconds

Classification problem

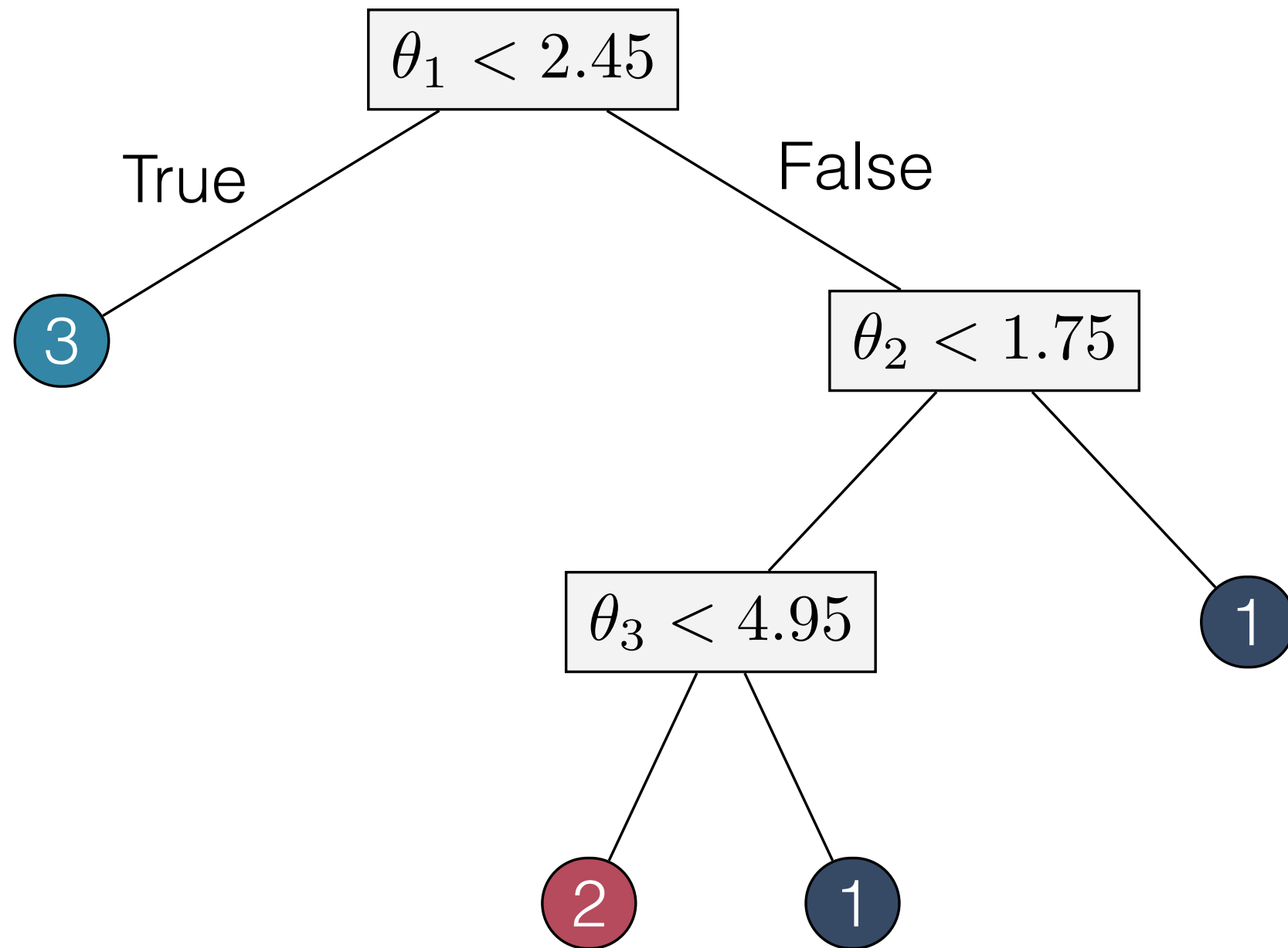
N data
 $(\theta_i, s(\theta_i))$

M labels
Strategies $\mathcal{S}$

Learn multiclass classifier



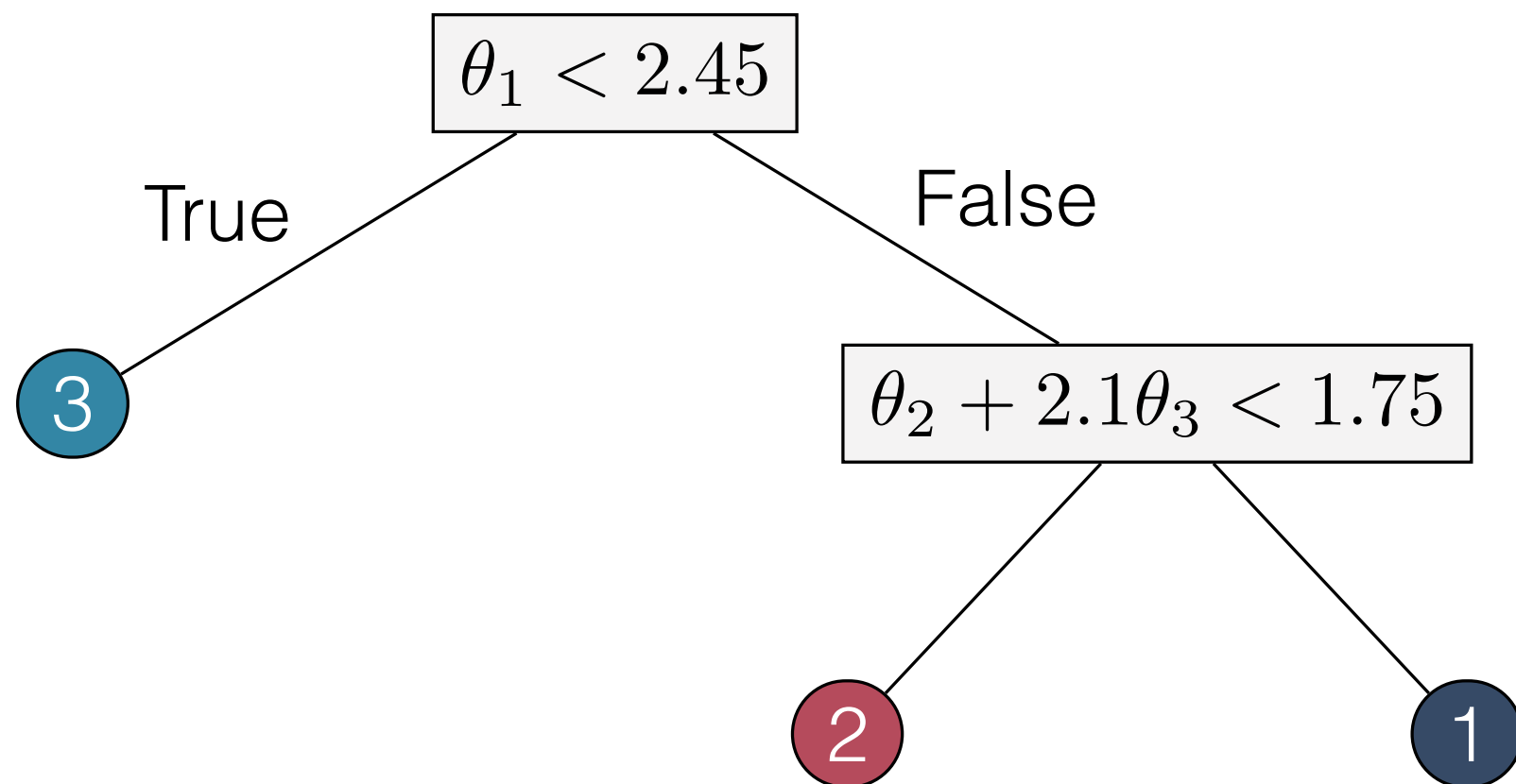
Optimal Classification Trees (OCT)



Interpretable AI

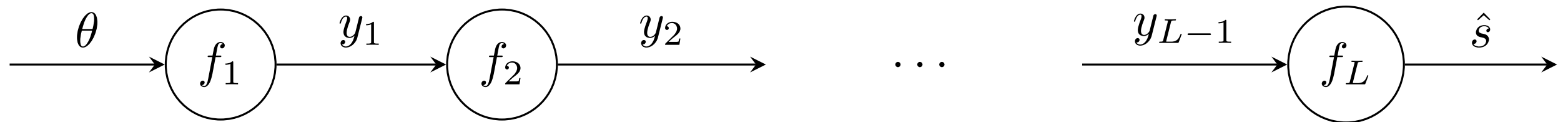
[Bertsimas and Dunn (2017)]

Optimal Classification Trees with Hyperplanes (OCT-H)



Interpretable AI
[Bertsimas and Dunn (2017)]

Neural Networks (NN)



Single layer

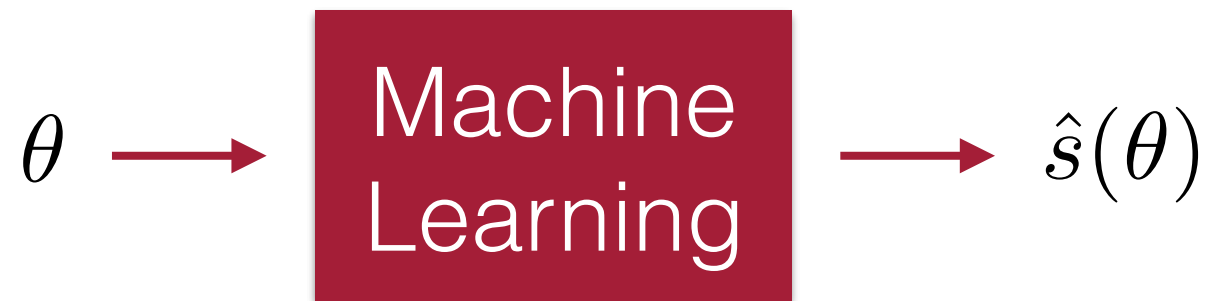
$$y_l = f(y_{l-1}) = (W_l y_{l-1} + b_l)_+$$

Strategies exploration

Classification problem

N data
 $(\theta_i, s(\theta_i))$

M labels
Strategies $\mathcal{S}$

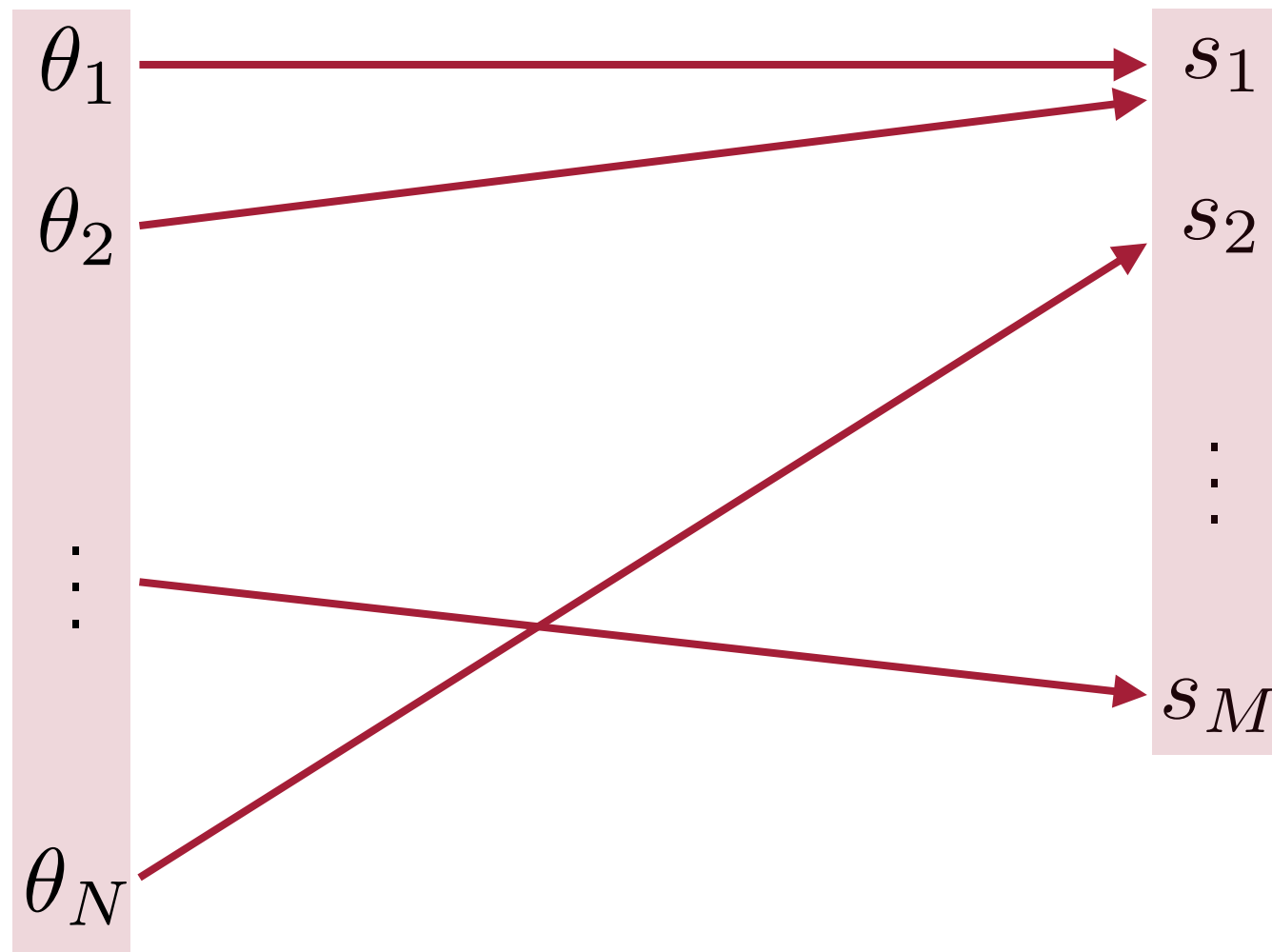


Have we seen enough data?

Will we find new strategies?

Parameters

Strategies



$\theta_{N+1}?$

Alan Turing already knew that

Decoded the Enigma Machine
in
World War II

Only some words (labels) needed



Good-Turing estimator

s_1	12 times
s_2	45 times
$\vdots$	$\vdots$
s_M	2 times

$GT \approx$ probability of unseen strategies

$$GT = \frac{N_1}{N}$$

strategies appeared once
samples

Sampling scheme

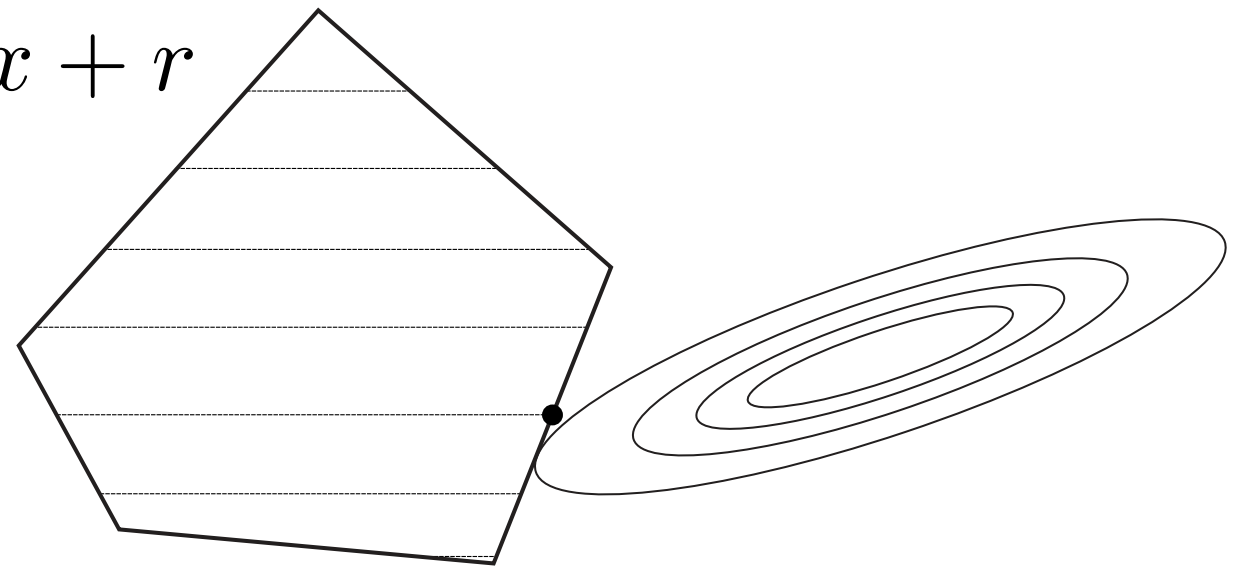
Repeat until $GT \leq \epsilon$

1. sample θ_i
2. compute $s(\theta_i)$
3. update estimator $GT \leftarrow \frac{N_1}{N}$

Speedups

Mixed-integer quadratic optimization

minimize $(1/2)x^T P x + q^T x + r$
 subject to $Ax \leq b$
 $x_{\mathcal{I}} \in \mathbf{Z}^d$



Strategy

minimize $(1/2)x^T P x + q^T x + r$
 subject to $A_{\mathcal{T}(\theta)} x = b_{\mathcal{T}(\theta)}$
 $x_{\mathcal{I}} = x_{\mathcal{I}}^*(\theta)$

Equality
Constrained
QP

Quick solution

KKT System

$$\begin{bmatrix} P & A_{\mathcal{T}(\theta)}^T & I_{\mathcal{I}}^T \\ A_{\mathcal{T}(\theta)} & 0 & 0 \\ I_{\mathcal{I}} & 0 & 0 \end{bmatrix} \begin{bmatrix} x \\ \nu \end{bmatrix} = \begin{bmatrix} -q \\ b_{\mathcal{T}(\theta)} \\ x_{\mathcal{I}}^*(\theta) \end{bmatrix}$$

$\mathcal{NP}$ -hard problem



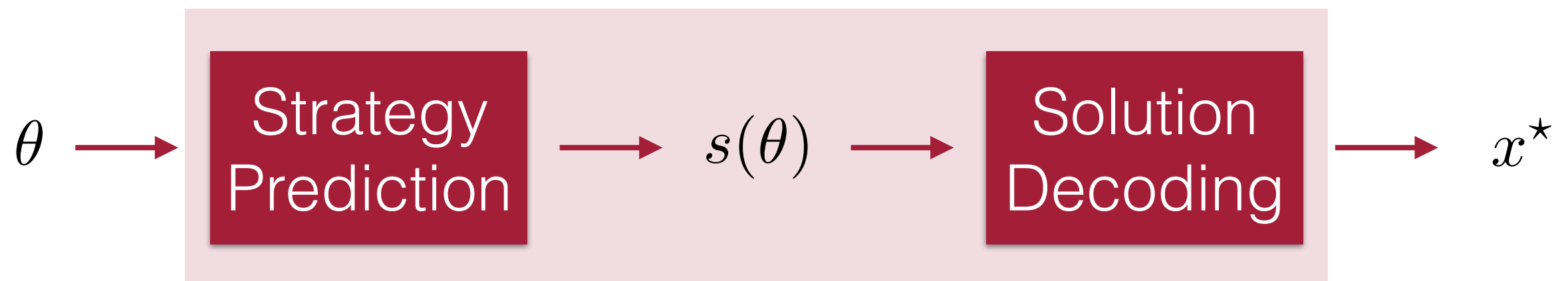
Linear System

MLOPT

Offline Learning

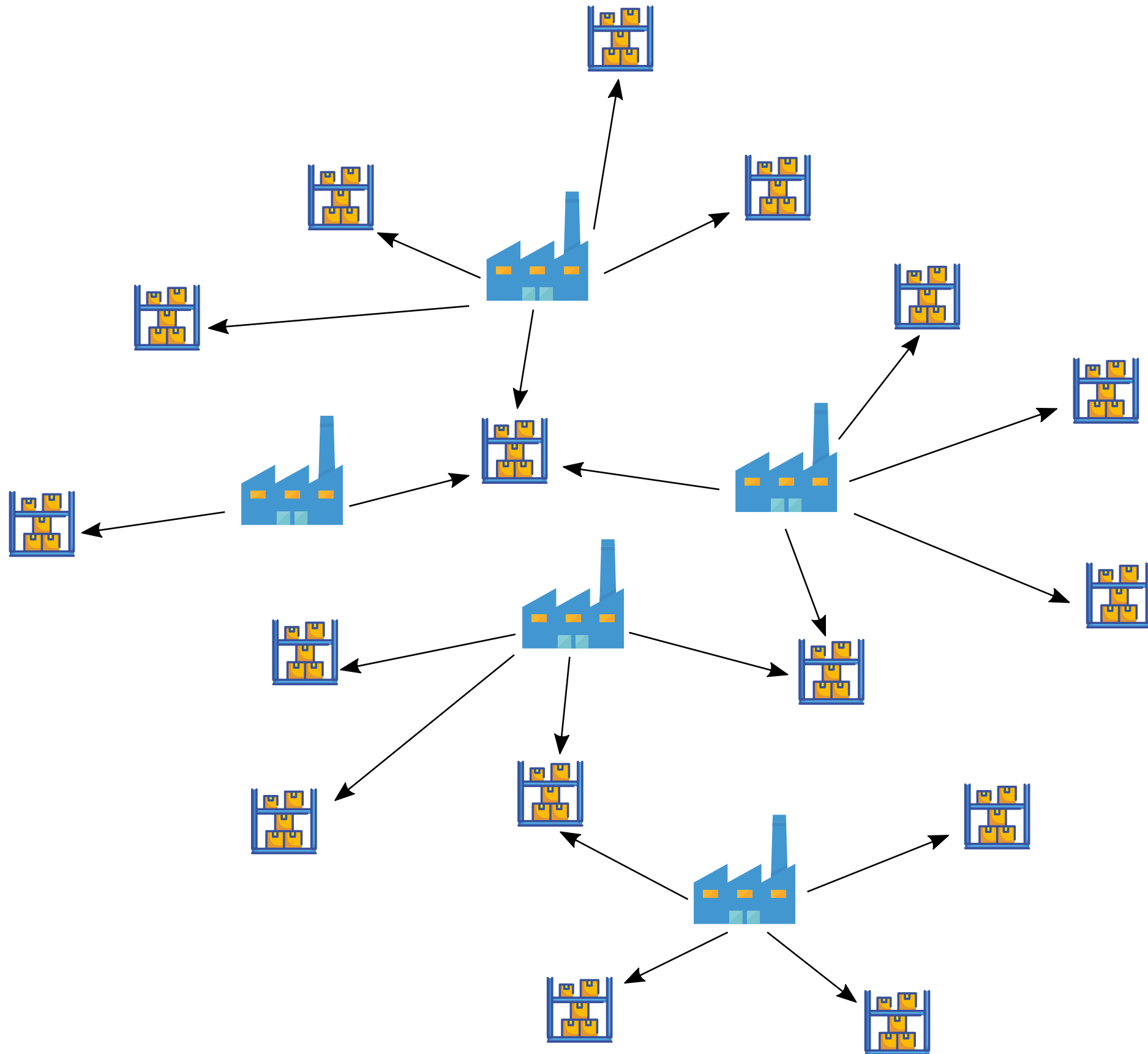


Online Optimization



Benchmarks

Facility location



Facility location

minimize $\sum_{i \in F} \sum_{j \in W} \overset{\text{Shipment cost}}{c_{ij} x_{ij}} + \sum_{i \in I} \overset{\text{Facility cost}}{f_i y_i}$

subject to $\sum_{i \in F} x_{ij} \geq \overset{\text{Demand}}{d_j} \quad \forall j \in W$

$\sum_{j \in W} x_{ij} \leq s_i y_i, \quad \forall i \in F$ Supply

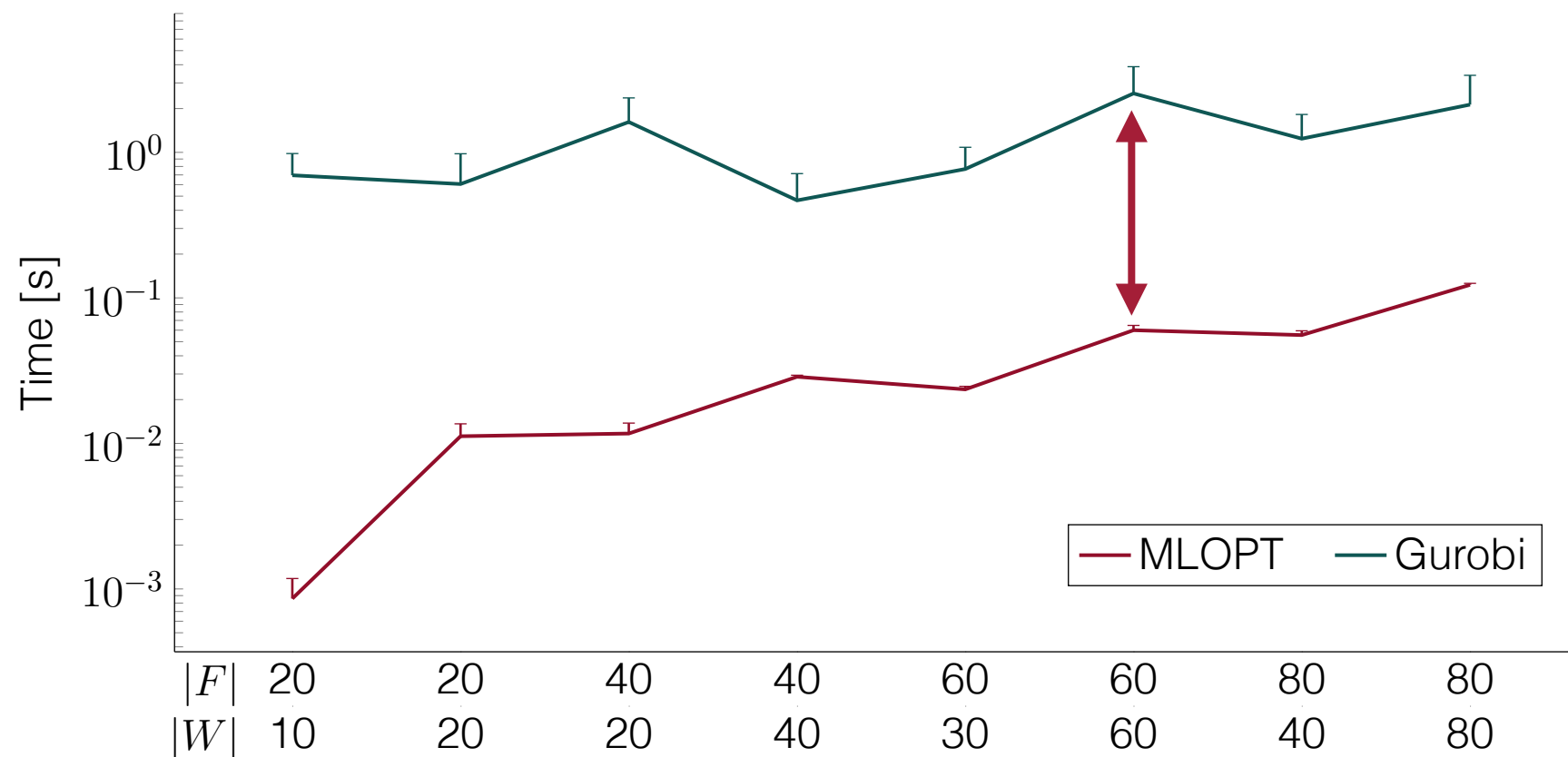
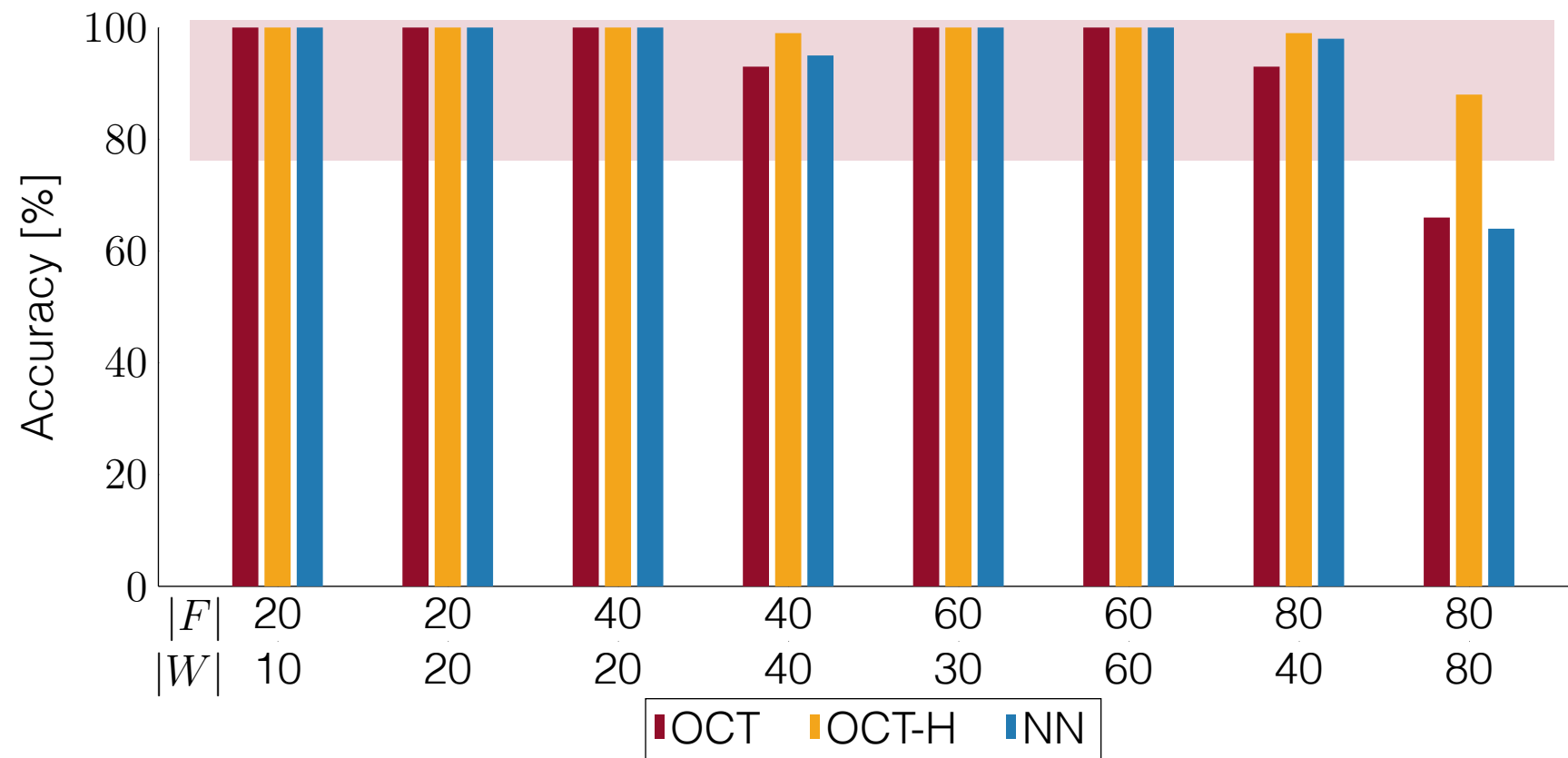
$\overset{\text{Shipment}}{x_{ij}} \geq 0, \overset{\text{Facility}}{y_i} \in \{0, 1\}$

Parameters

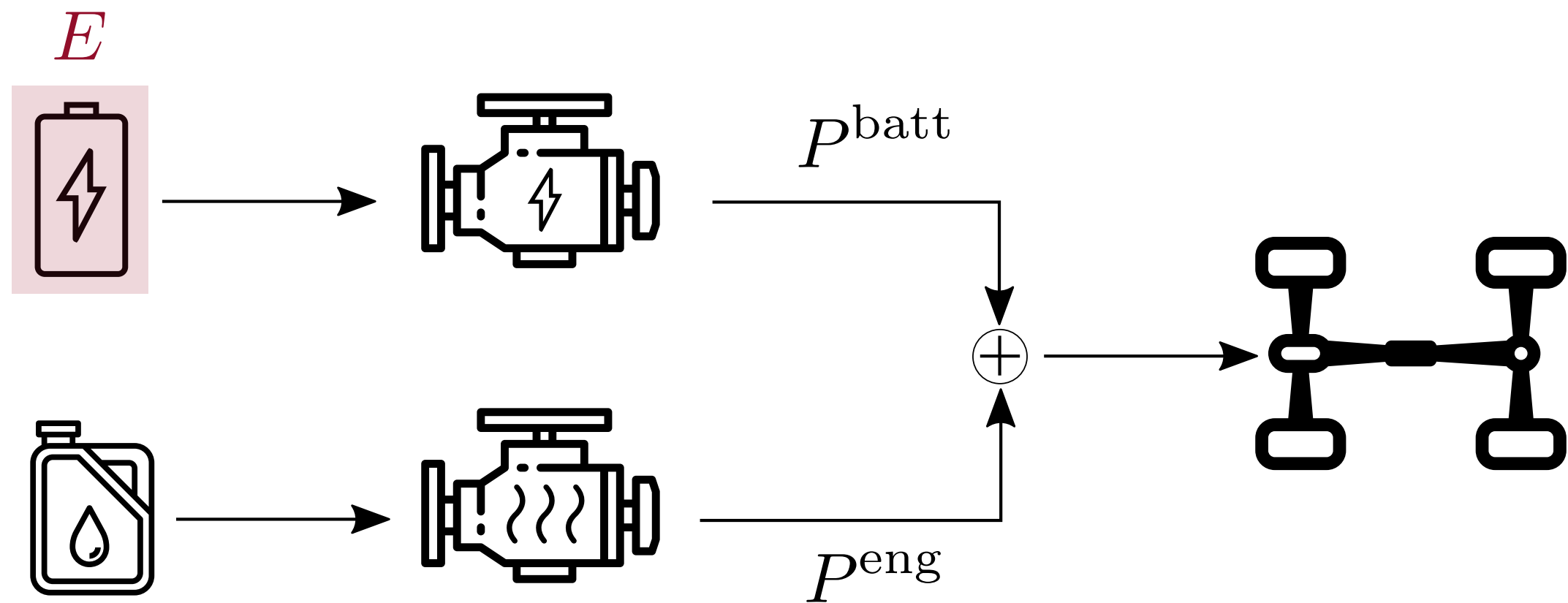
Shipment

Facility

Facility location results



Hybrid-vehicle control



Hybrid-vehicle control

maximize
$$\sum_{t=0}^{T-1} f(P_t^{\text{eng}}, z_t) + \delta(z_t - z_{t-1})_+$$

subject to

$$P_t^{\text{batt}} + P_t^{\text{eng}} \geq P_t^{\text{des}}$$

Desired Power

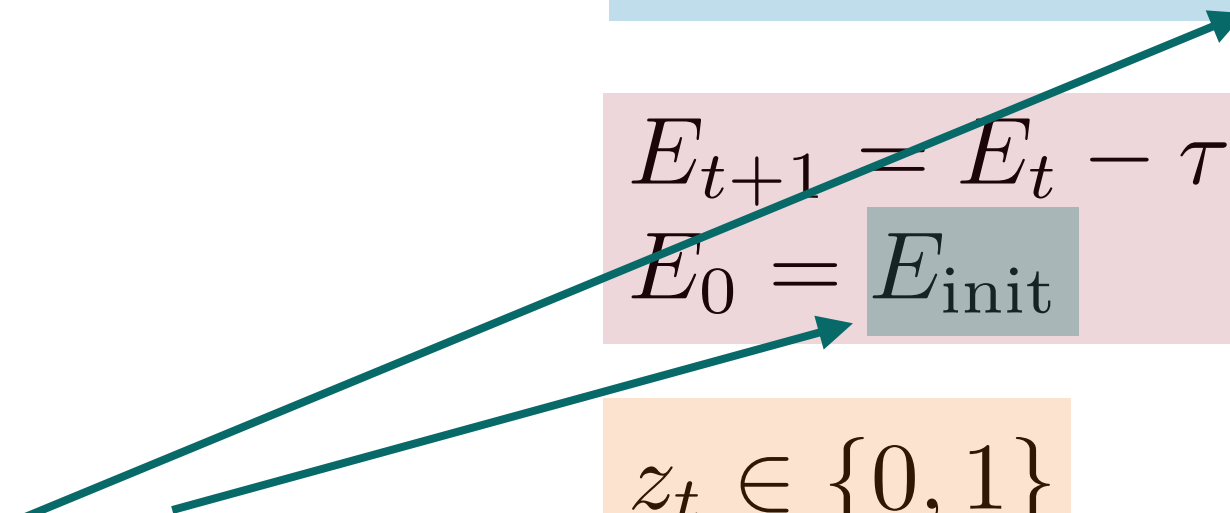
$$E_{t+1} = E_t - \tau P_t^{\text{batt}}$$
$$E_0 = E_{\text{init}}$$

Battery
Dynamics

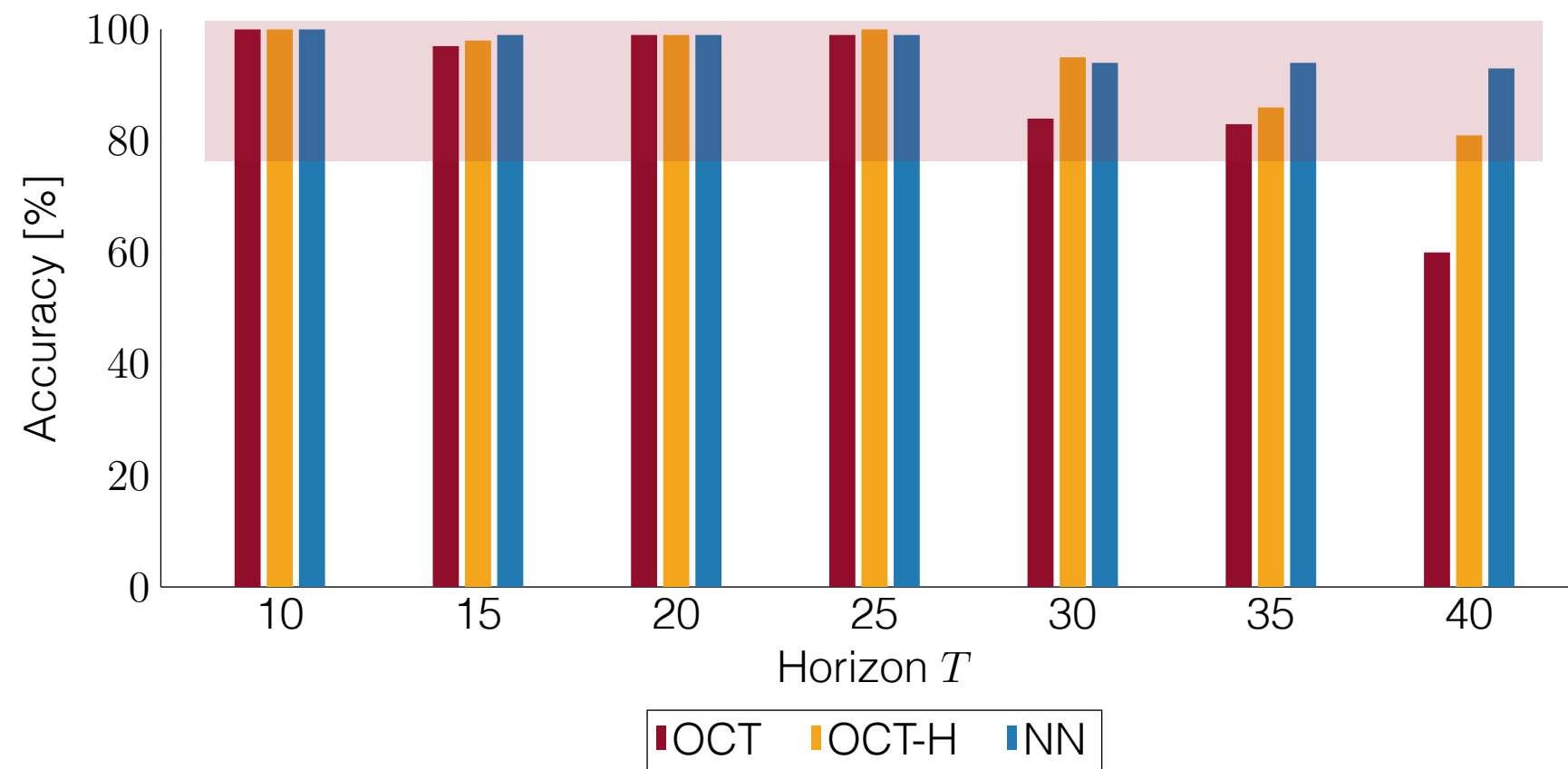
$$z_t \in \{0, 1\}$$

Engine Switch

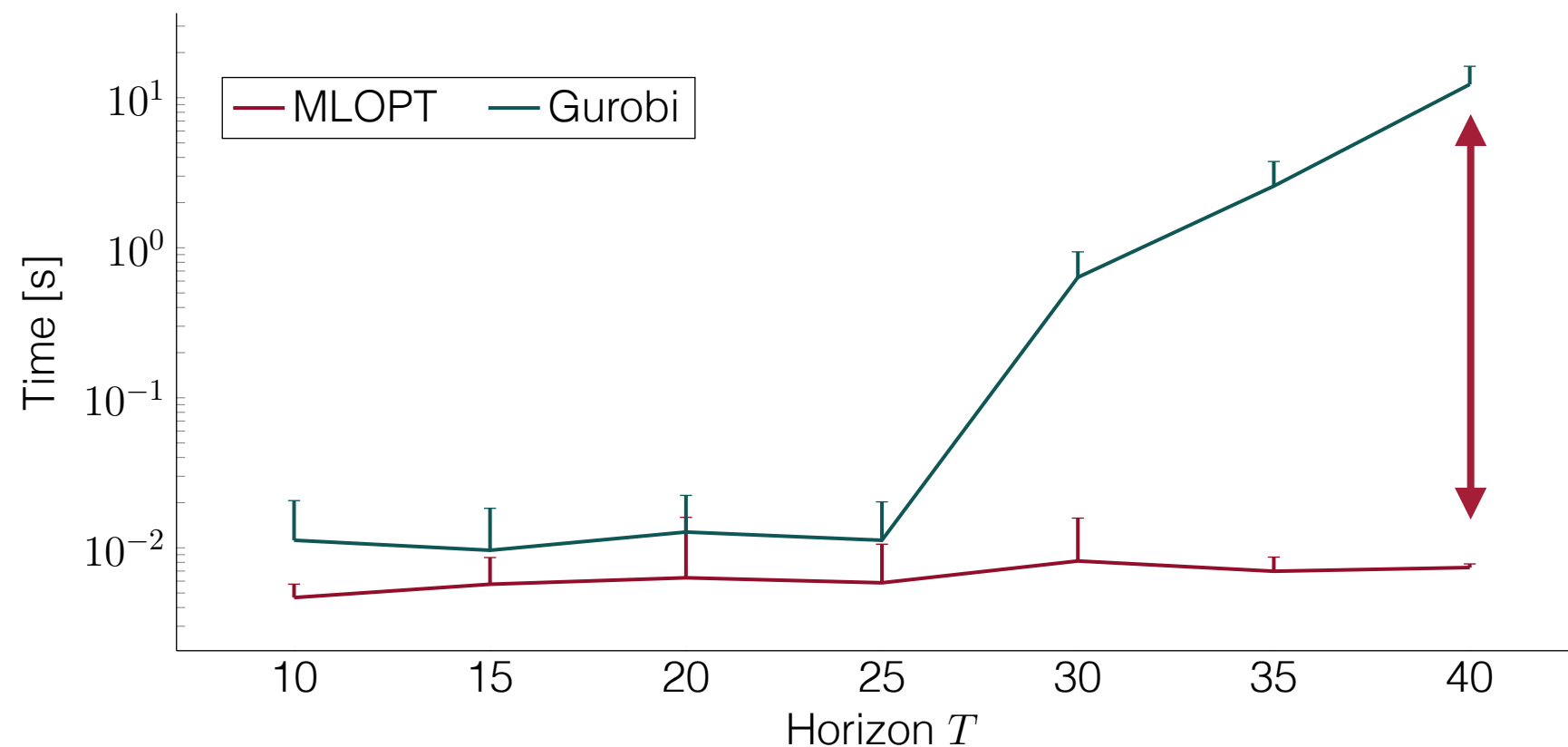
Parameters



Model-predictive control results

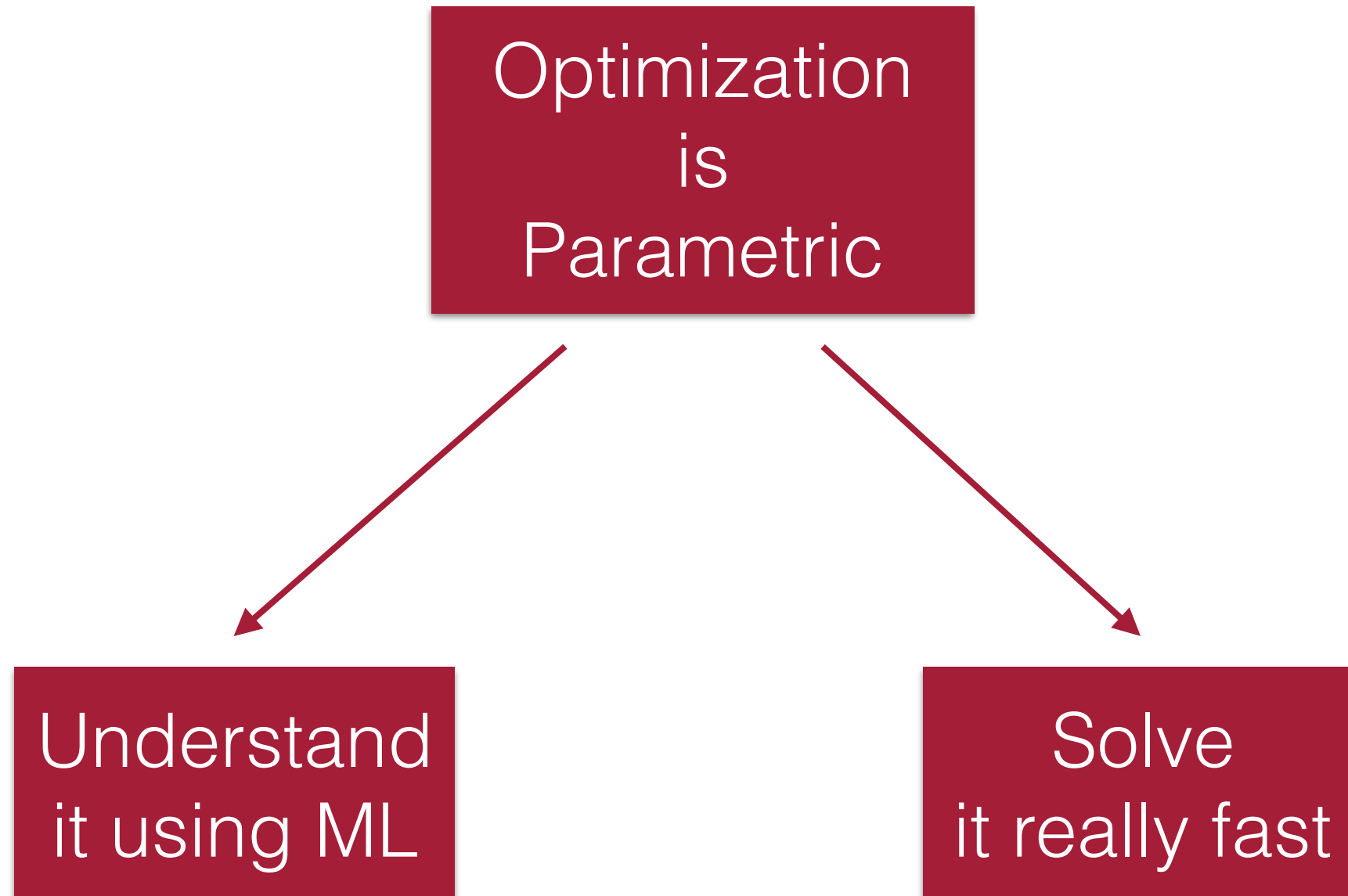


High
accuracy



1000x
faster

Conclusions



Bertsimas, D., Stellato, B., *The Voice of Optimization*, arXiv:1812.09991, 2018

Backup

Good-Turing estimator

Approximation

$$GT = \frac{N_1}{N} \approx \mathbf{P}(\theta_{N+1} \in \mathbf{R}^p \mid s(\theta_{N+1}) \notin \mathcal{S}(\Theta_N))$$

Bound

$$\mathbf{P}(\theta_{N+1} \in \mathbf{R}^p \mid s(\theta_{N+1}) \notin \mathcal{S}(\Theta_N)) \leq GT + c\sqrt{(1/N) \ln(3/\beta)}$$

Sampling scheme

Algorithm 1 Strategies exploration

```
1: given  $\epsilon, \beta, \Theta = \emptyset, \mathcal{S} = \emptyset, u = \infty$ 
2: for  $k = 1, \dots$ , do
3:   Sample  $\theta_k$  and compute  $s(\theta_k)$            ▷ Sample parameter and strategy.
4:    $\Theta \leftarrow \Theta \cup \{\theta_k\}$                ▷ Update set of samples.
5:   if  $s(\theta_k) \notin \mathcal{S}$  then
6:      $\mathcal{S} \leftarrow \mathcal{S} \cup \{s(\theta_k)\}$        ▷ Update strategy set if new strategy found
7:   end if
8:   if  $G + c\sqrt{(1/k) \ln(3/\beta)} \leq \epsilon$  then           ▷ Break if bound less than  $\epsilon$ 
9:     break
10:  end if
11: end for
12: return  $k, \Theta, \mathcal{S}$ 
```
